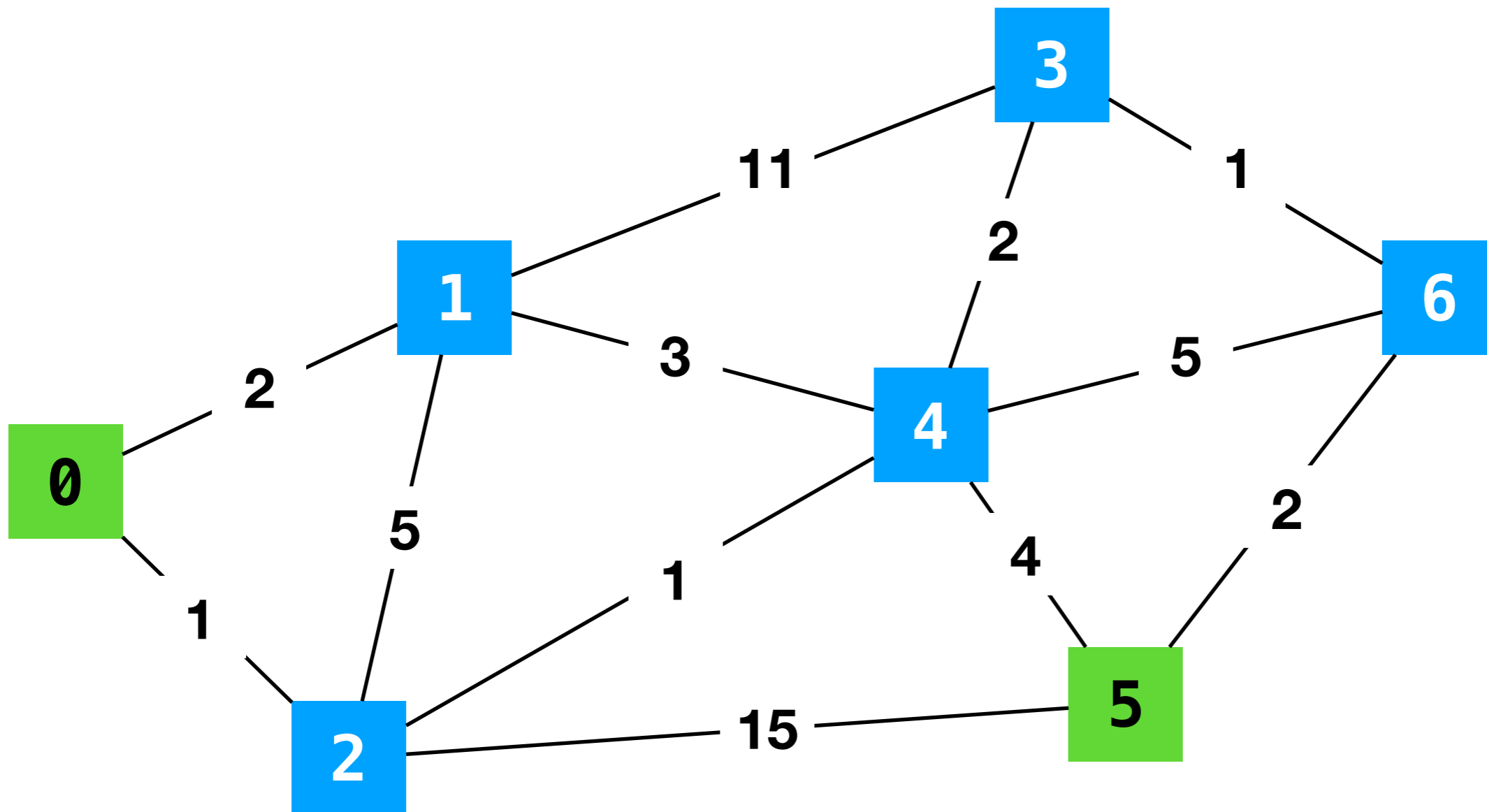


# CS 61BL Lab 17

Ryan Purpura

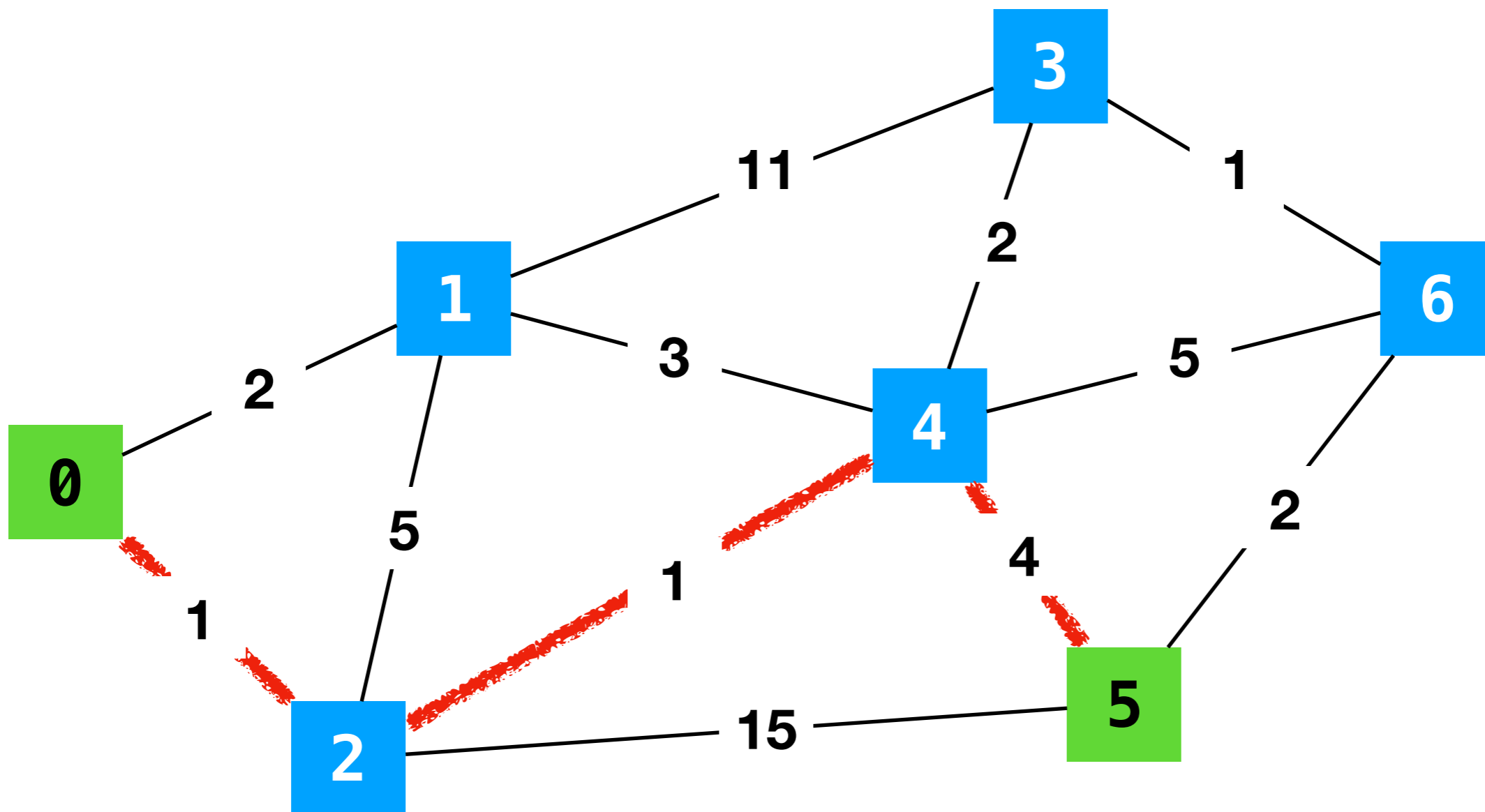
# Best Paths

- What is the best path from node 0 to node 5?



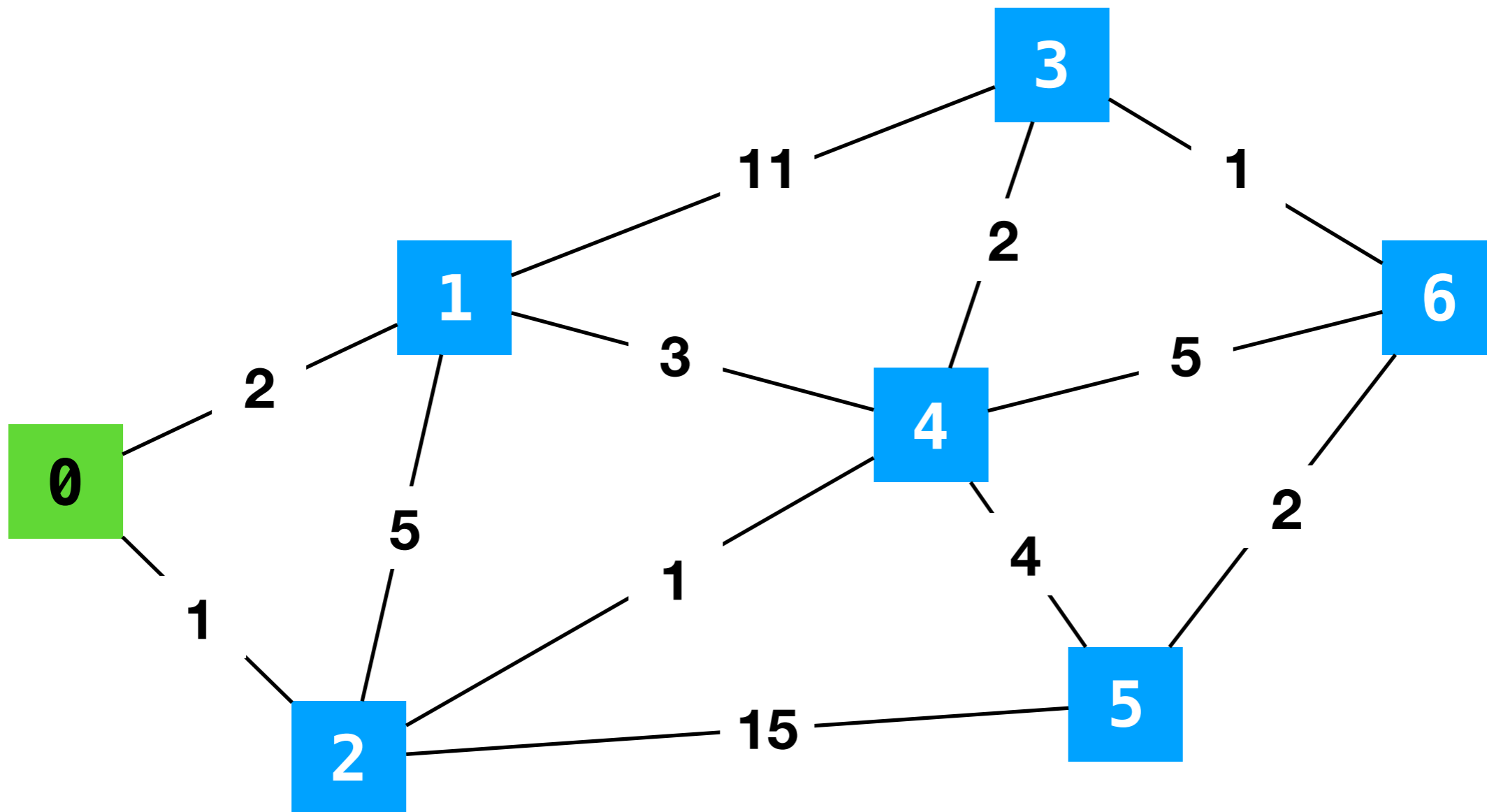
# Best Paths

- What is the best path from node 0 to node 5?



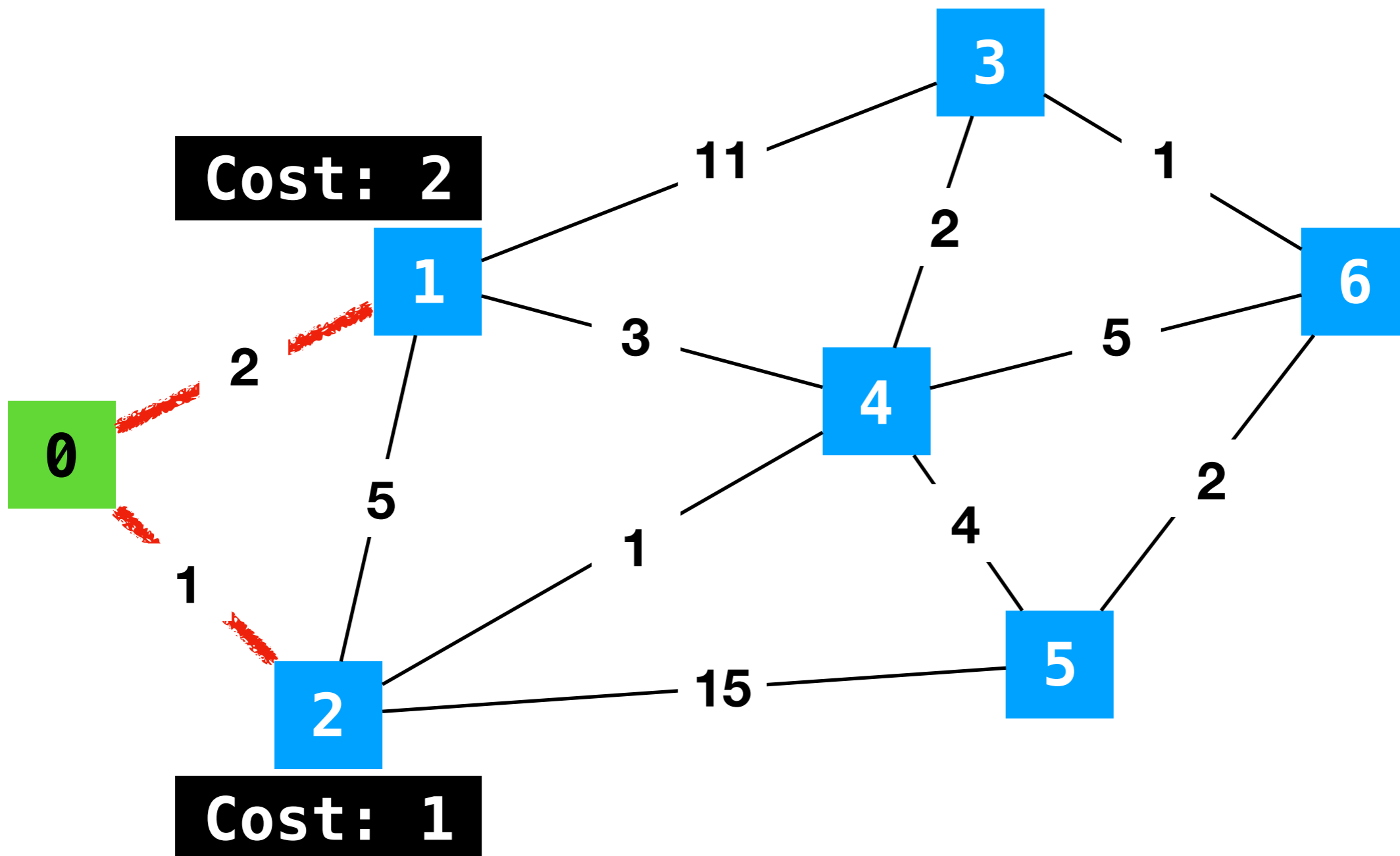
# Best Paths

- What is the best path from node 0 to every other node?



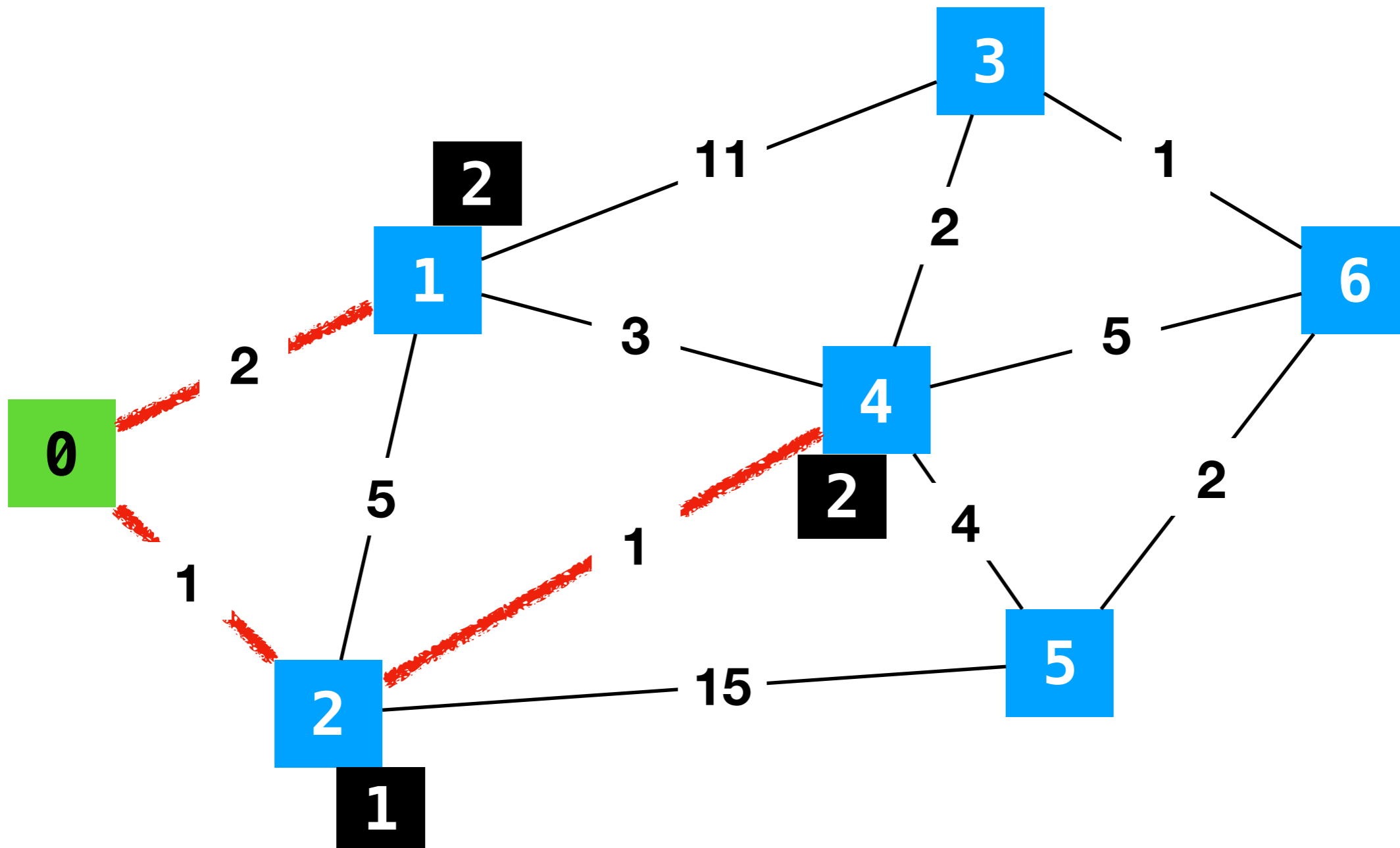
# Best Paths (by inspection)

- What are the best paths from node 0 to every other node?



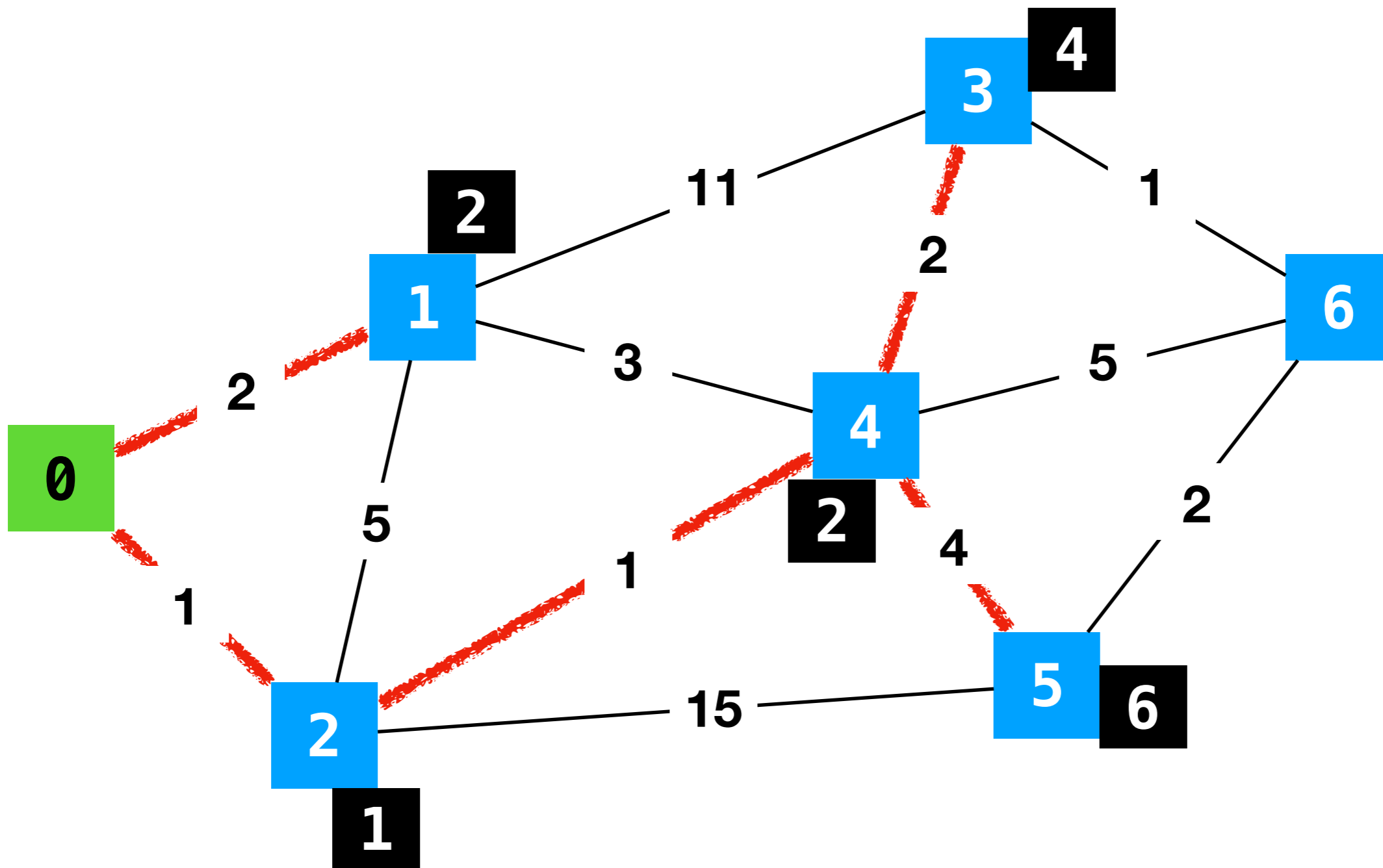
# Best Paths (by inspection)

- What are the best paths from node 0 to every other node?



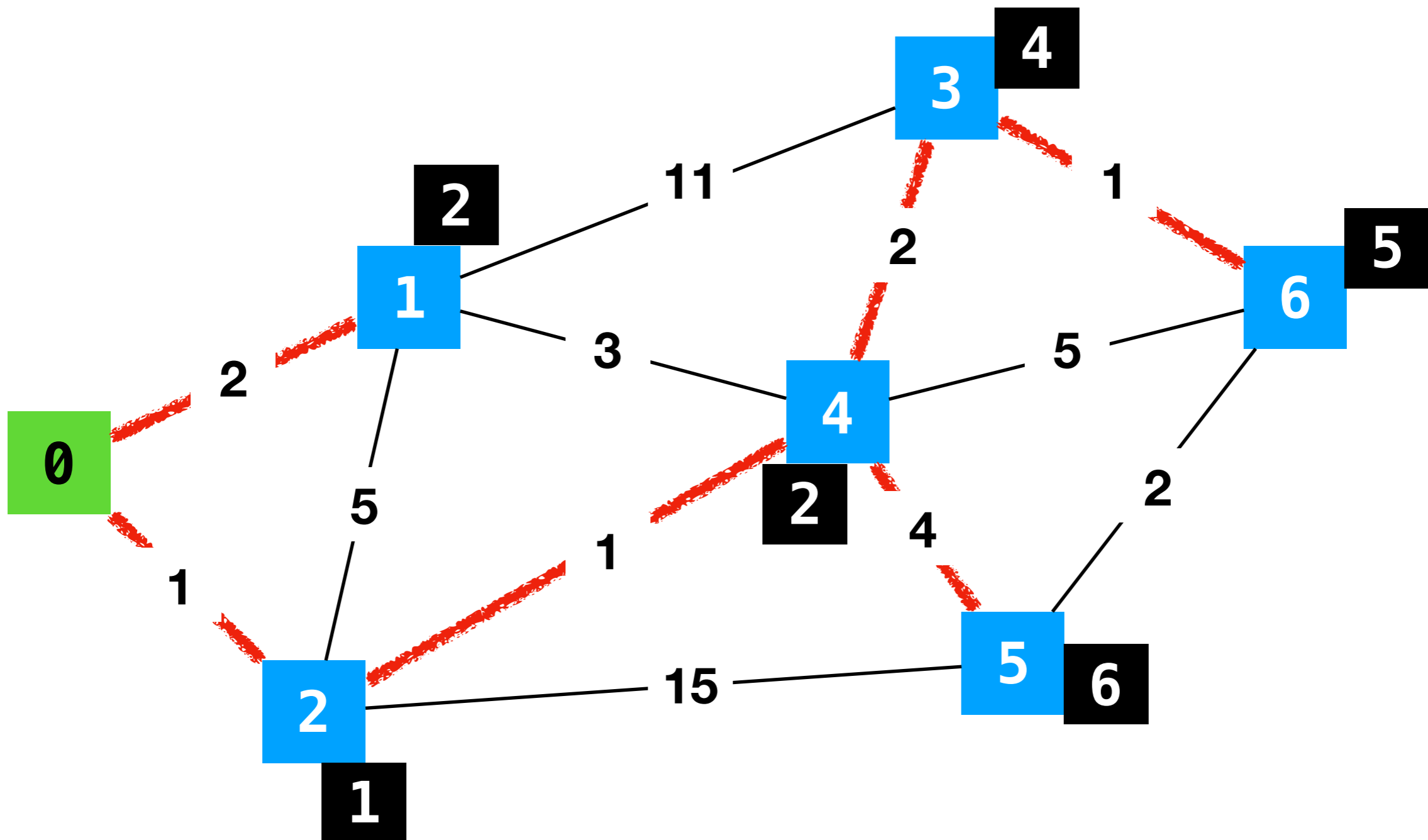
# Best Paths (by inspection)

- What are the best paths from node 0 to every other node?



# Best Paths (by inspection)

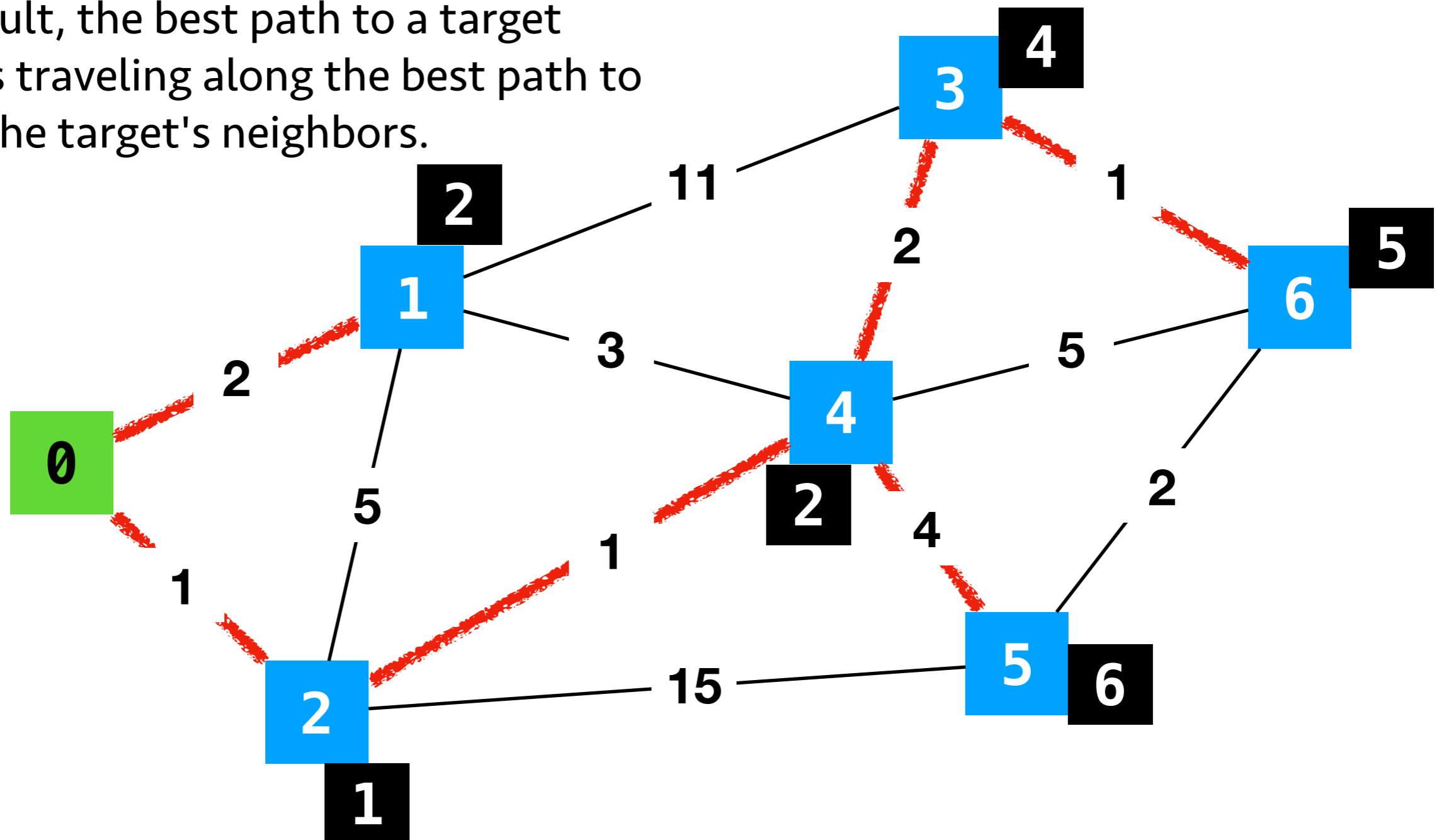
- What are the best paths from node 0 to every other node?



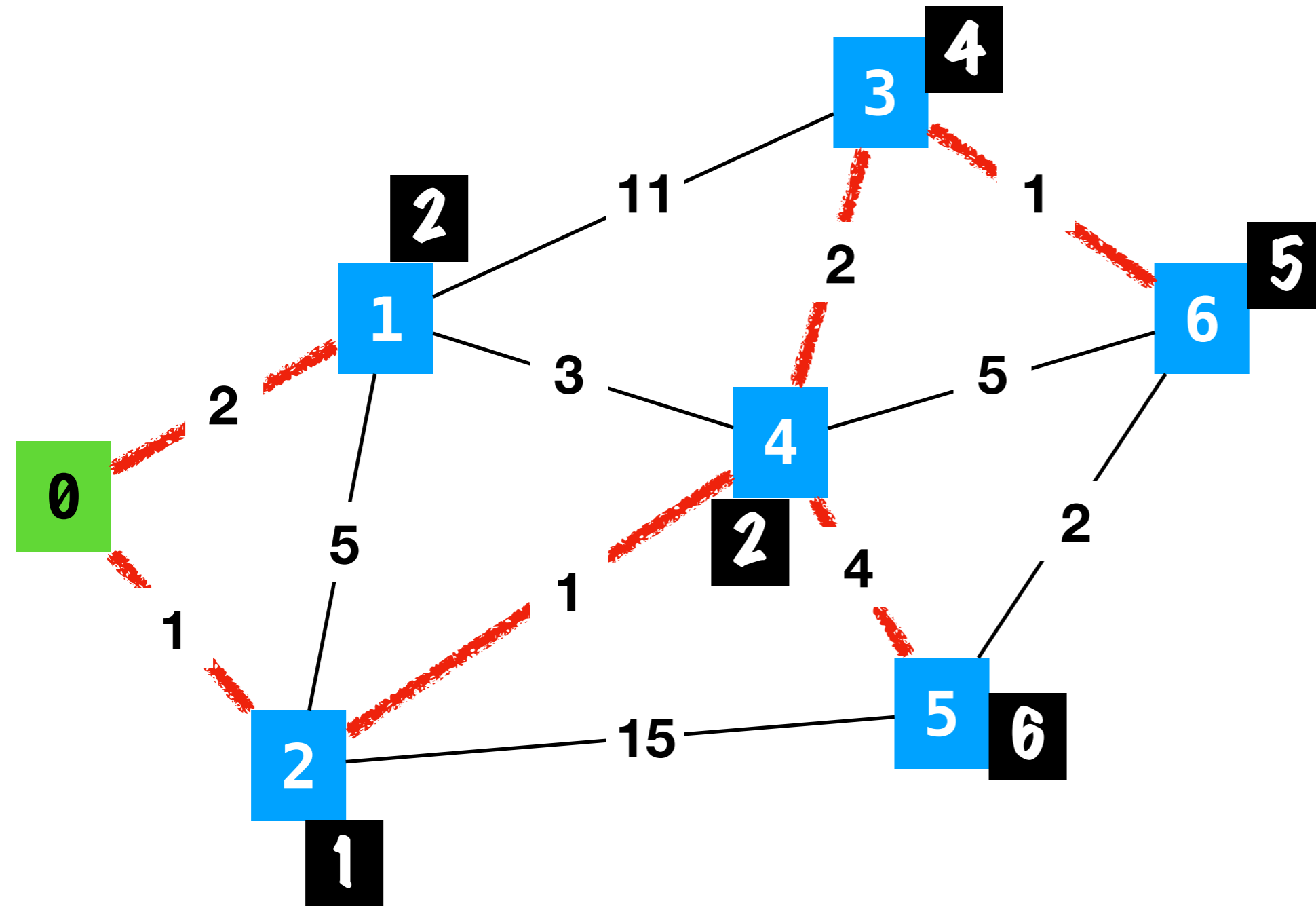


# Observations

- Notice the result is a tree (i.e., no cycles)! Why is that?
- As a result, the best path to a target involves traveling along the best path to one of the target's neighbors.



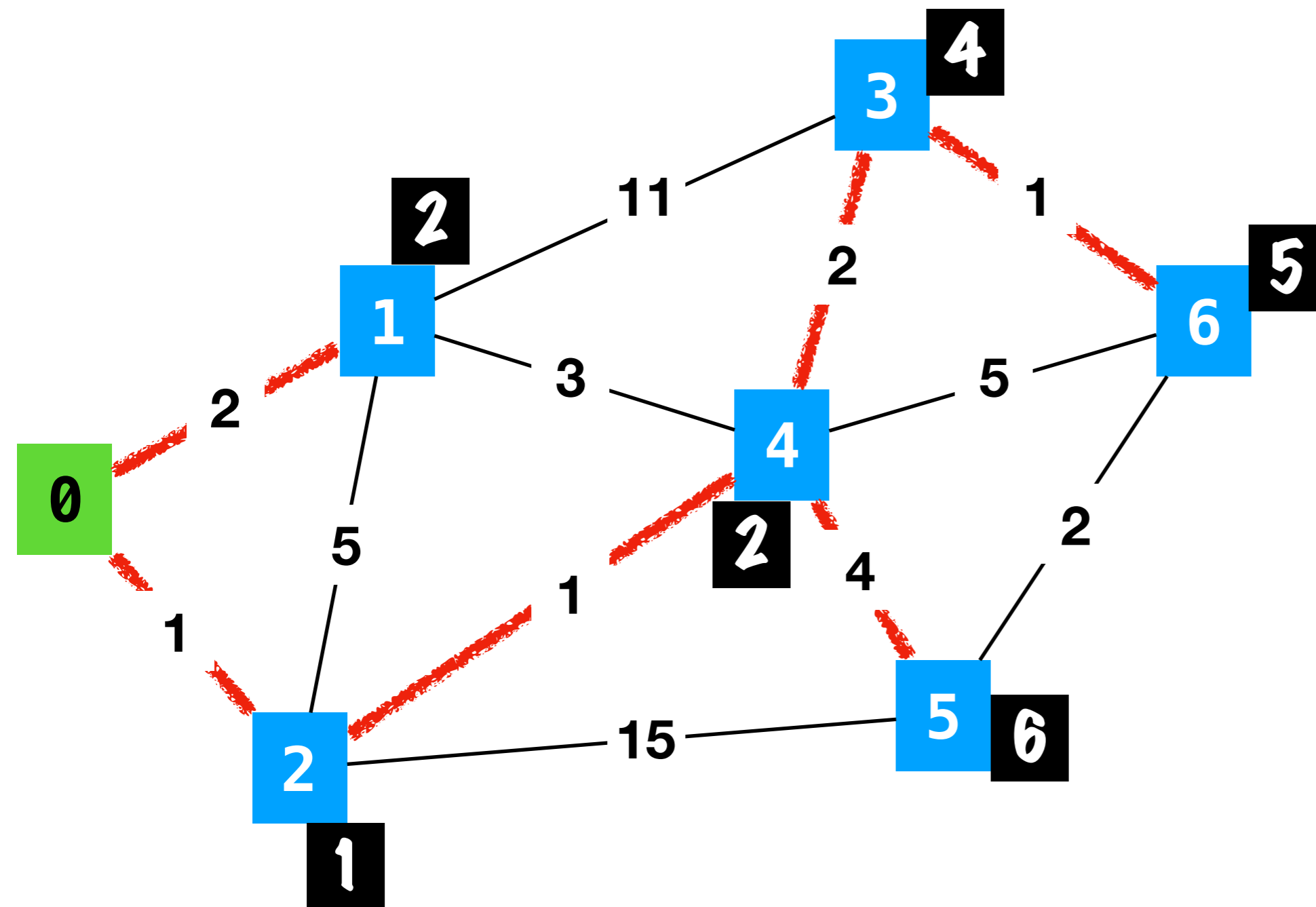
# Representing the Shortest-Path Tree



v	total cost	prev Node
0	0	--
1	2	0
2	1	0
3	4	4
4	2	2
5	6	4
6	5	3

# Our goal

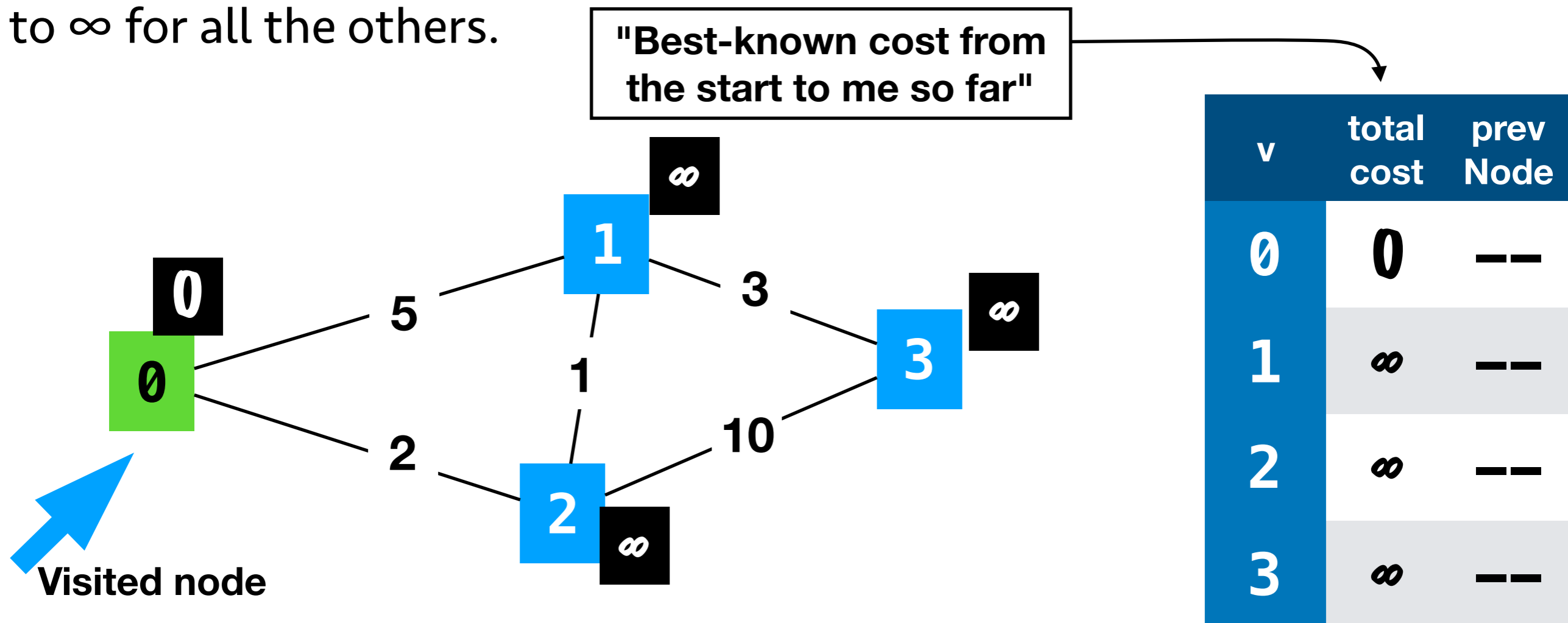
- We want to algorithmically generate the shortest-path tree.



v	total cost	prev Node
0	0	--
1	2	0
2	1	0
3	4	4
4	2	2
5	6	4
6	5	3

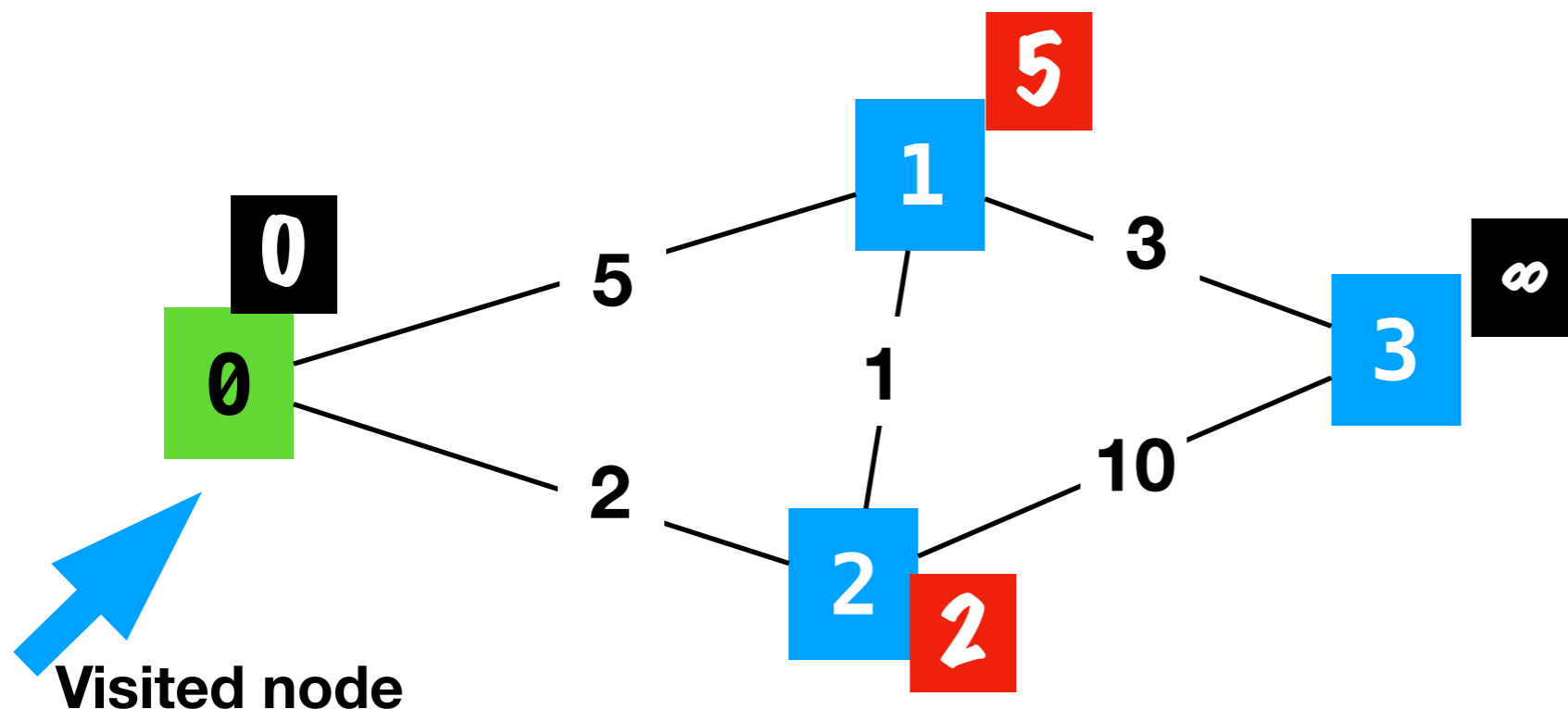
# The Idea: Dijkstra's Algorithm

- Visit nodes in closest-first order, beginning at the start node.
- When we visit a node, we will update its neighbors' costs and prevNode if going through the visited node results in a cheaper path.
- The starting cost should be initialized to 0 for the starting node and to  $\infty$  for all the others.



# The Idea

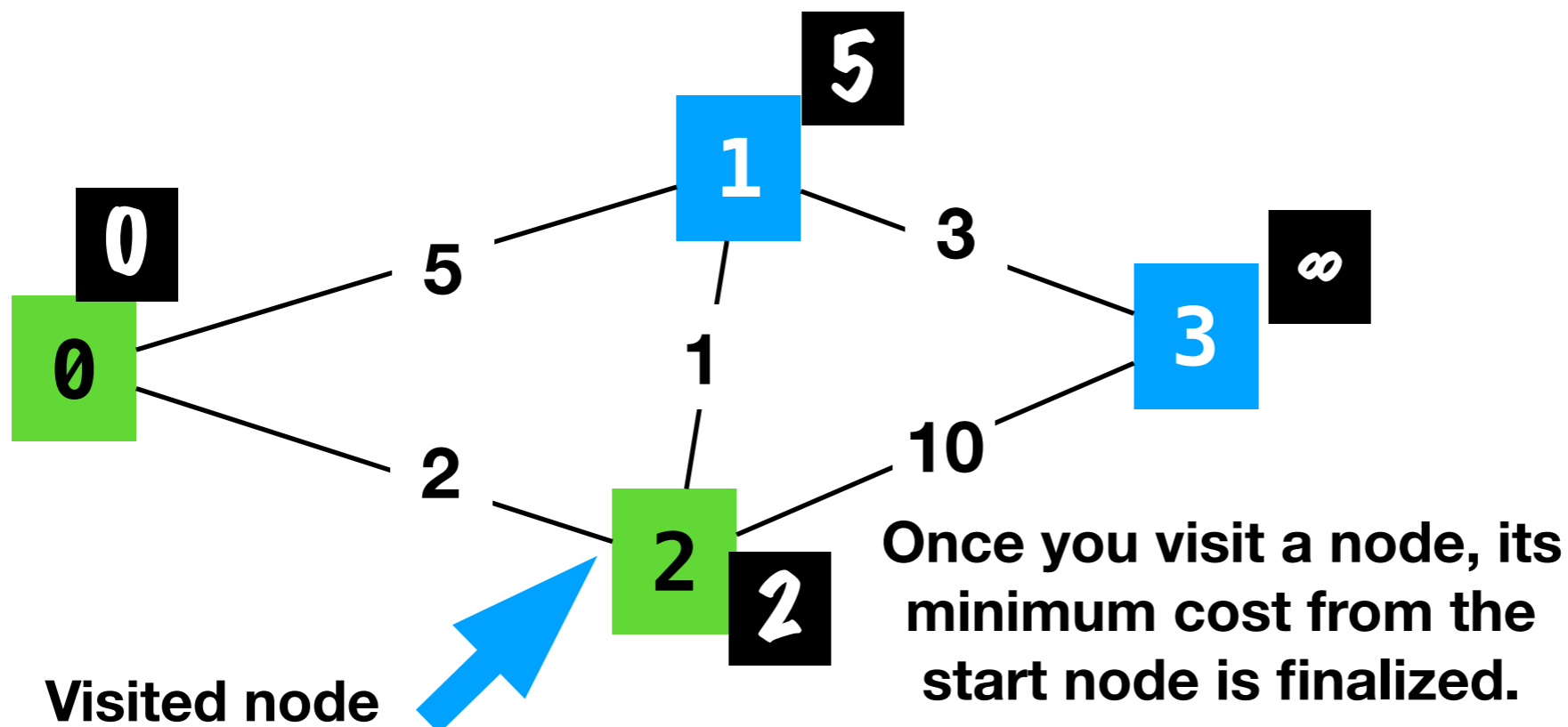
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- When we visit a node, we will update its neighbors' costs and prevNode if going through the visited node results in a cheaper path.
- The starting cost should be initialized to 0 for the starting node and to  $\infty$  for all the others.



v	total cost	prev Node
0	0	--
1	5	0
2	2	0
3	$\infty$	--

# The Idea

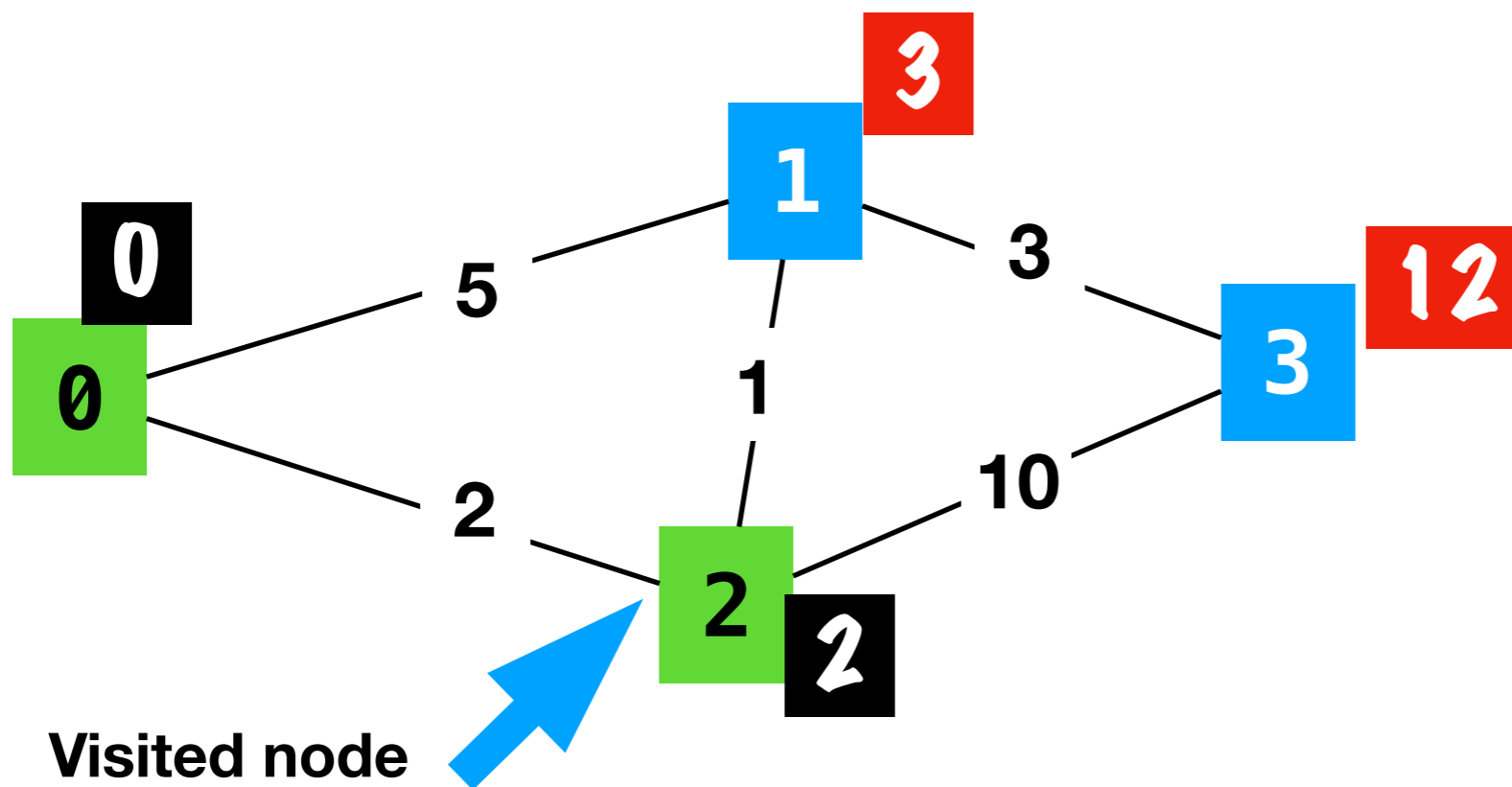
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v	total cost	prev Node
0	0	--
1	5	0
2	2	0
3	$\infty$	--

# The Idea

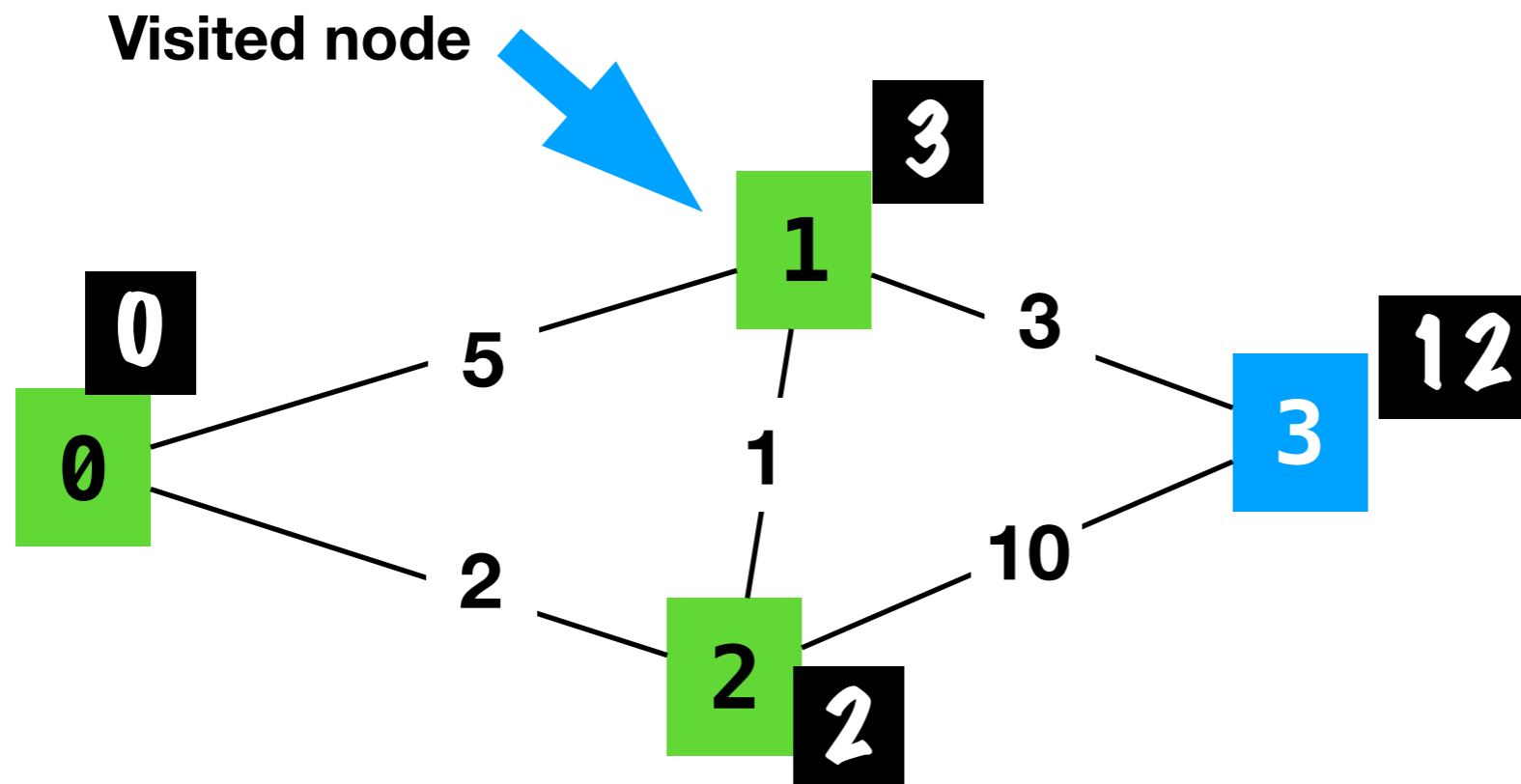
- Visit nodes in closest-first order, starting at the start node.
- When we visit a node, we will update its neighbors' costs and prevNode if going through the visited node results in a cheaper path.
- The starting cost should be initialized to 0 for the starting node and to  $\infty$  for all the others.



v	total cost	prev Node
0	0	--
1	3	2
2	2	0
3	12	2

# The Idea

- Visit nodes in closest-first order, starting at the start node.
- When we visit a node, we will update its neighbors' costs and prevNode if going through the visited node results in a cheaper path.
- The starting cost should be initialized to 0 for the starting node and to  $\infty$  for all the others.

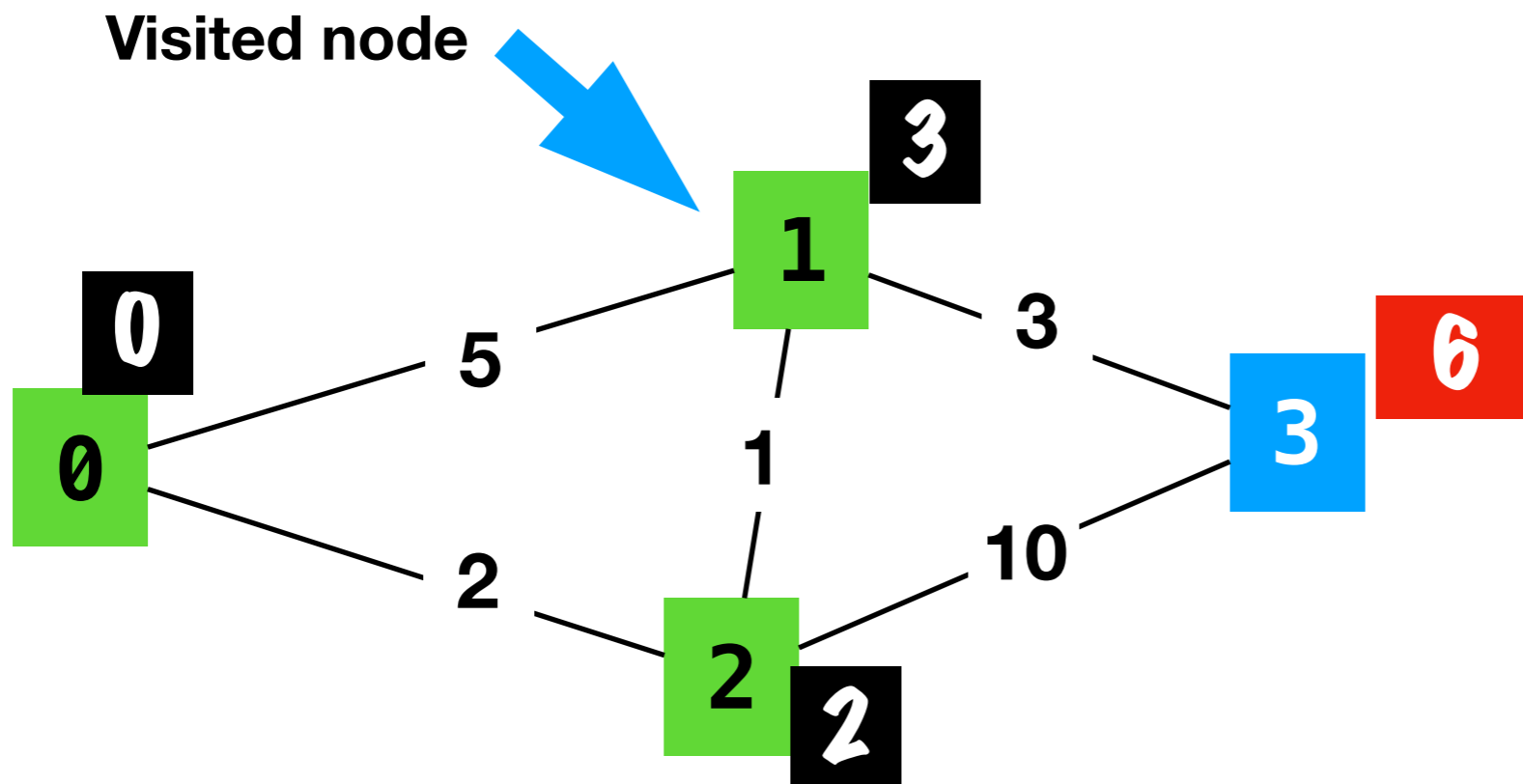


v	total cost	prev Node
0	0	--
1	3	2
2	2	0
3	12	2



# The Idea

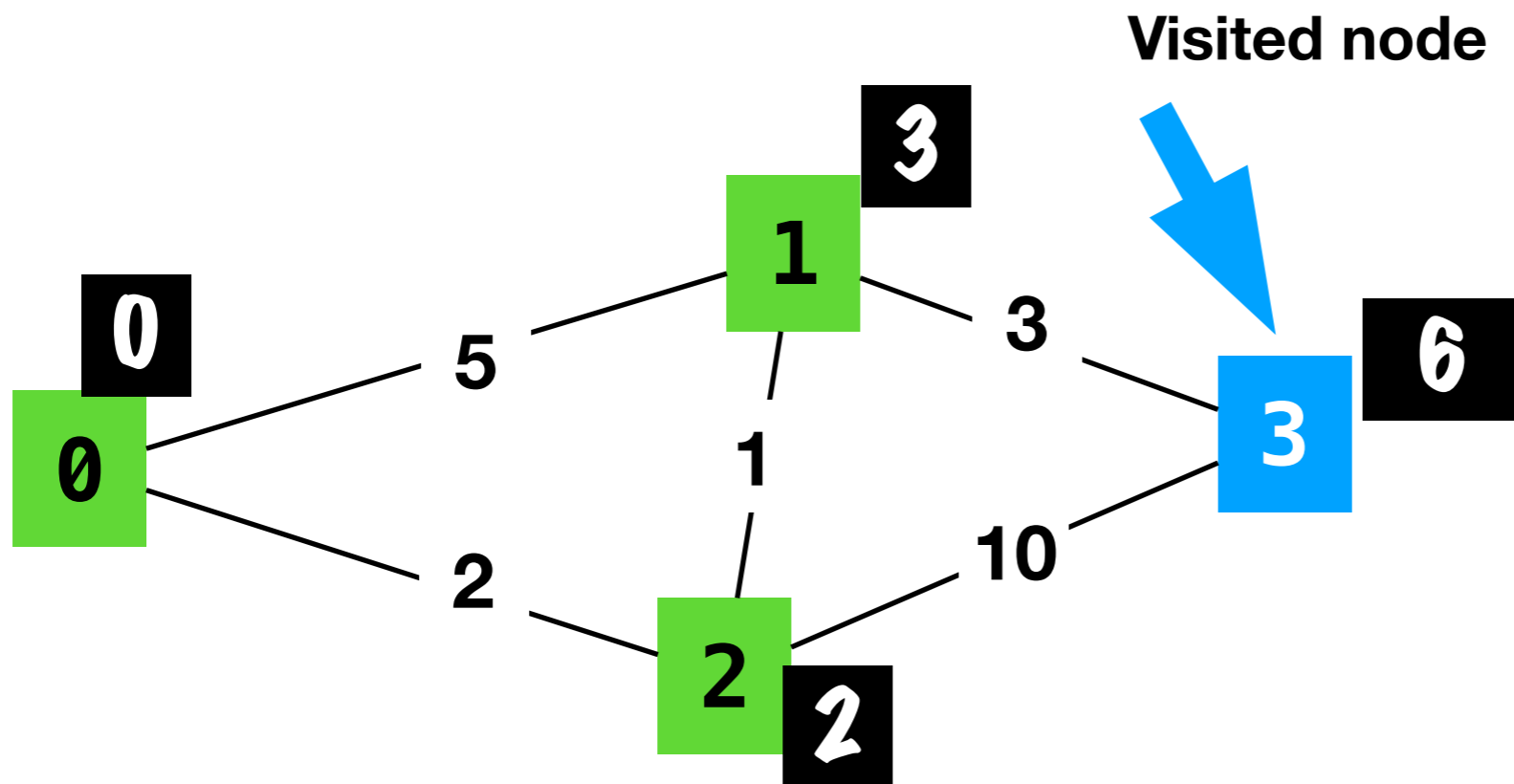
- Visit nodes in closest-first order, starting at the start node.
- When we visit a node, we will update its neighbors' costs and prevNode if going through the visited node results in a cheaper path.
- The starting cost should be initialized to 0 for the starting node and to  $\infty$  for all the others.



v	total cost	prev Node
0	0	--
1	3	2
2	2	0
3	6	1

# The Idea

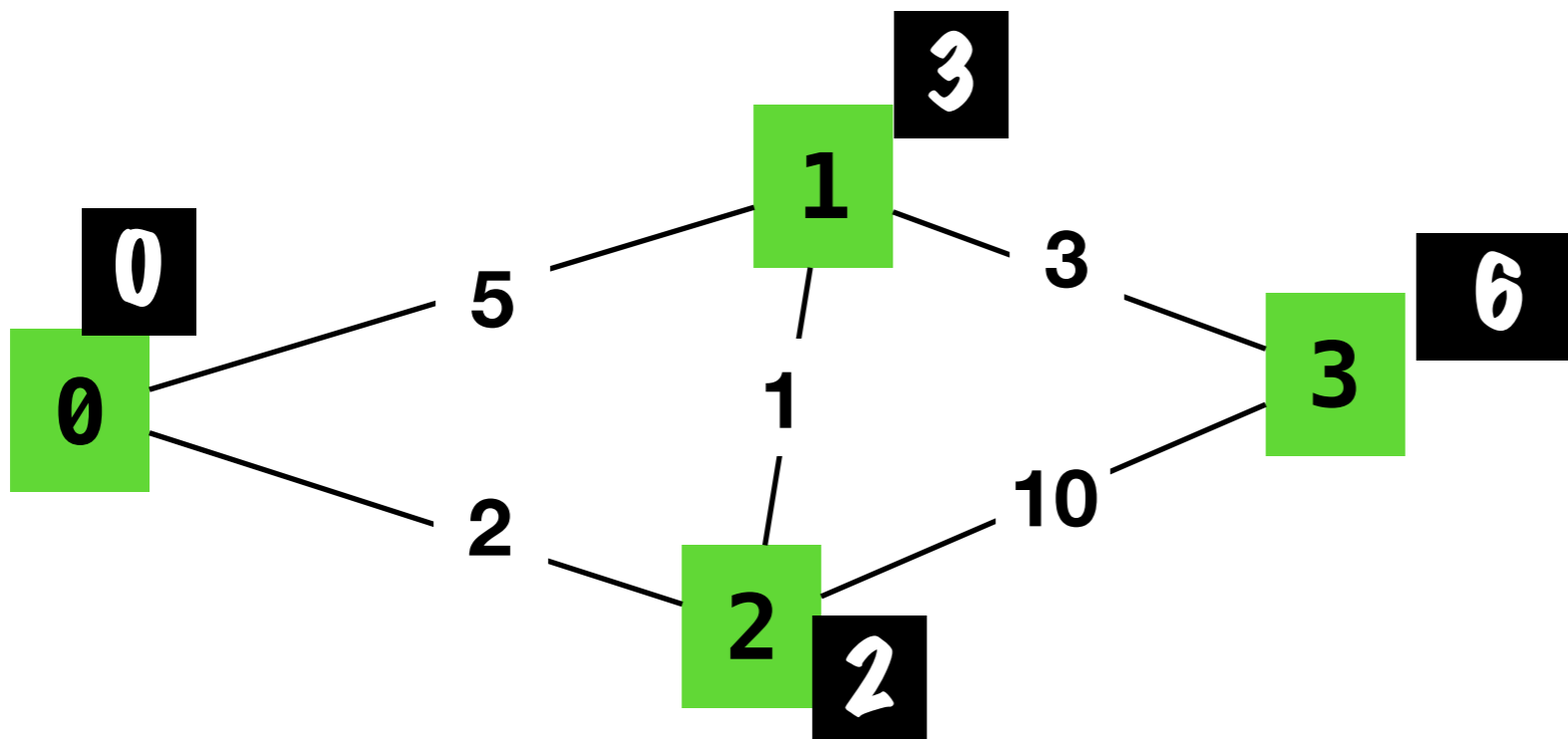
- Visit nodes in closest-first order, starting at the start node.
- When we visit a node, we will update its neighbors' costs and prevNode if going through the visited node results in a cheaper path.
- The starting cost should be initialized to 0 for the starting node and to  $\infty$  for all the others.



v	total cost	prev Node
0	0	--
1	3	2
2	2	0
3	6	1

# The Idea

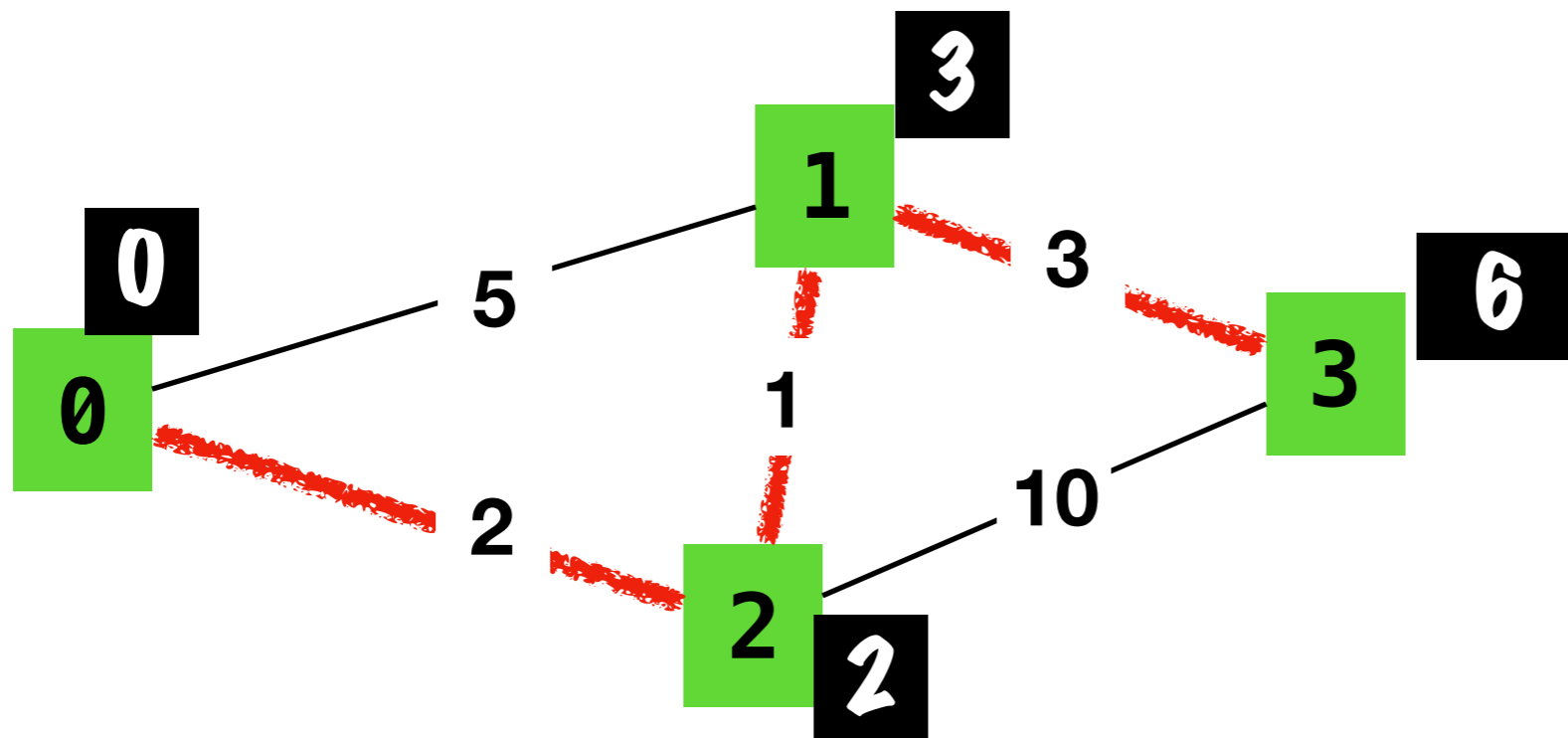
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v	total cost	prev Node
0	0	--
1	3	2
2	2	0
3	6	1

# The Idea

- Visit nodes in closest-first order, starting at the start node.
- When we visit a node, we will update its neighbors' costs and prevNode if going through the visited node results in a cheaper path.
- The starting cost should be initialized to 0 for the starting node and to  $\infty$  for all the others.

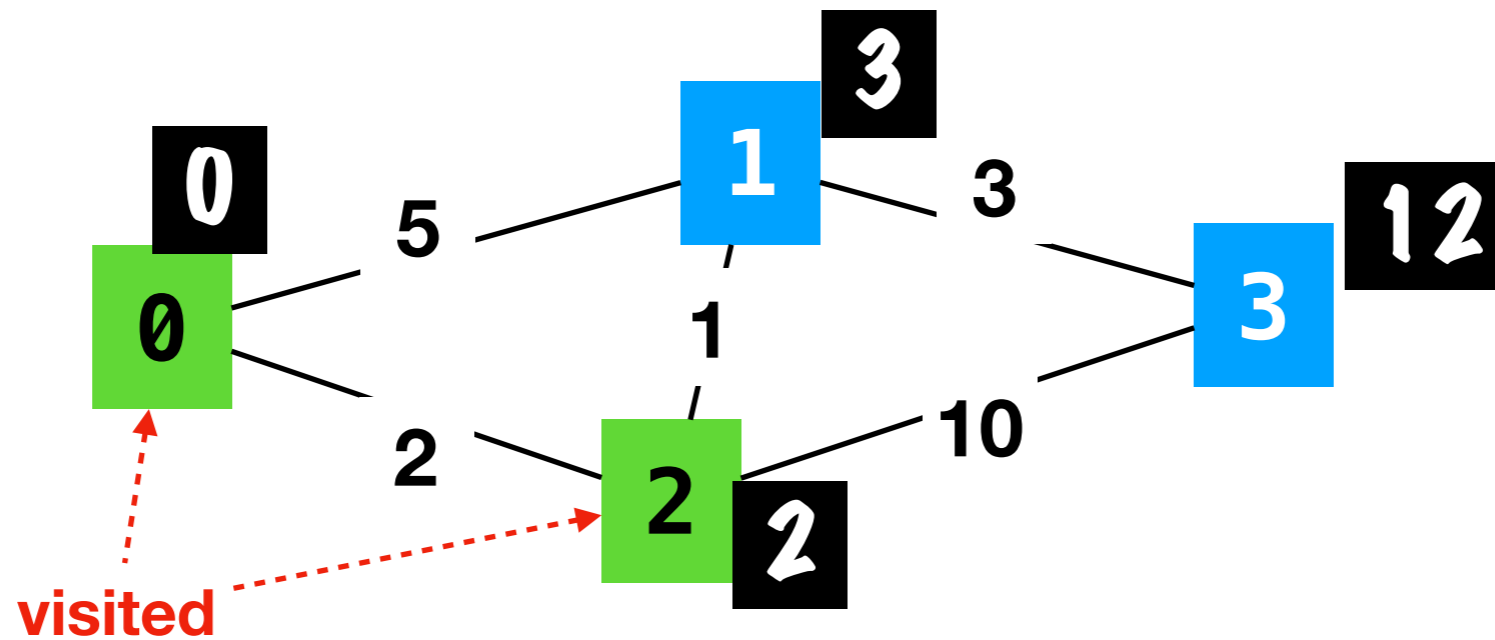


v	total cost	prev Node
0	0	--
1	3	2
2	2	0
3	6	1

# The Fringe

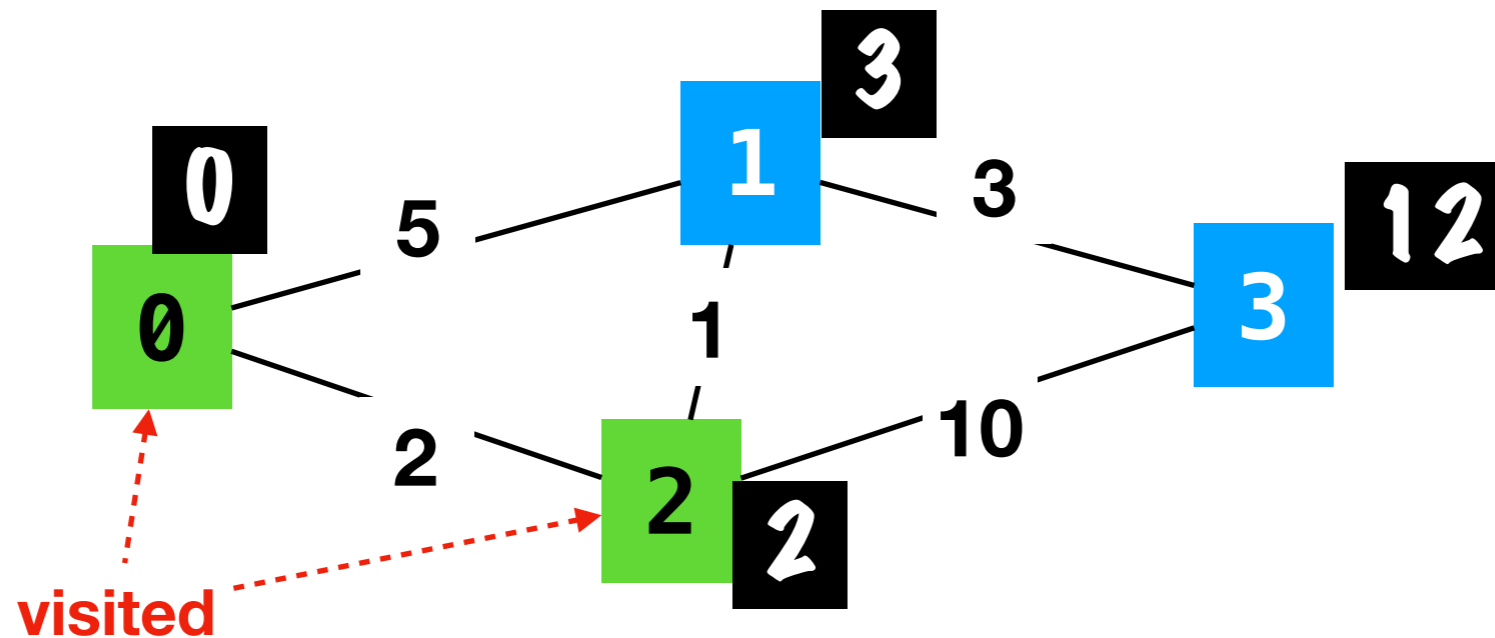
# The Fringe

- How to get which unvisited node is closest to the start?



# The Fringe

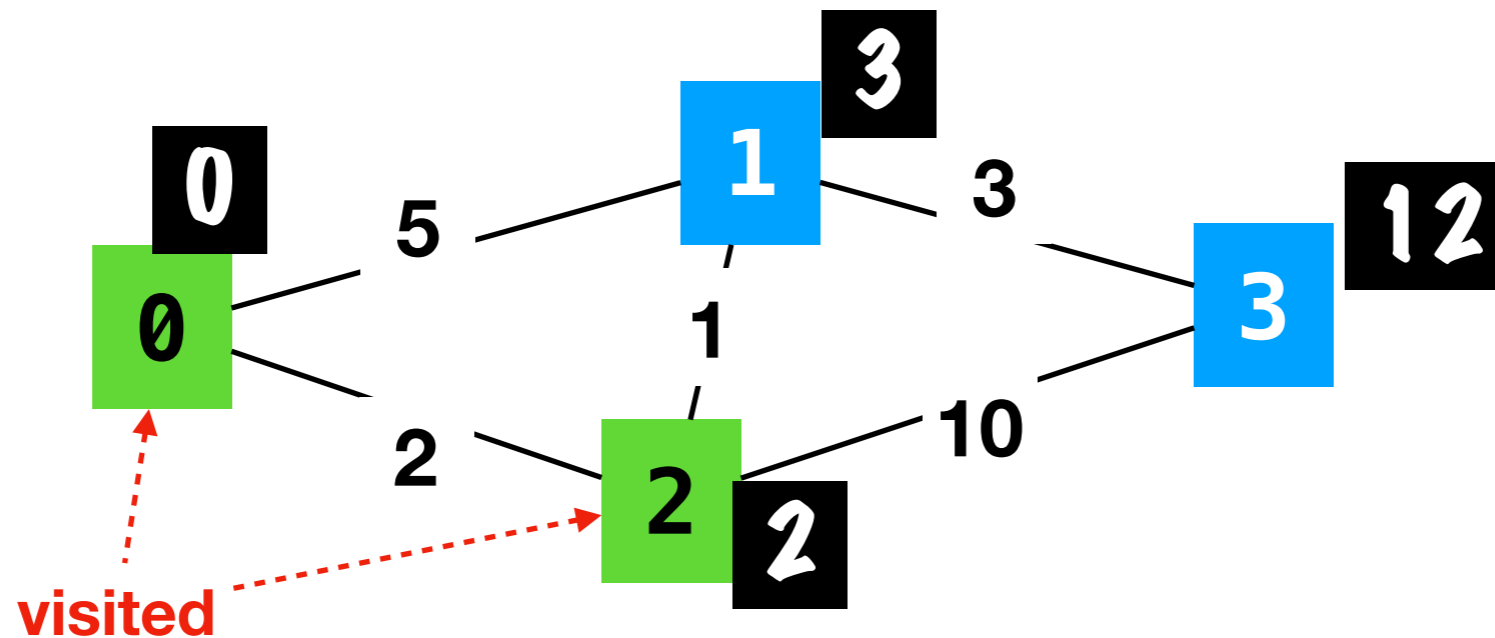
- How to get which unvisited node is closest to the start?



- This is handled by the fringe, which will contain all unvisited nodes and will give us the node to visit next.

# The Fringe

- How to get which unvisited node is closest to the start?

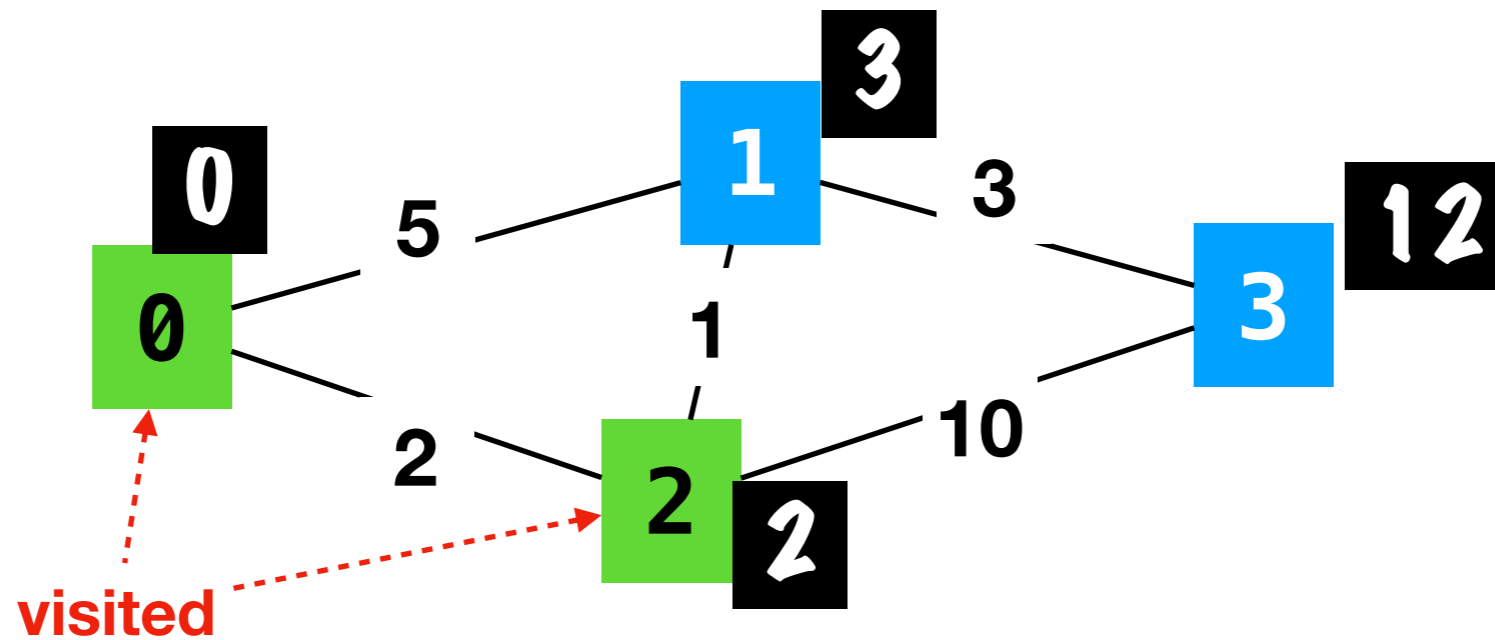


- This is handled by the fringe, which will contain all unvisited nodes and will give us the node to visit next.
- What abstract data type should our fringe be?



# The Fringe

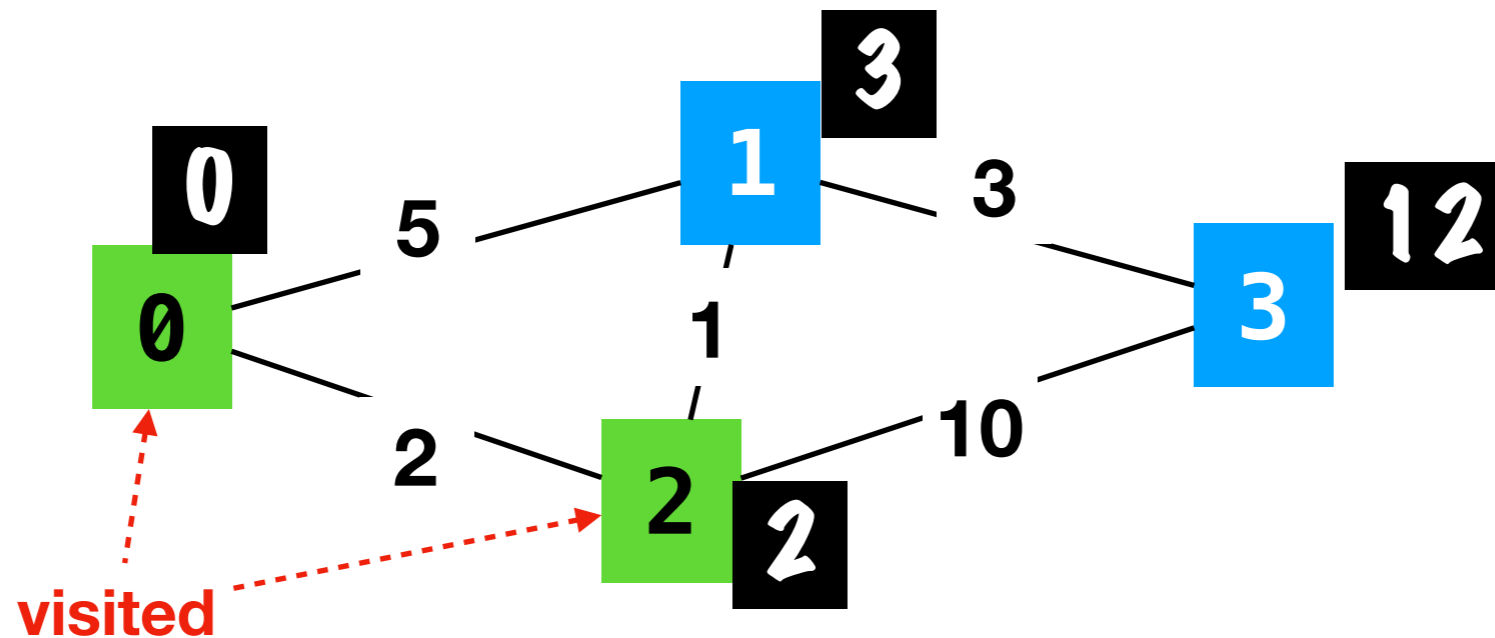
- How to get which unvisited node is closest to the start?



- This is handled by the fringe, which will contain all unvisited nodes and will give us the node to visit next.
- What abstract data type should our fringe be?
- (Min) priority queue, with priority equal to the best-known cost so far.

# The Fringe

- How to get which unvisited node is closest to the start?



- This is handled by the fringe, which will contain all unvisited nodes and will give us the node to visit next.
- What abstract data type should our fringe be?
- (Min) priority queue, with priority equal to the best-known cost so far.

$$PQ = \{ (1, 3), (3, 12) \}$$

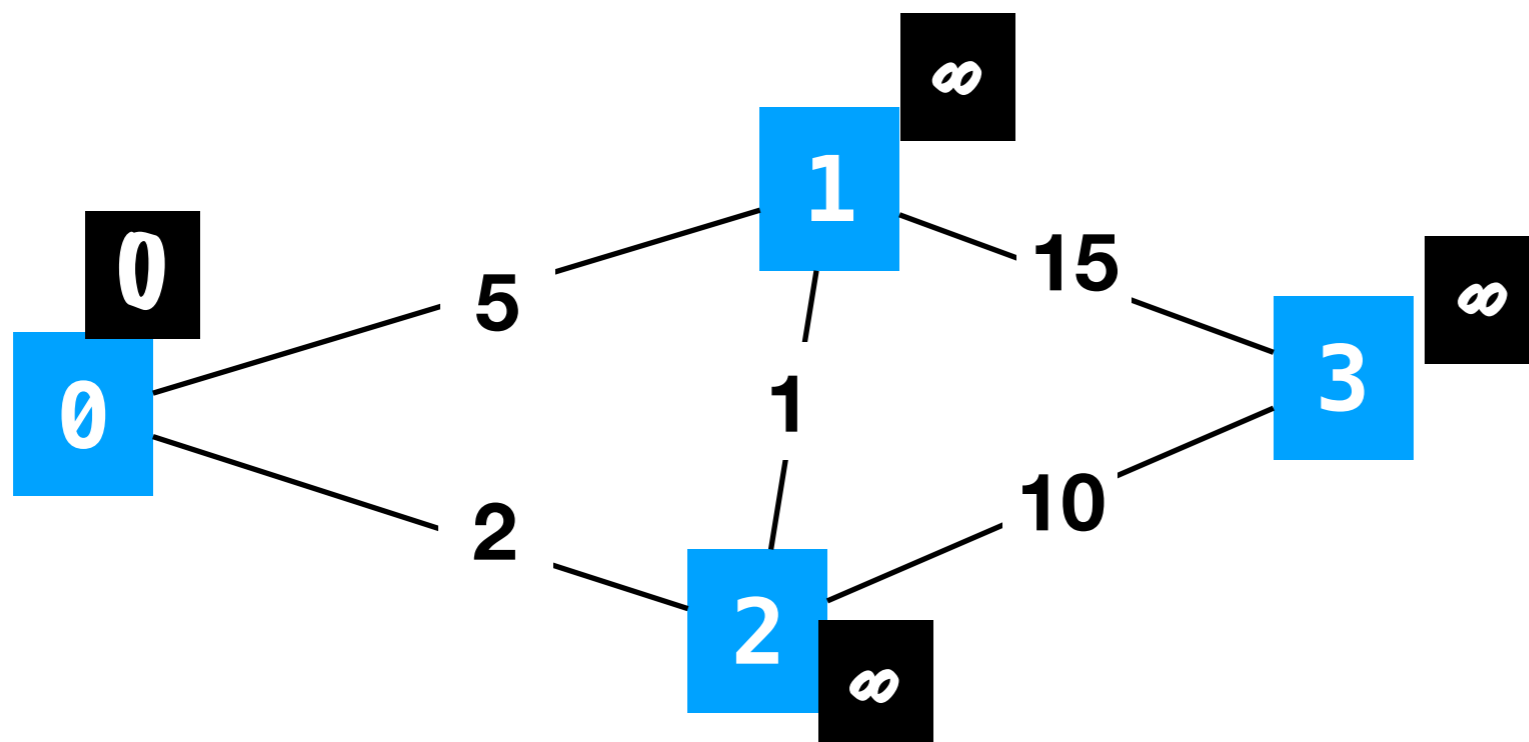
# The Fringe: The Specifics

- Recall the start node begins with a cost of 0, while all other nodes begin with a "best-known" cost of  $\infty$ .
  - Therefore, we will insert all of our nodes into our fringe with the start node with priority 0 and all others with priority  $\infty$ .
- If we find a cheaper path to a node's neighbor, we need to update that neighbor's priority in the fringe.

# The Idea

- First, we initialize our fringe.

fringe = { ( 0 , 0 ), ( 1 , ∞ ), ( 2 , ∞ ), ( 3 , ∞ ) }

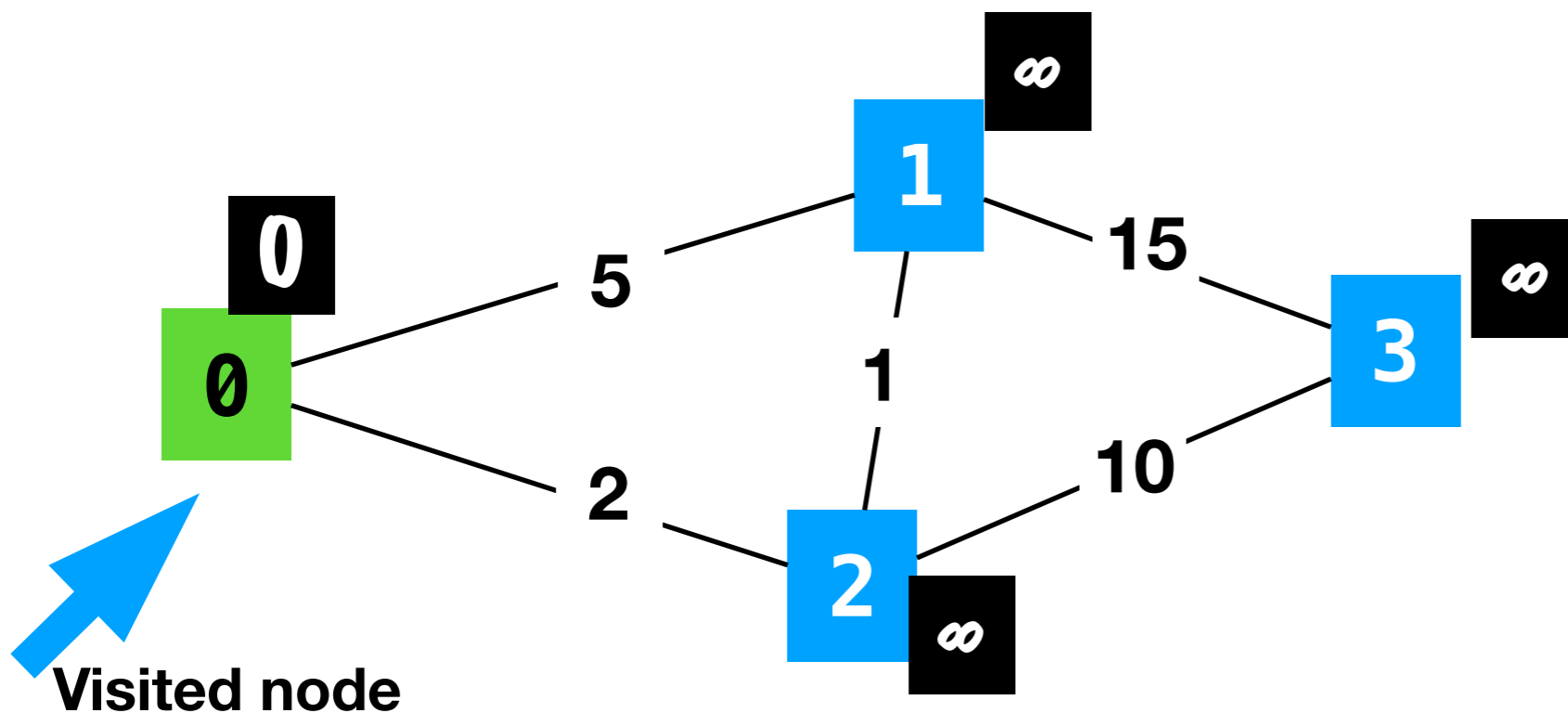


v	total cost	prev Node
0	0	--
1	∞	--
2	∞	--
3	∞	--

# The Idea

- Remove the minimum item from the fringe. This is our current node.

fringe = { ( 1,  $\infty$  ), ( 2,  $\infty$  ), ( 3,  $\infty$  ) }

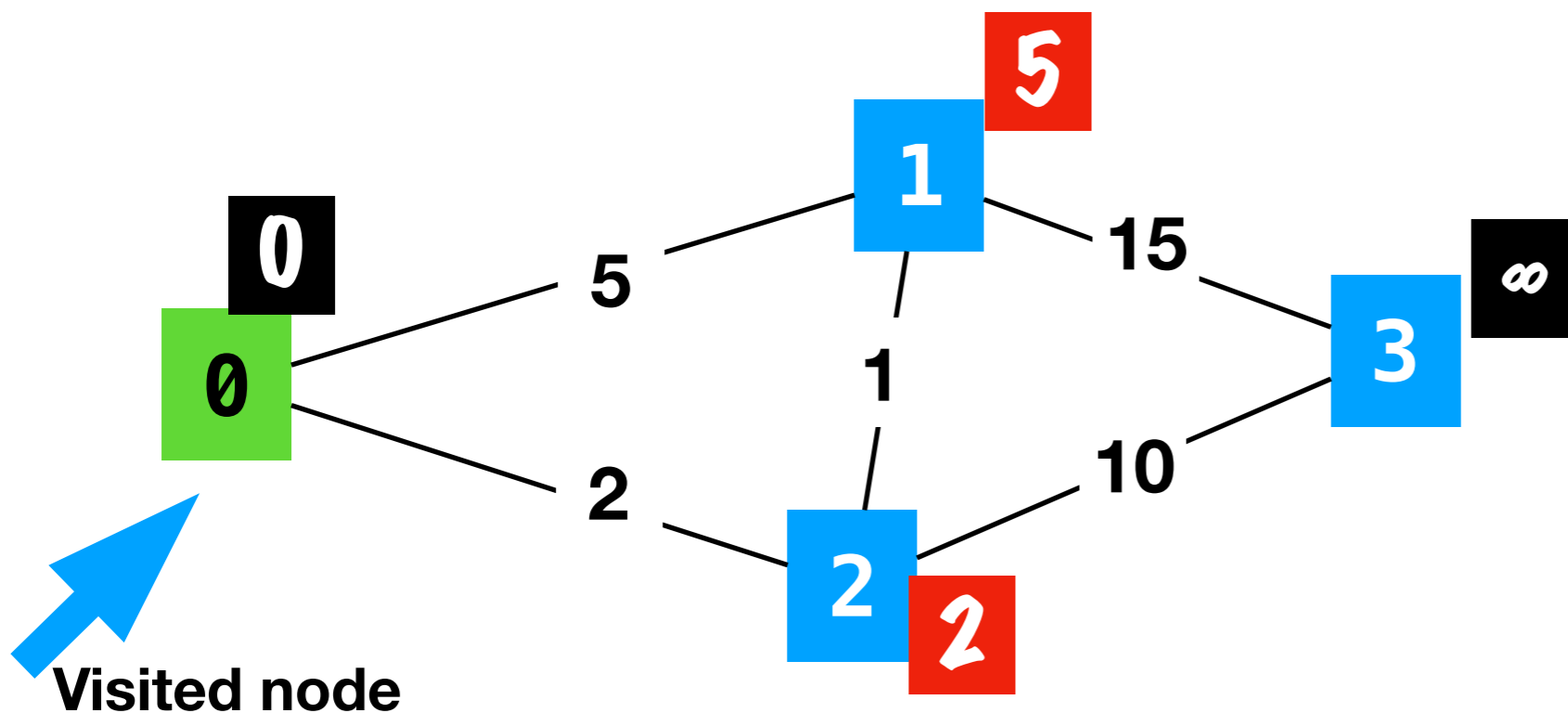


v	total cost	prev Node
0	0	--
1	$\infty$	--
2	$\infty$	--
3	$\infty$	--

# The Idea

- Update the values in the fringe if we update the best-known cost to a node.

fringe = { ( 2 , 2 ), ( 1 , 5 ), ( 3 , ∞ ) }

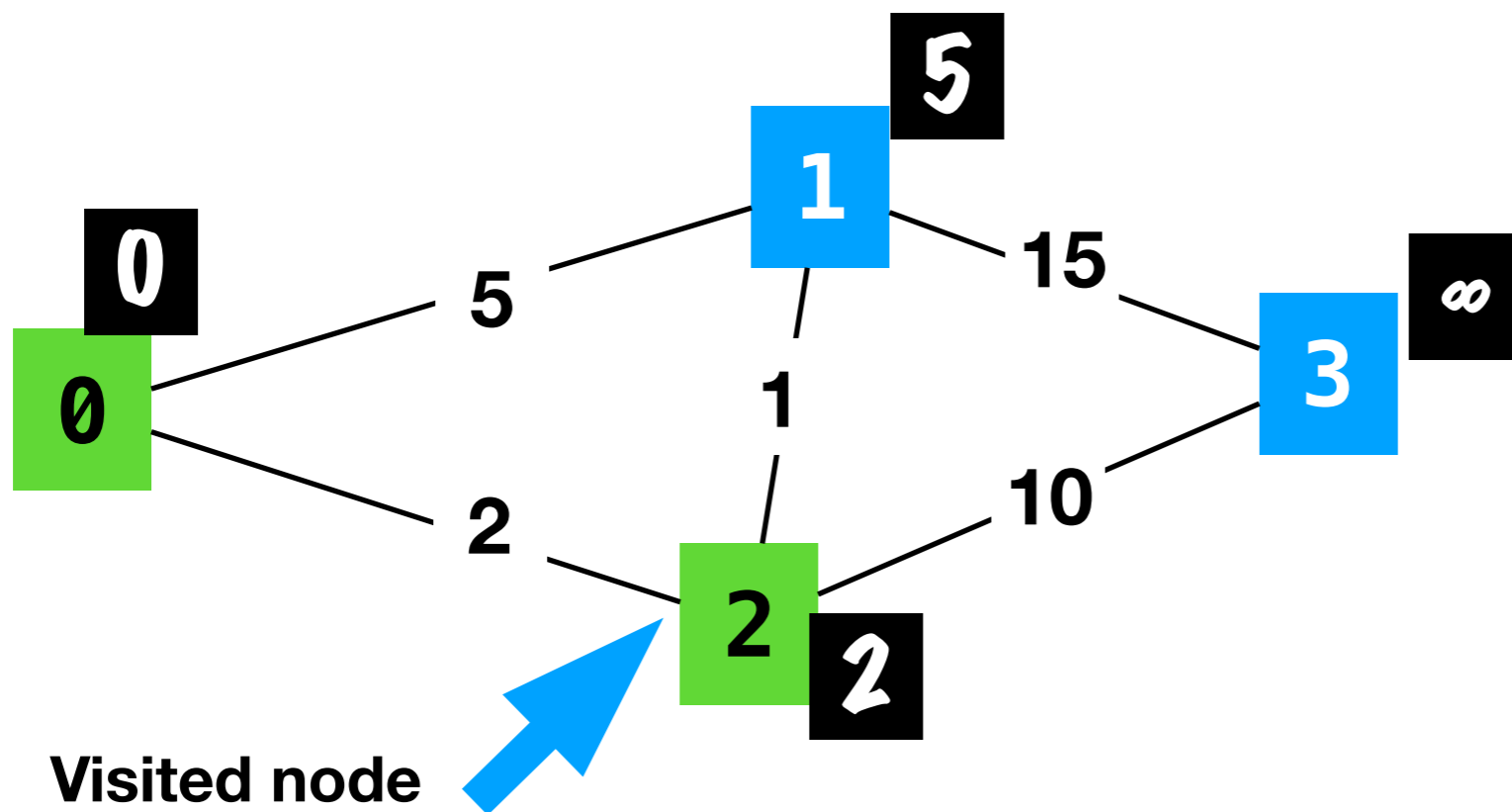


v	total cost	prev Node
0	0	--
1	5	0
2	2	0
3	∞	--

# The Idea

- Remove the minimum item from the fringe. This is our current node.

fringe = { ( **1**, **5** ), ( **3**,  $\infty$  ) }

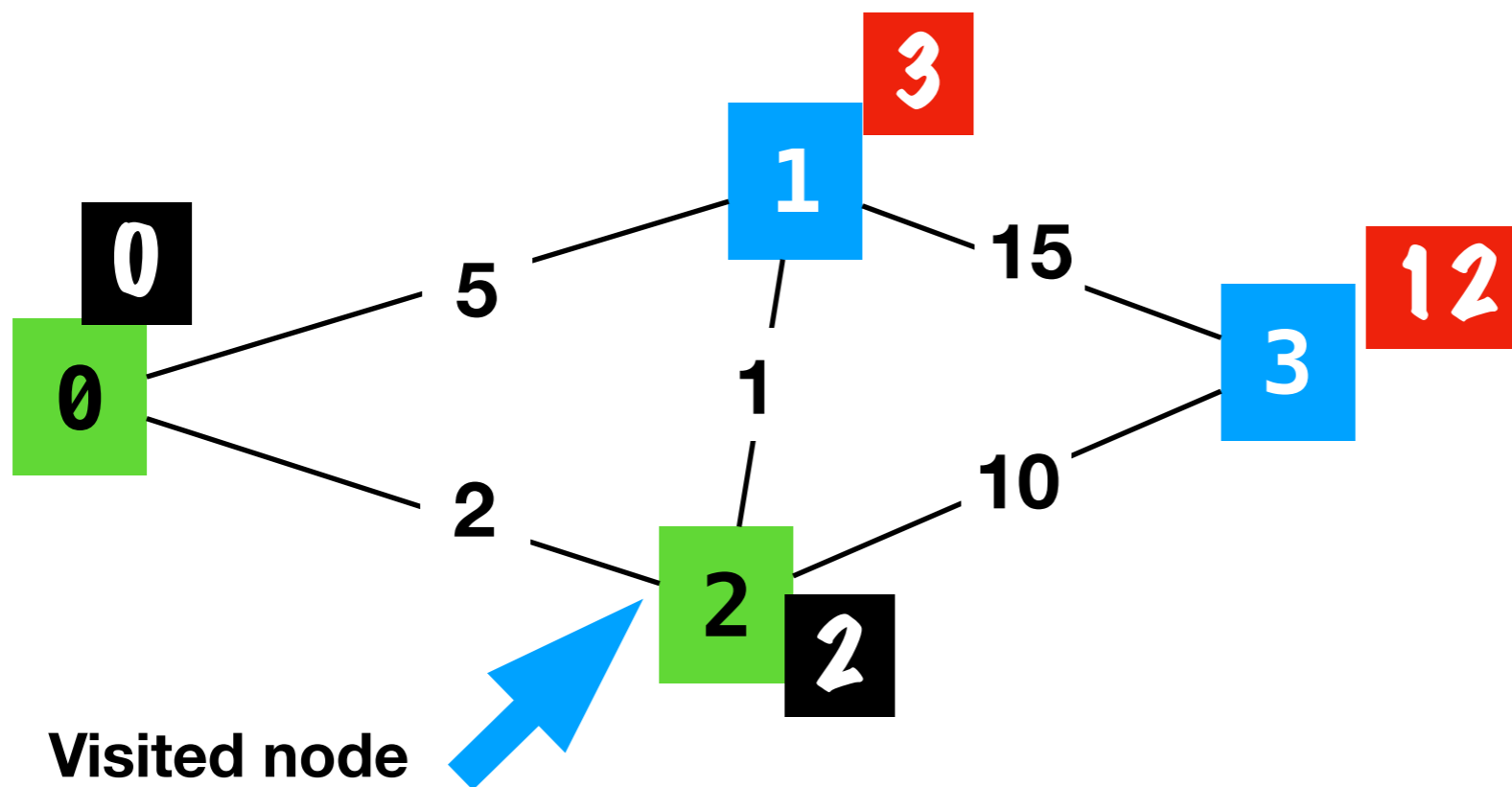


v	total cost	prev Node
0	0	--
1	5	0
2	2	0
3	$\infty$	--

# The Idea

- Update the values in the fringe if we update the best-known cost to a node.

fringe = { ( **1** , **3** ), ( **3** , **12** ) }



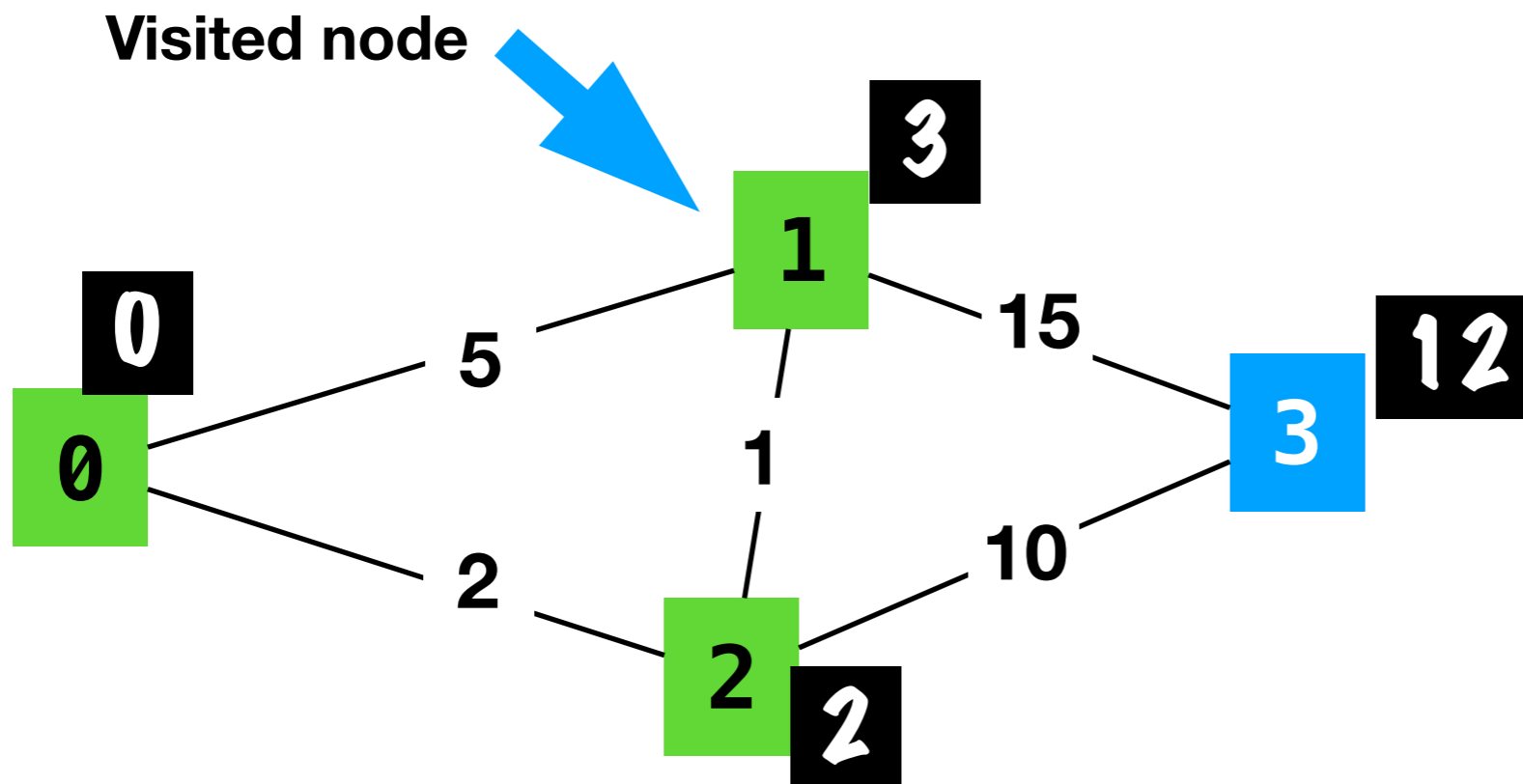
v	total cost	prev Node
0	0	--
1	3	2
2	2	0
3	12	2



# The Idea

- Remove the minimum item from the fringe. This is our current node.

fringe = { ( **3** , **12** ) }

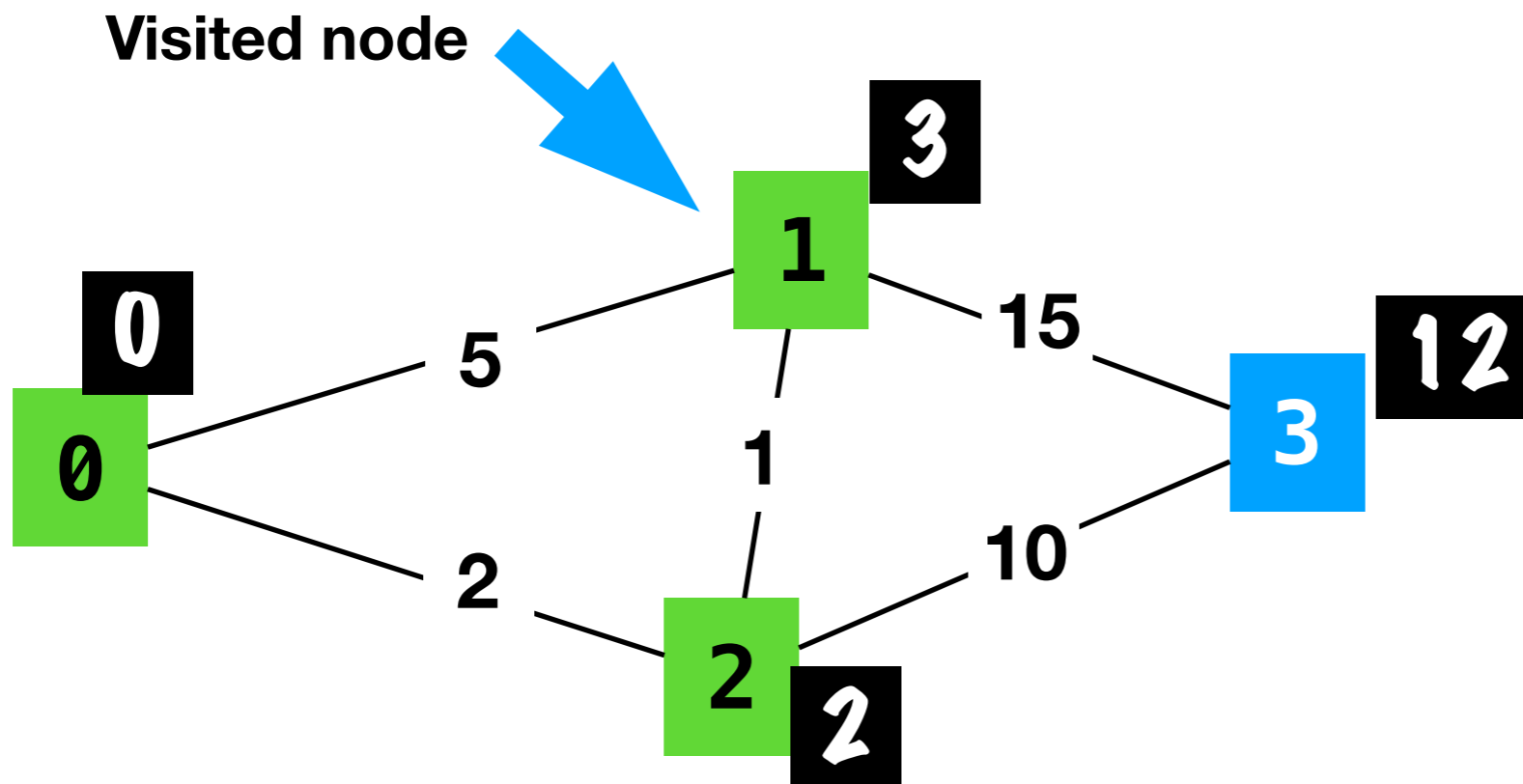


v	total cost	prev Node
0	0	--
1	3	2
2	2	0
3	12	2

# The Idea

- No update to neighbor since the path through the current node is not better than the neighbor's current cost.

fringe = { ( **3** , **12** ) }

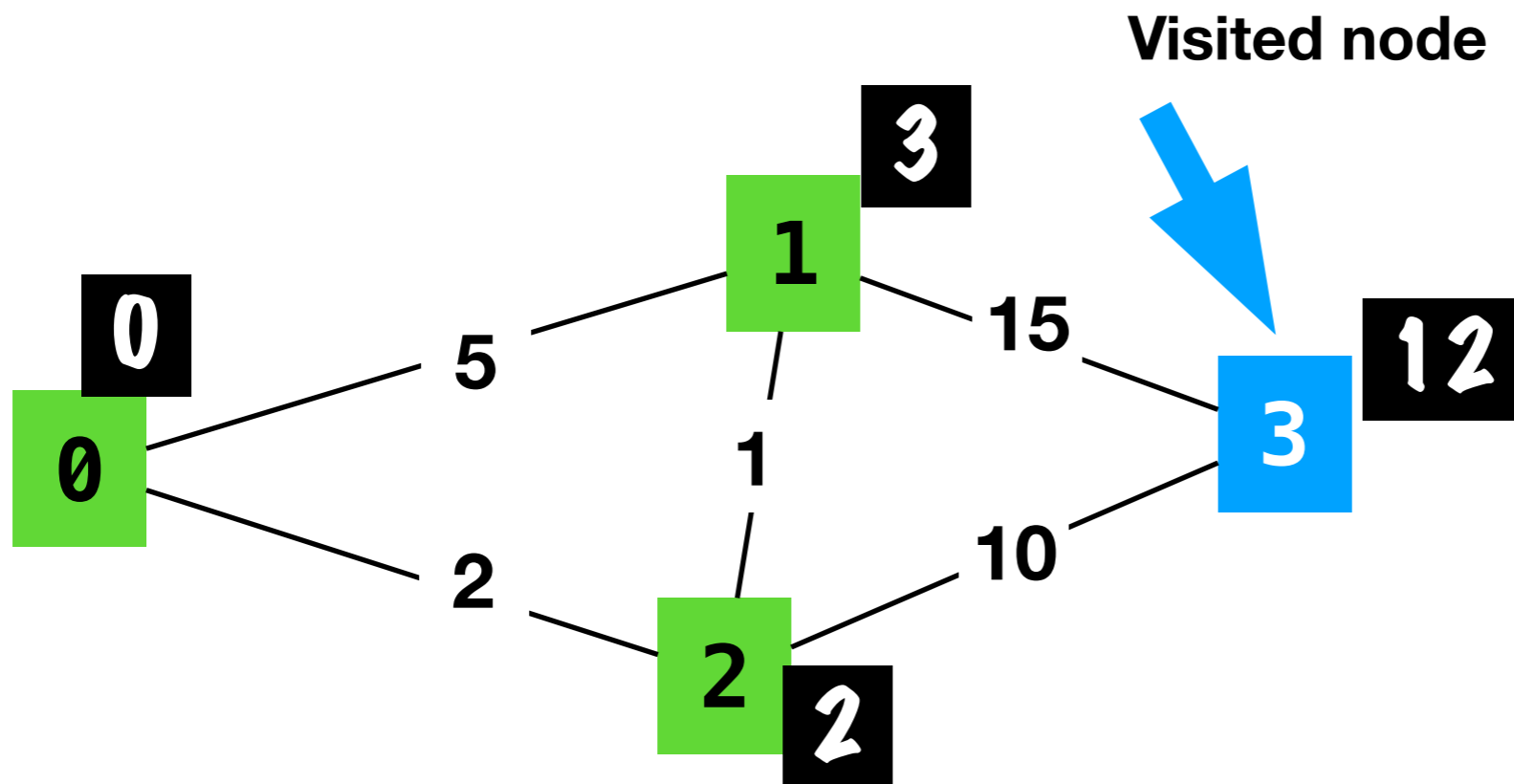


v	total cost	prev Node
0	0	--
1	3	2
2	2	0
3	12	2

# The Idea

- Remove the minimum item from the fringe. This is our current node.

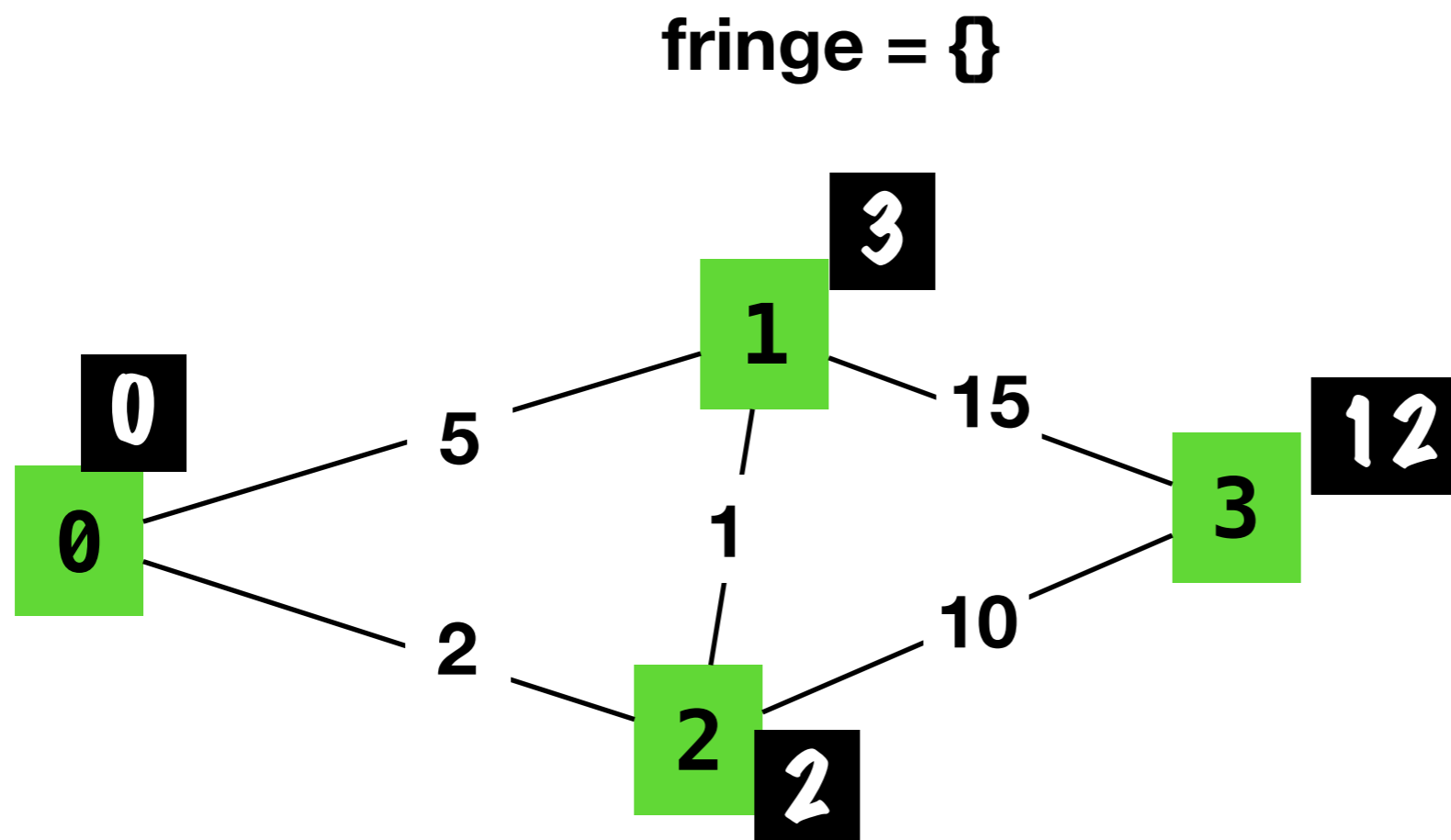
fringe = {}



v	total cost	prev Node
0	0	--
1	3	2
2	2	0
3	12	2

# The Idea

- Our fringe is empty, so we are done.



v	total cost	prev Node
0	0	--
1	3	2
2	2	0
3	12	2