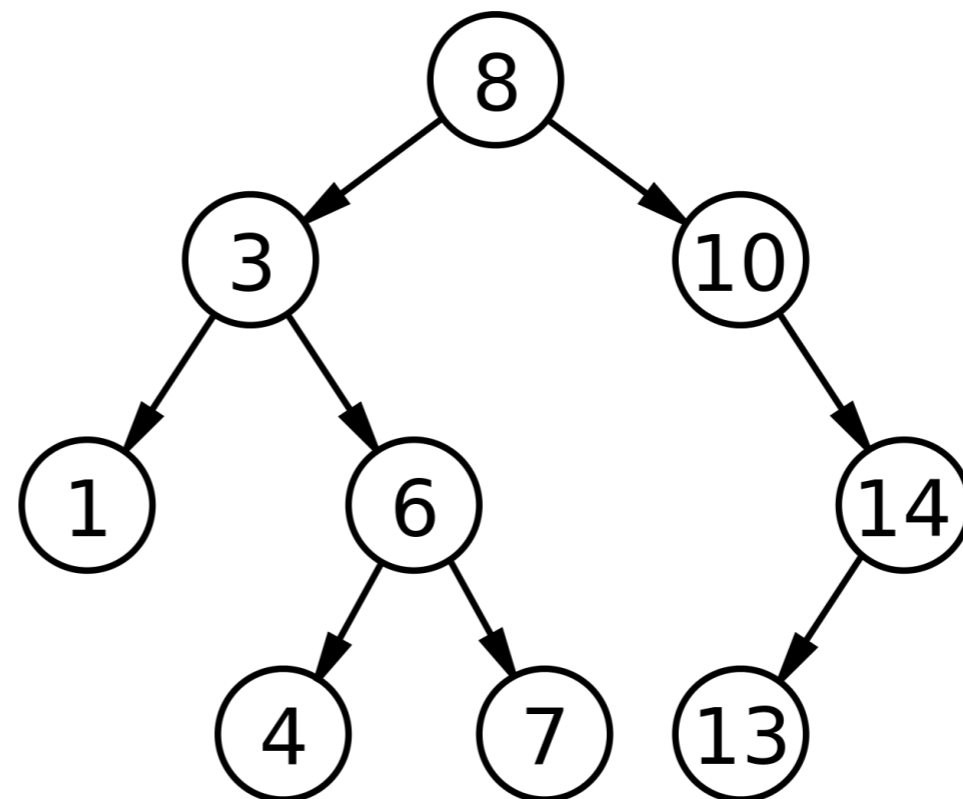


# CS 61BL Lab 11

Ryan Purpura

# Announcements

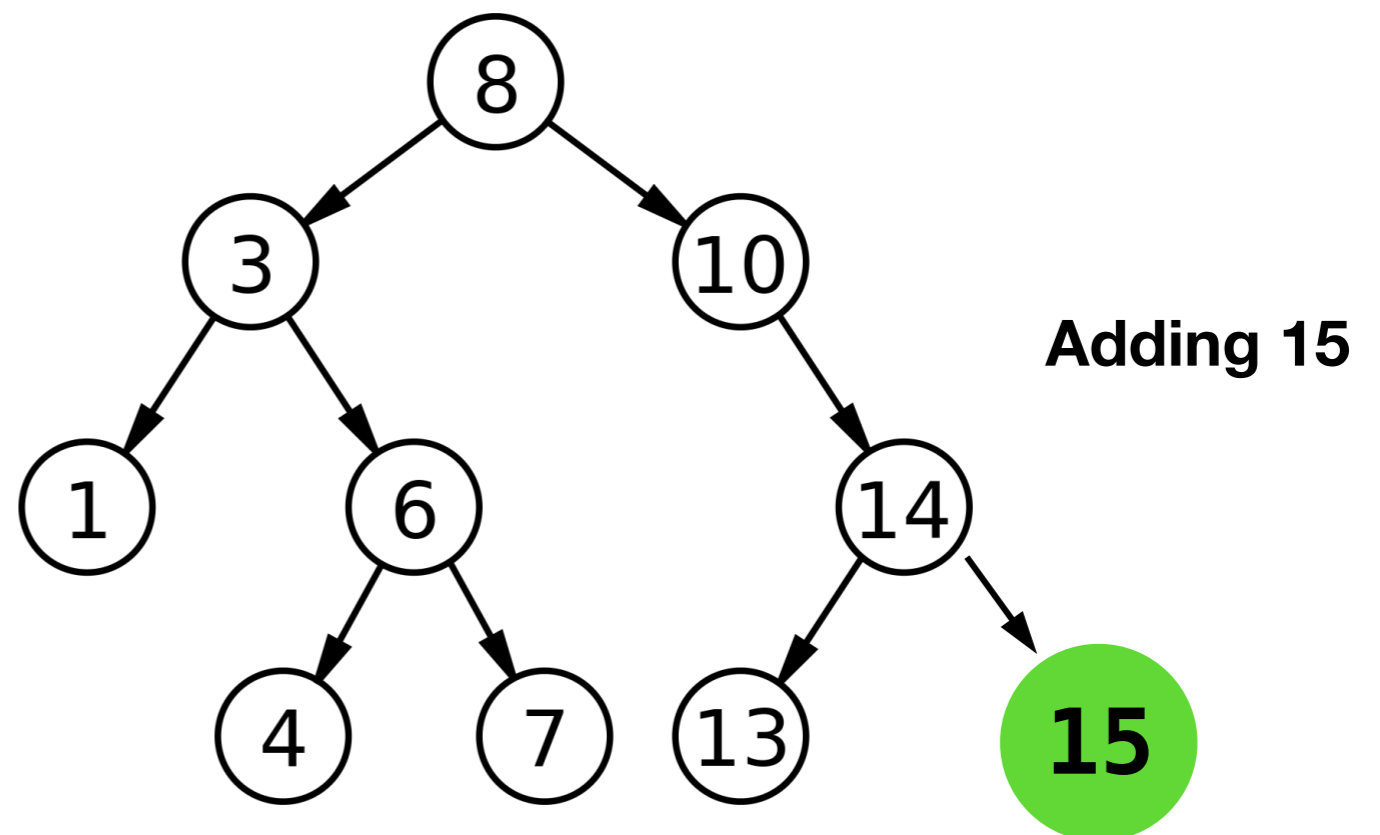
- Midterm 1 scores are out; regrade requests are open and will close Friday at the end of the day.
- Gitlet is due next Friday, July 26 at the end of the day.
- One clarification from Friday's lab: I said that when adding new nodes, you always add the new node to a leaf; this is not always true. For example:



**Adding 15**

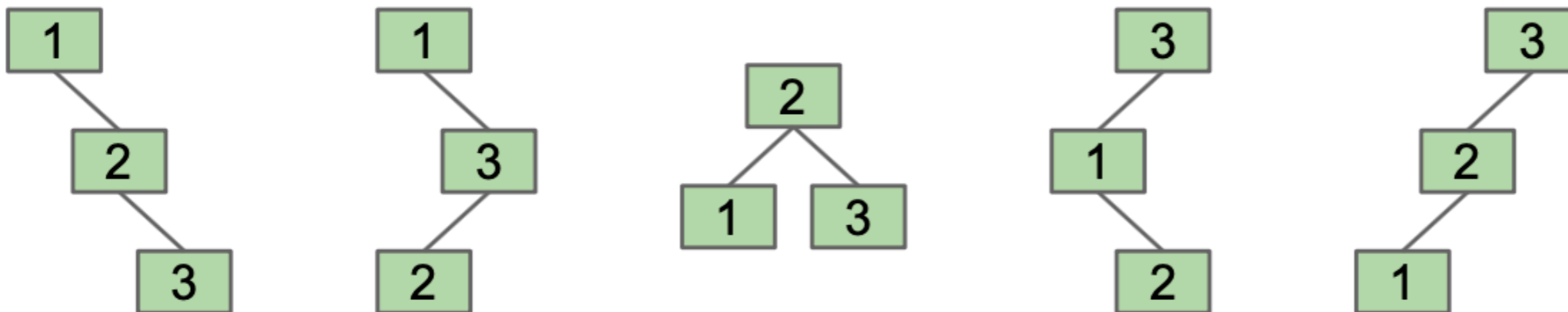
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# Tree Rotations

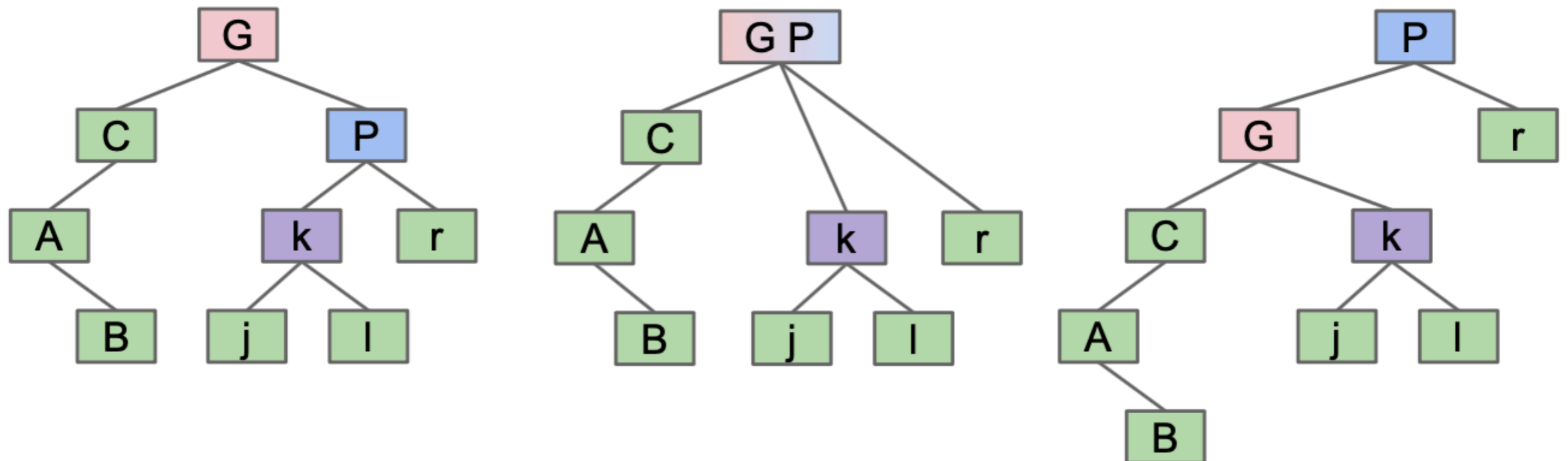
- The same elements can be organized into different binary search trees.



- We can use "tree rotations" to convert between equivalent binary search trees.

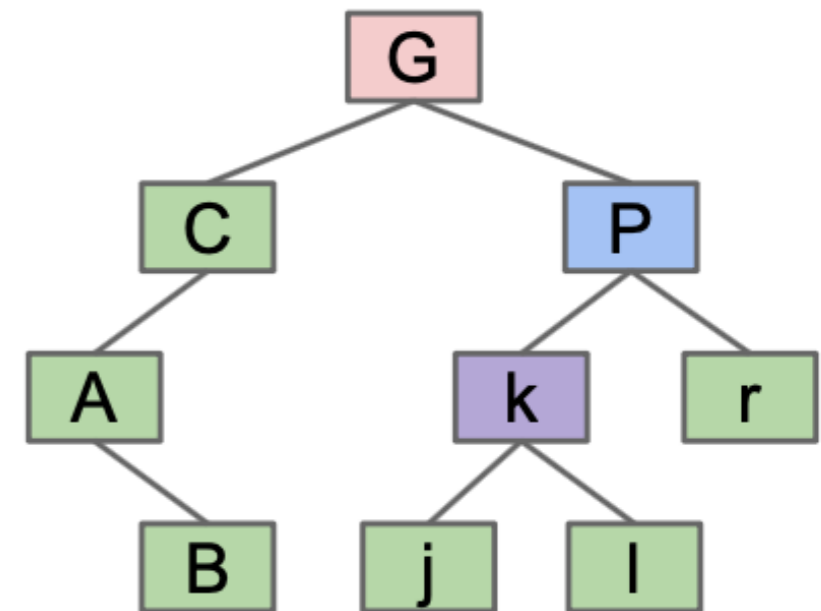
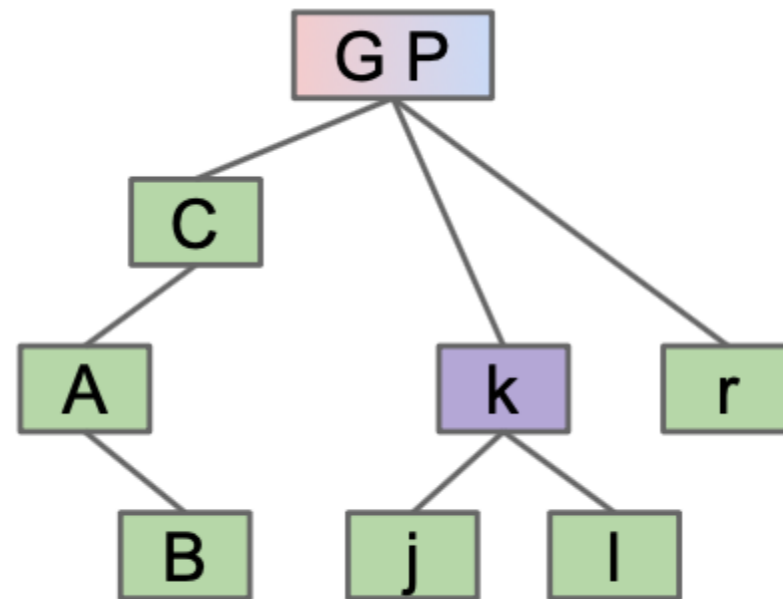
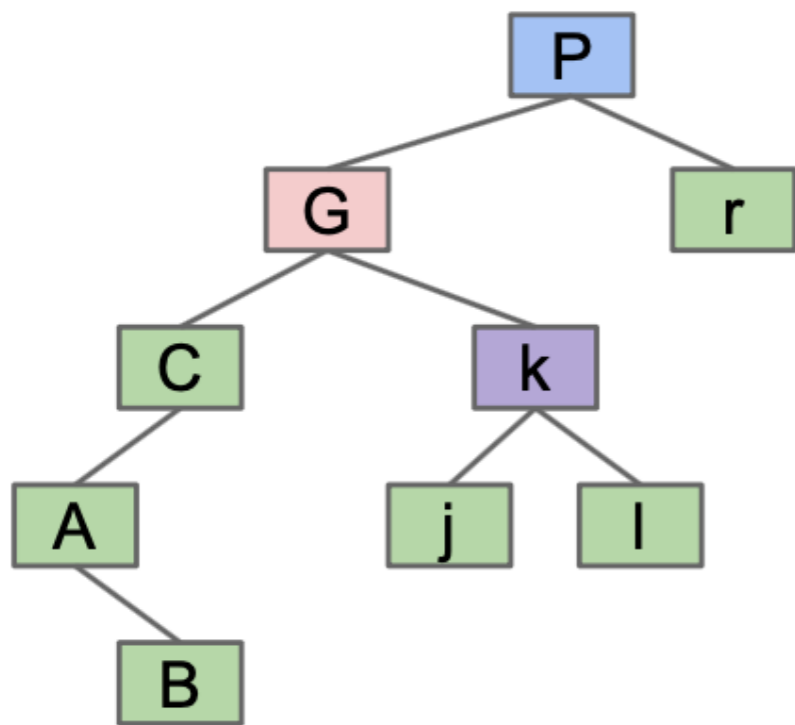
# Rotate Left

- rotateLeft(G):  
Let  $x$  be the right child of  $G$ . Make  $G$  the new left child of  $x$ .
- Can think of it as temporarily merging  $G$  and  $P$  and then sending  $G$  down and to the left. Notice what happens to  $k$ !



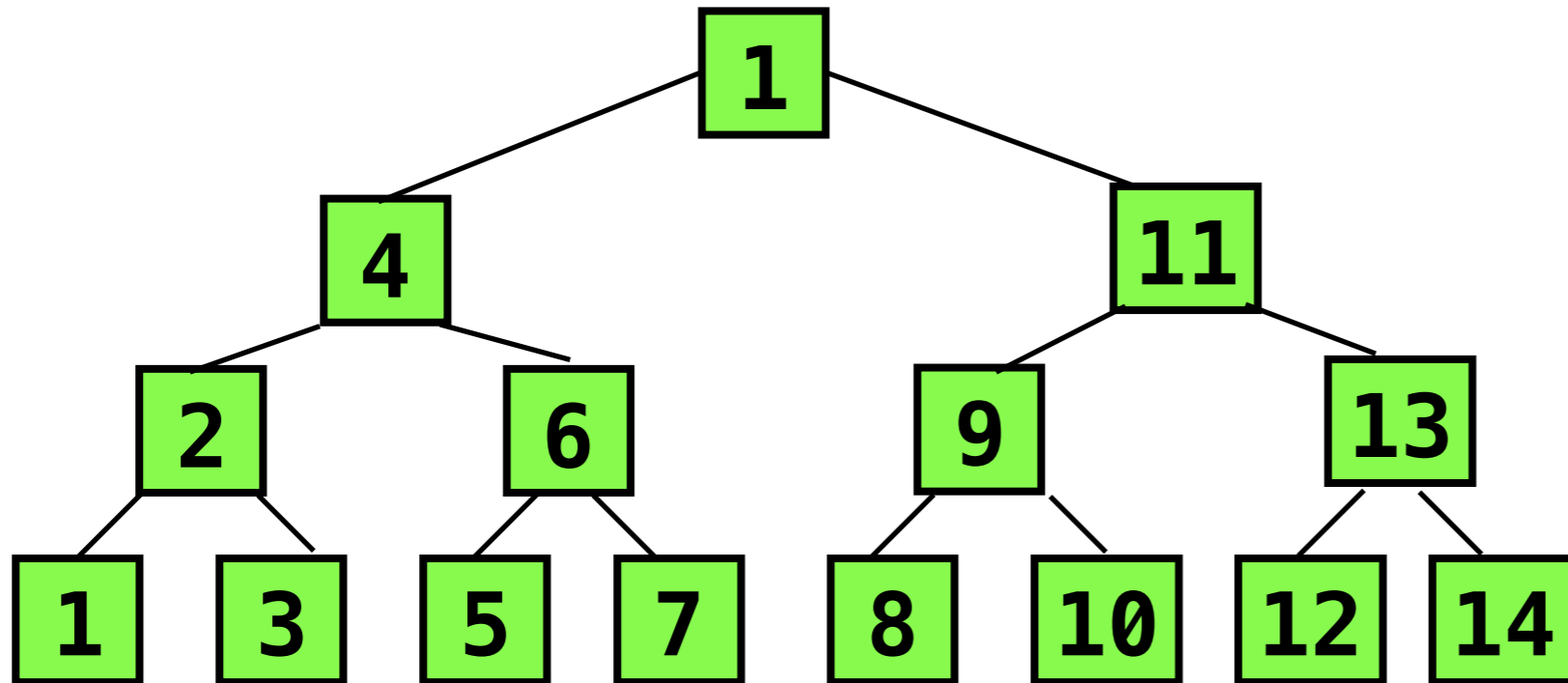
# Rotate Right

- rotateRight(P):  
Let  $x$  be the left child of  $P$ . Make  $P$  the new right child of  $x$ .
- Can think of it as temporarily merging  $G$  and  $P$  and then sending  $P$  down and to the right. Notice what happens to  $k$ !

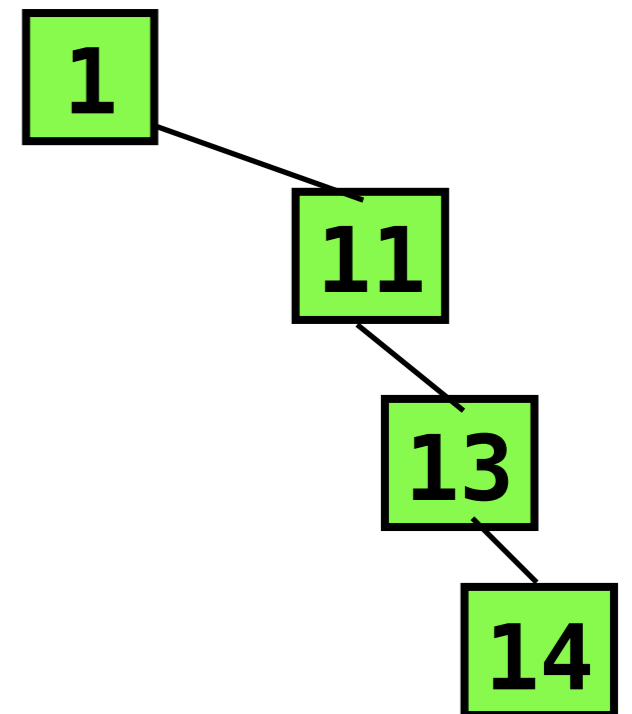


# Bushy vs. Spindly Trees

- Binary trees come in many shapes and sizes:



Nice bushy tree.



Very spindly tree:  
basically a linked  
list!

# Making Bushy Trees

- Idea 1: Smarter insertion order.
- Task: Make a balanced search tree with values 1, 2, 3, 4, 5, 6, 7
- Bad approach: add in order.



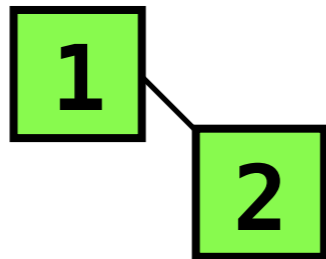
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**1**

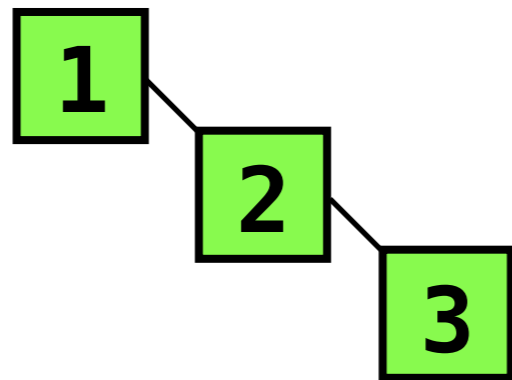
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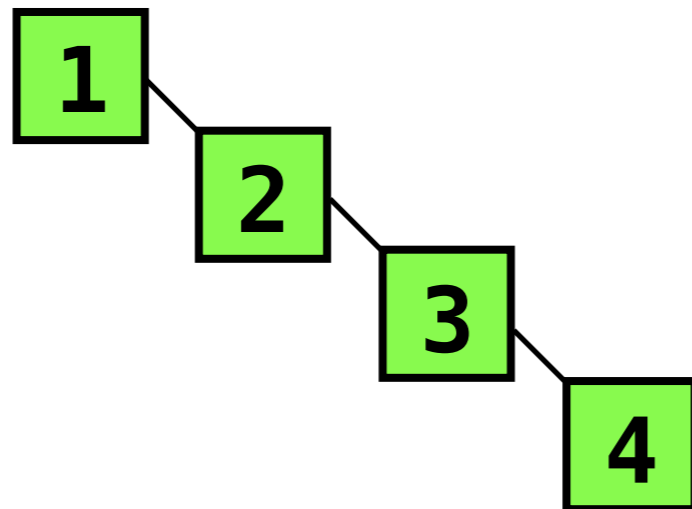
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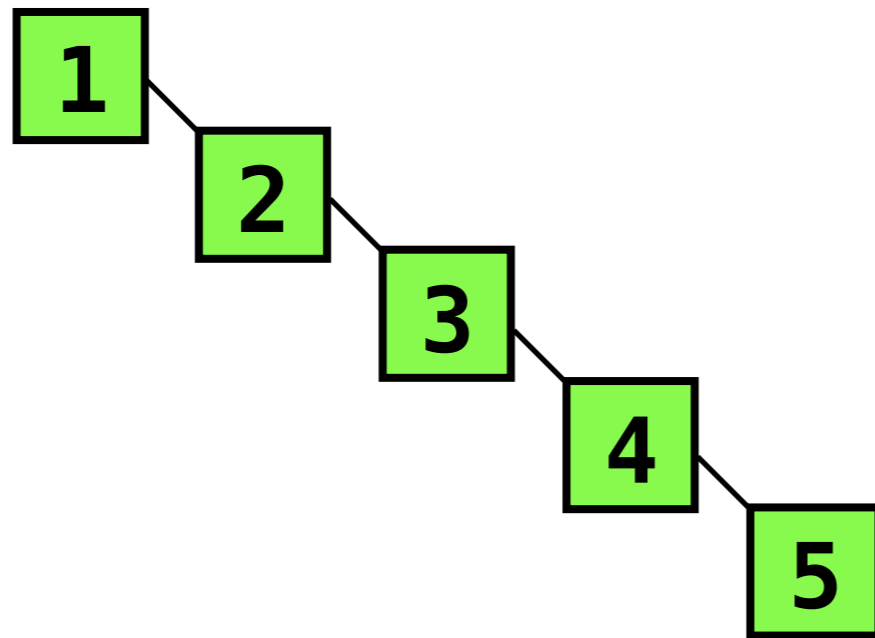
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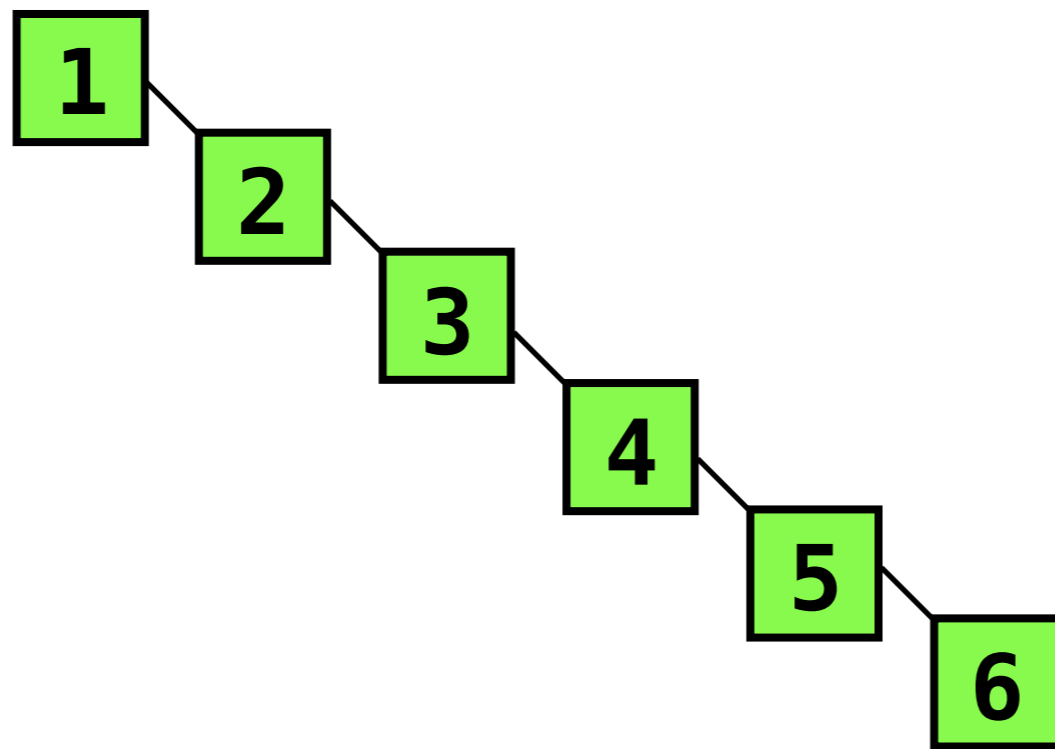
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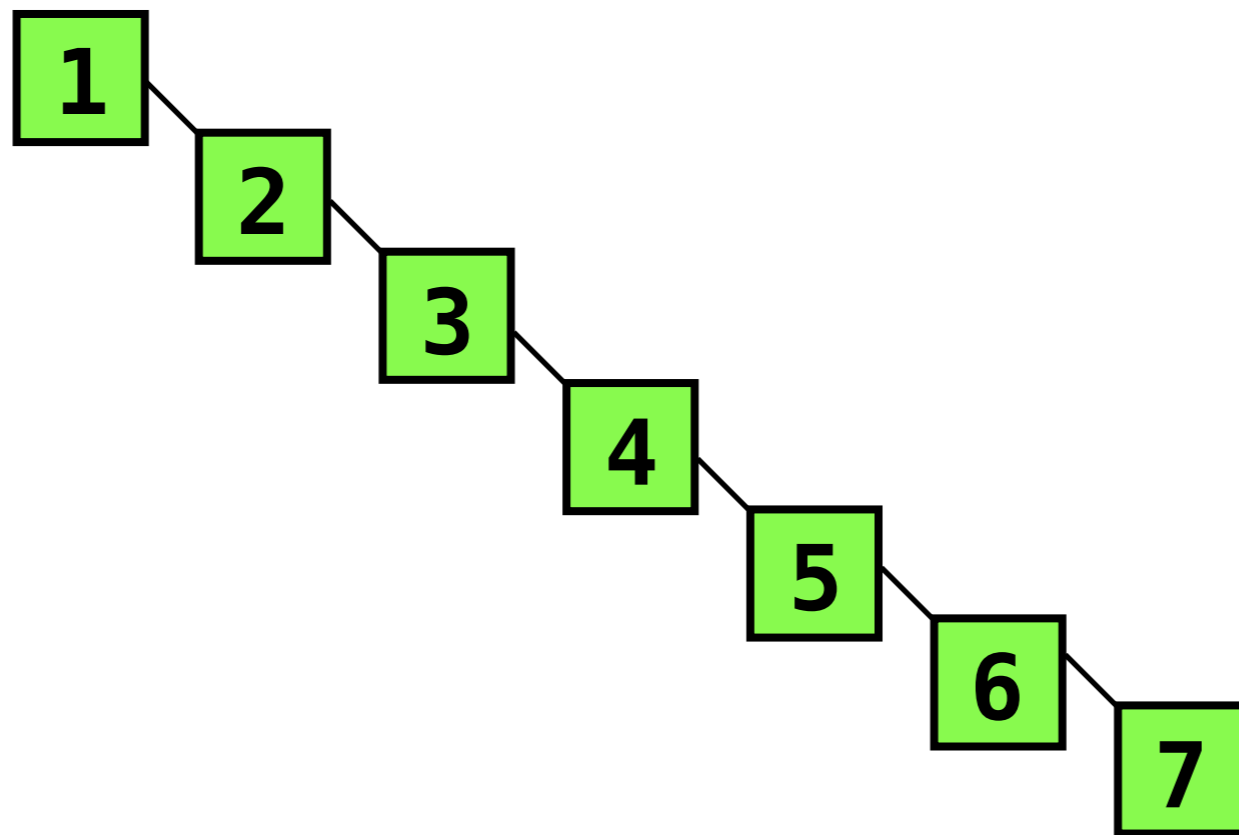
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# Better Approach

- Start in the middle, so that we make sure each node will have an equal number of elements for both children.
- So start with 4.



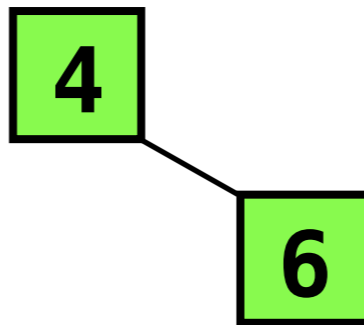
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4

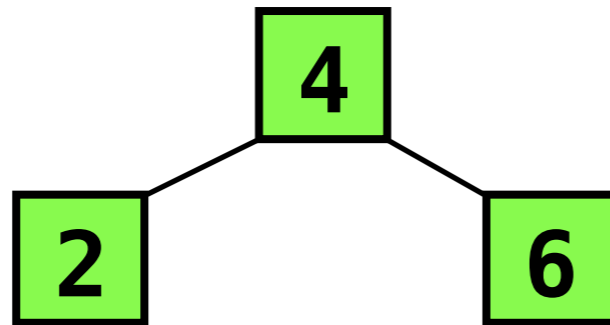
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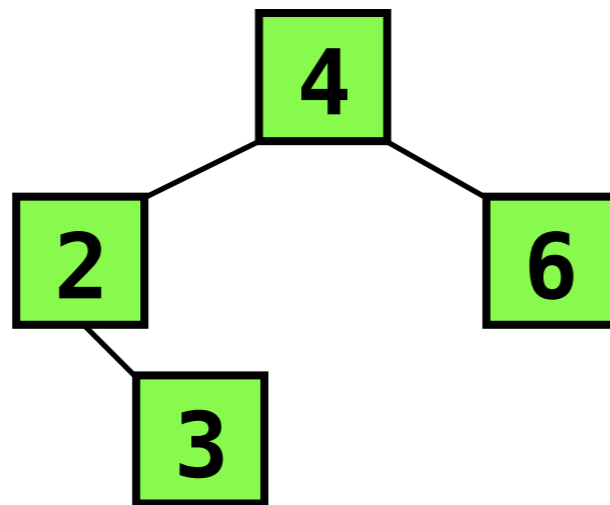
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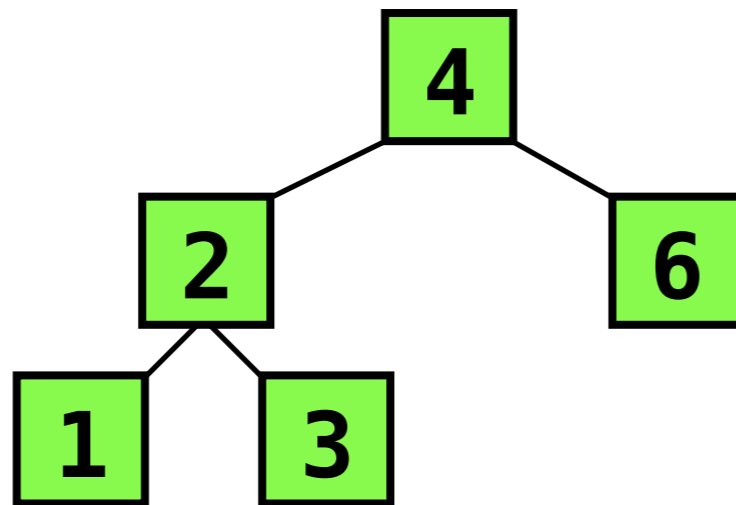
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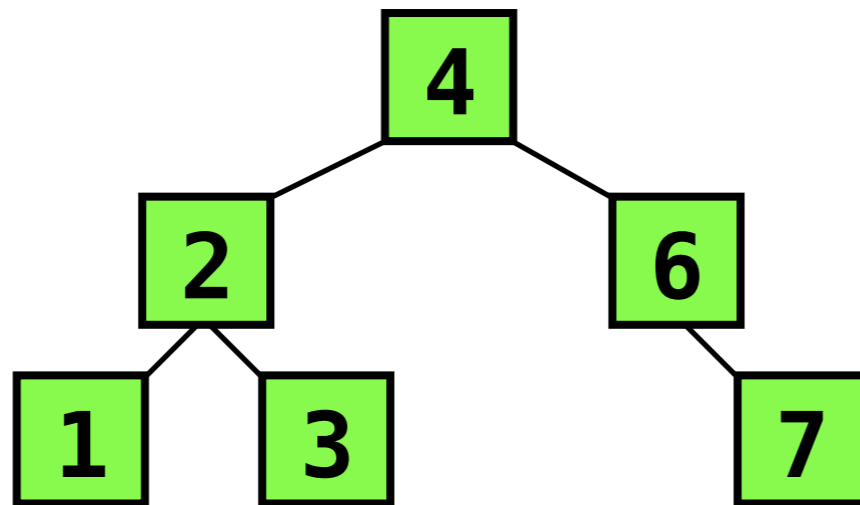
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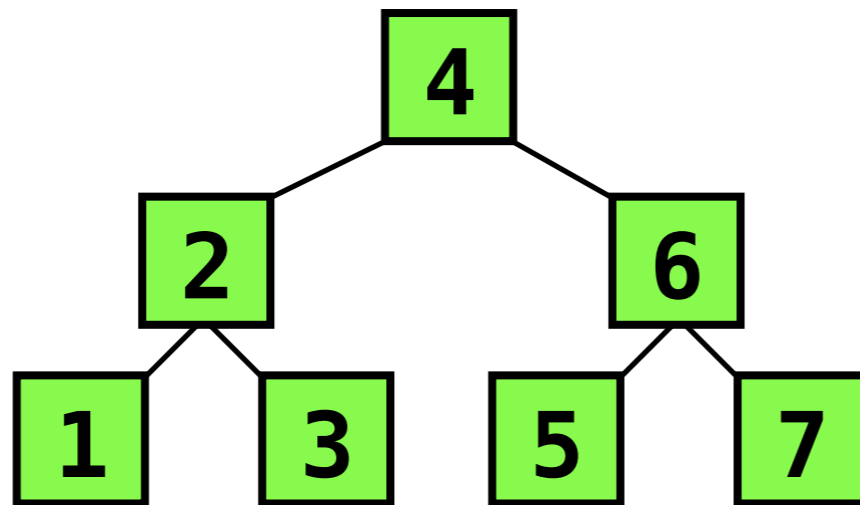
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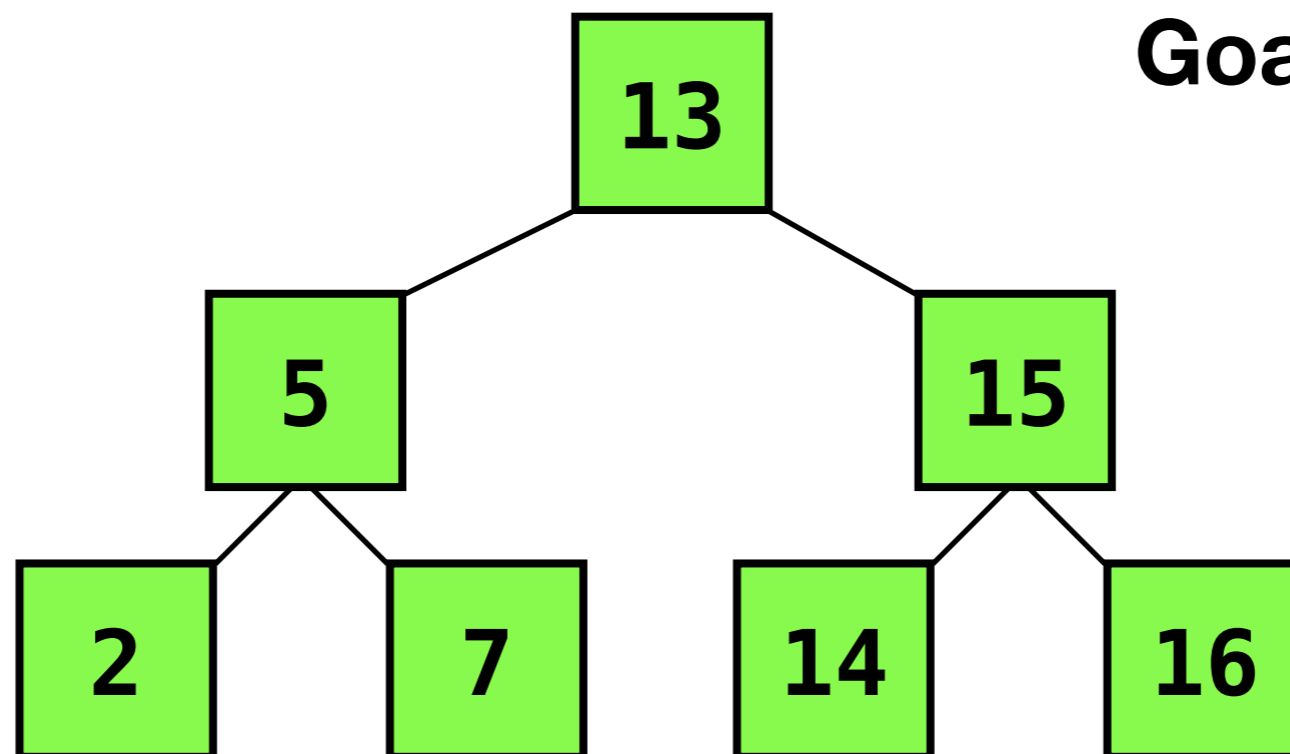
# What's the issue with this?

- We don't always have all the data in advance.



# B-Trees

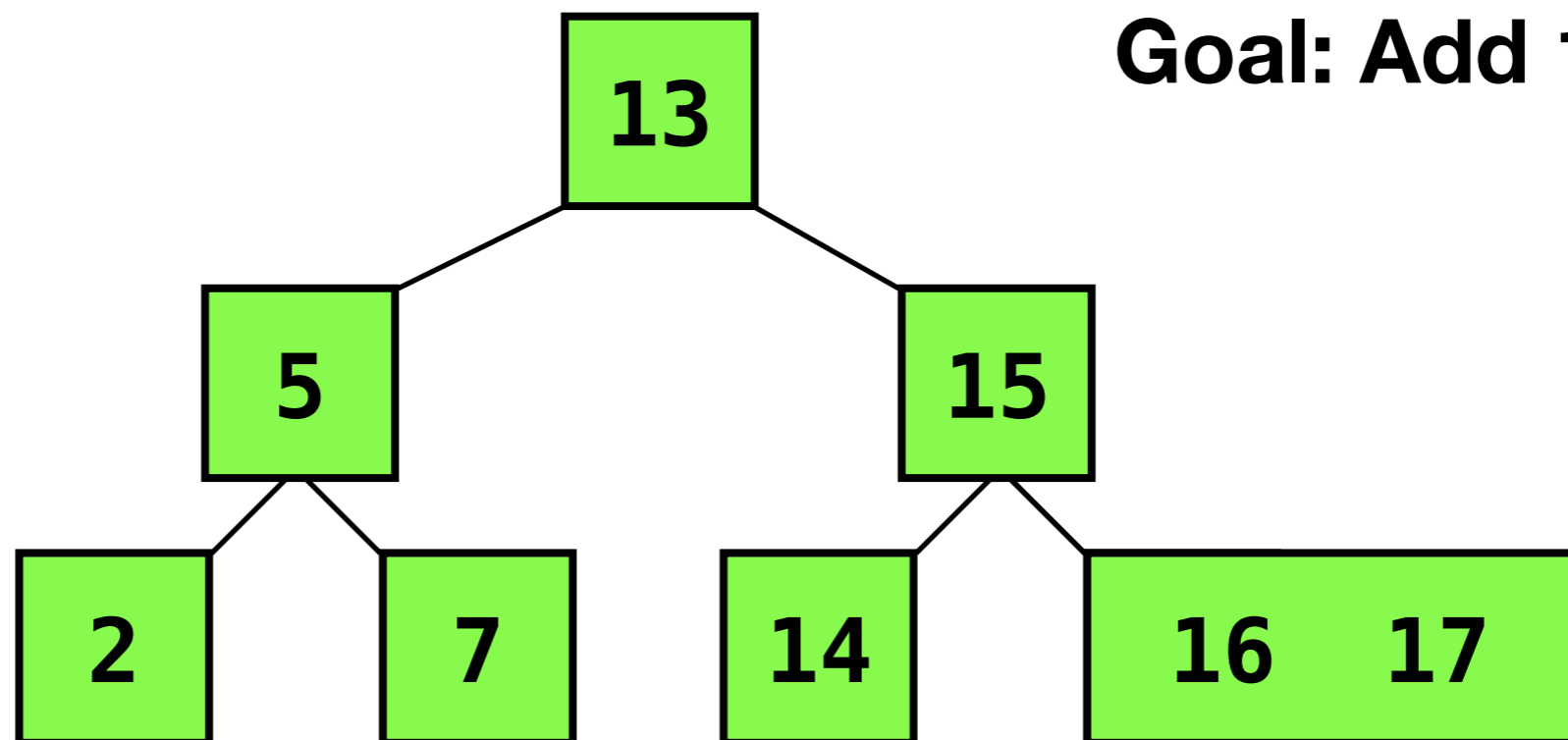
- The problem is adding leaves.
- So let's never add leaves!
- ...what do we do with insertions then?



**Goal: Add 17.**

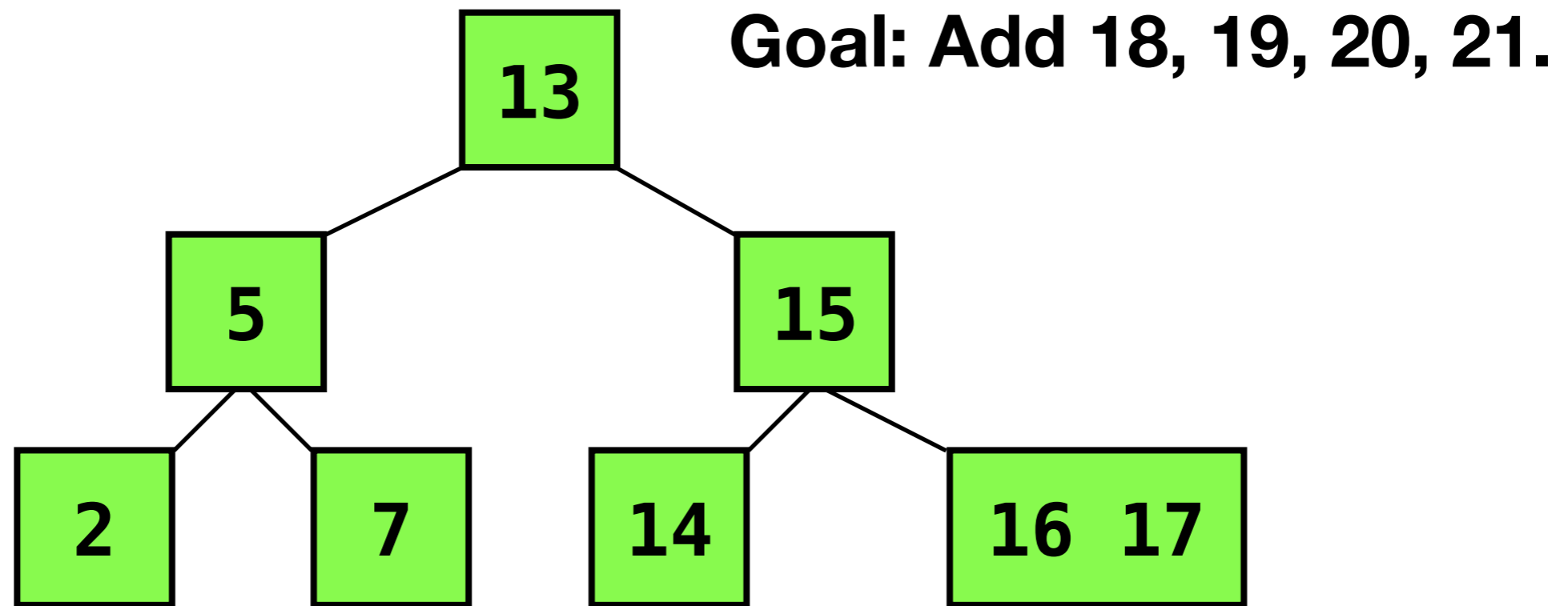
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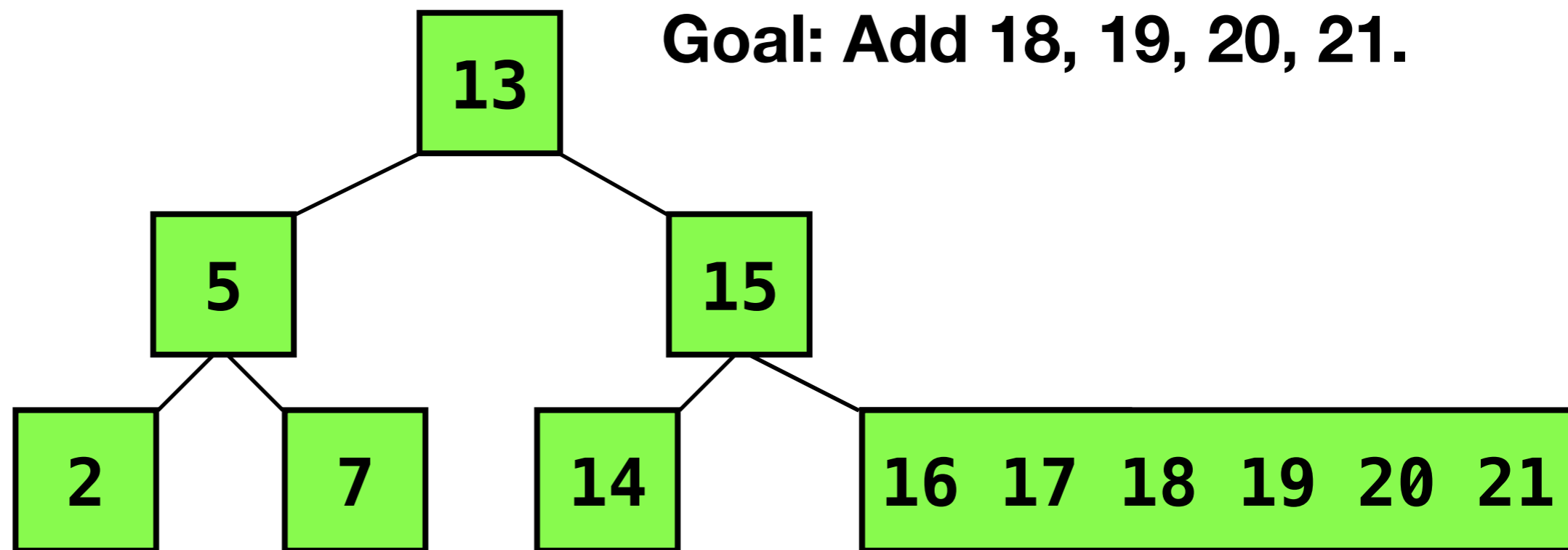


**Goal: Add 17.**

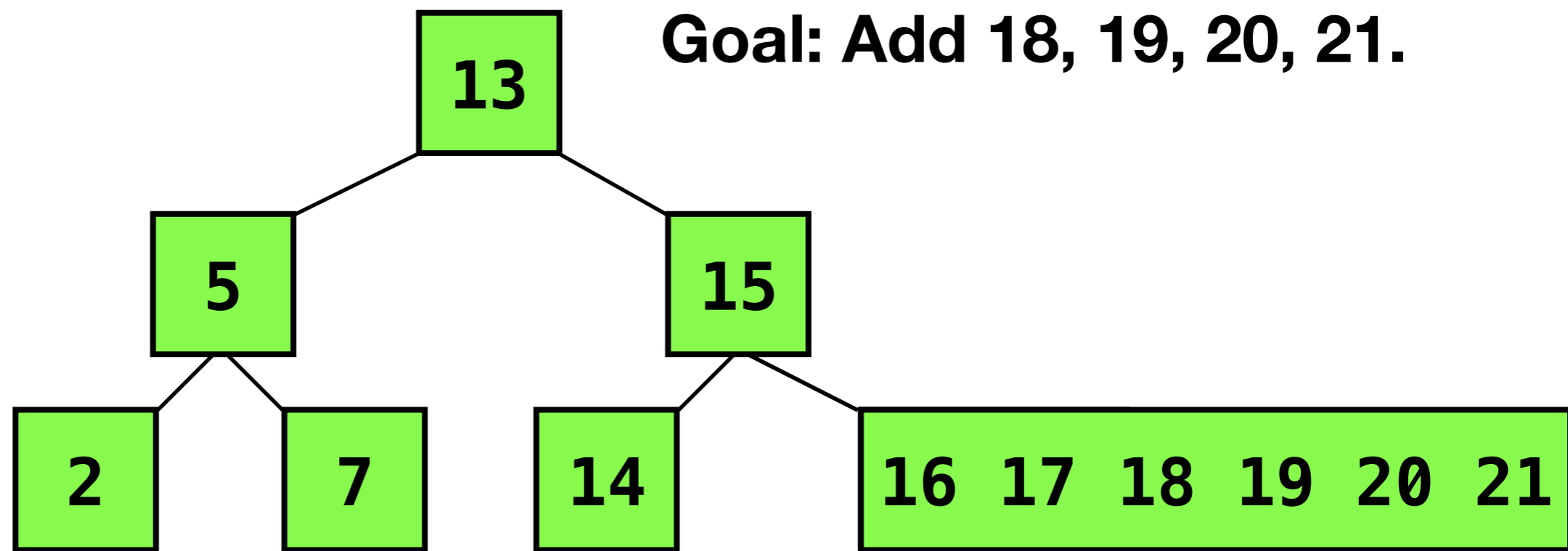
# Let's keep going!



# Let's keep going!



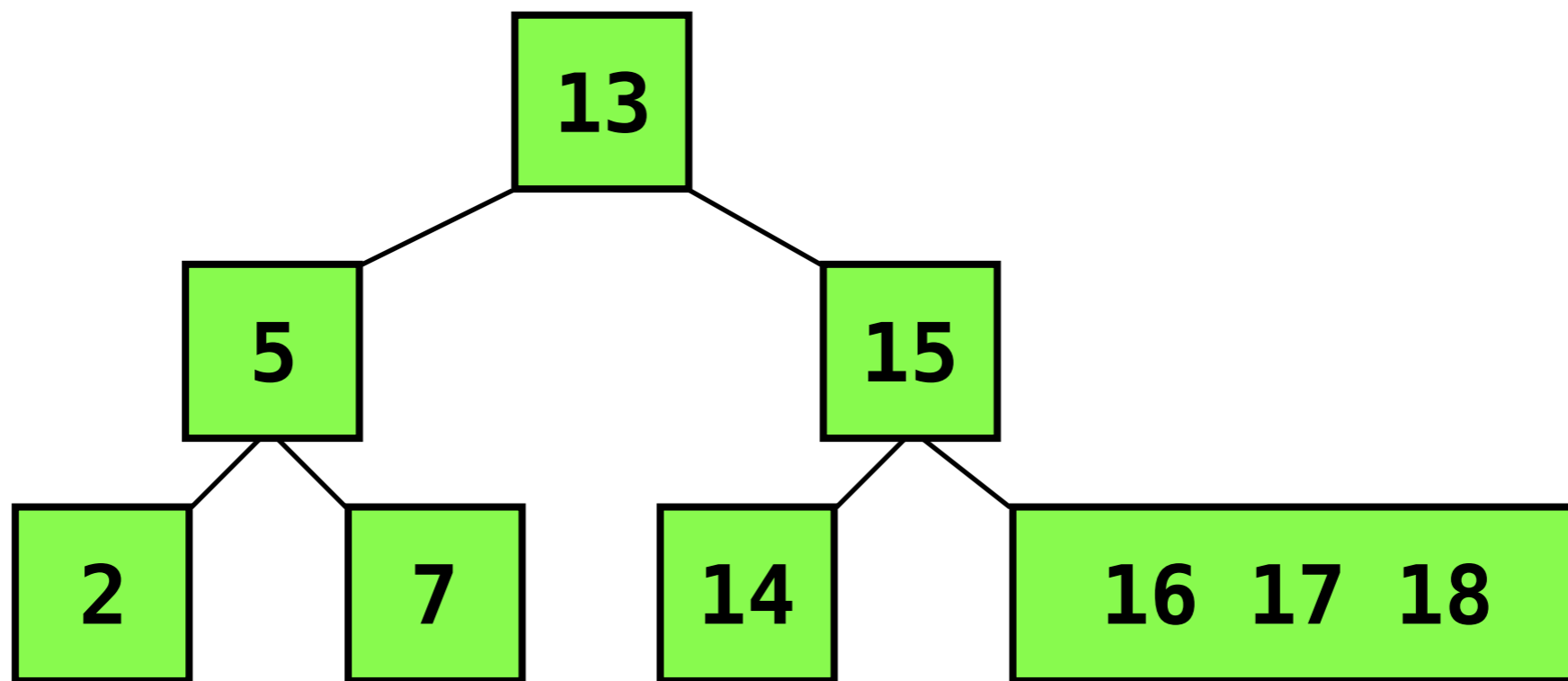
# Let's keep going!



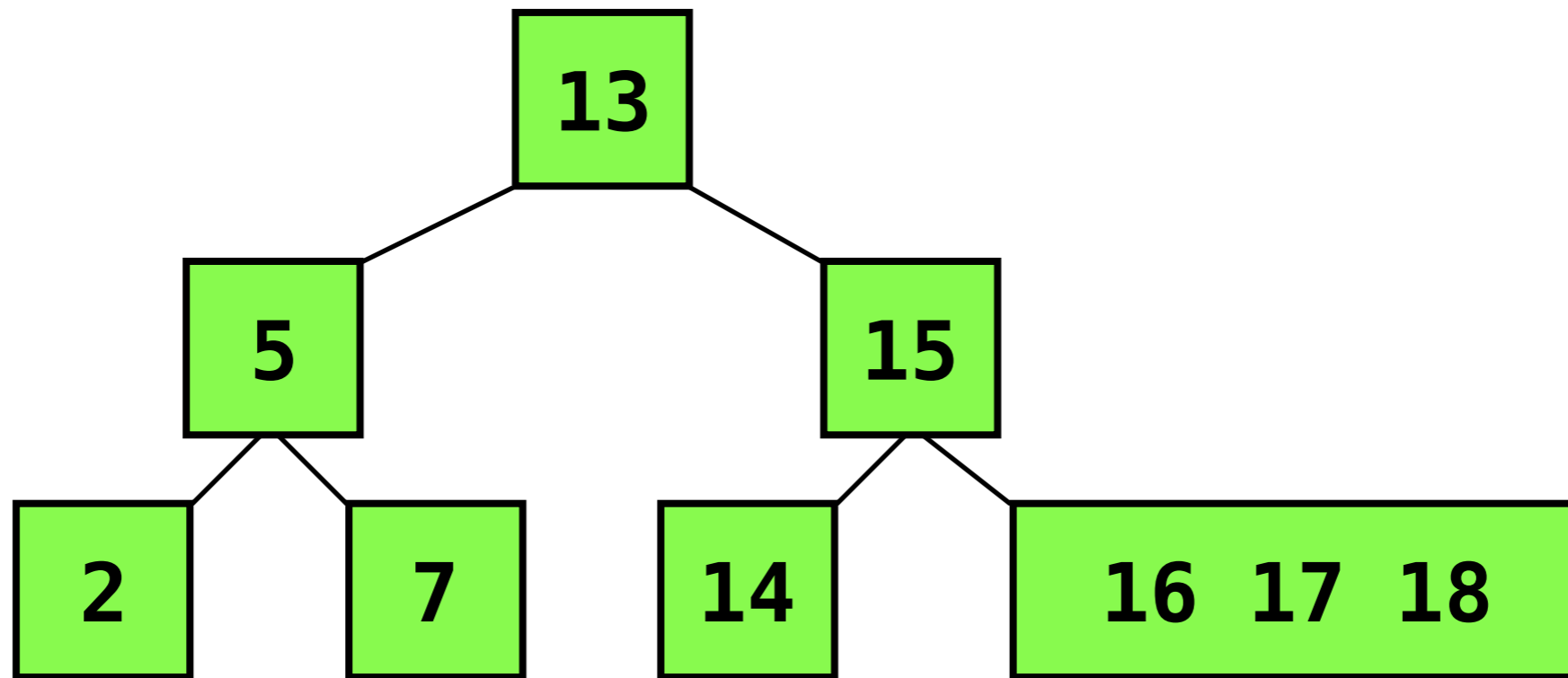
**This is getting unwieldy...**

# Item Limit

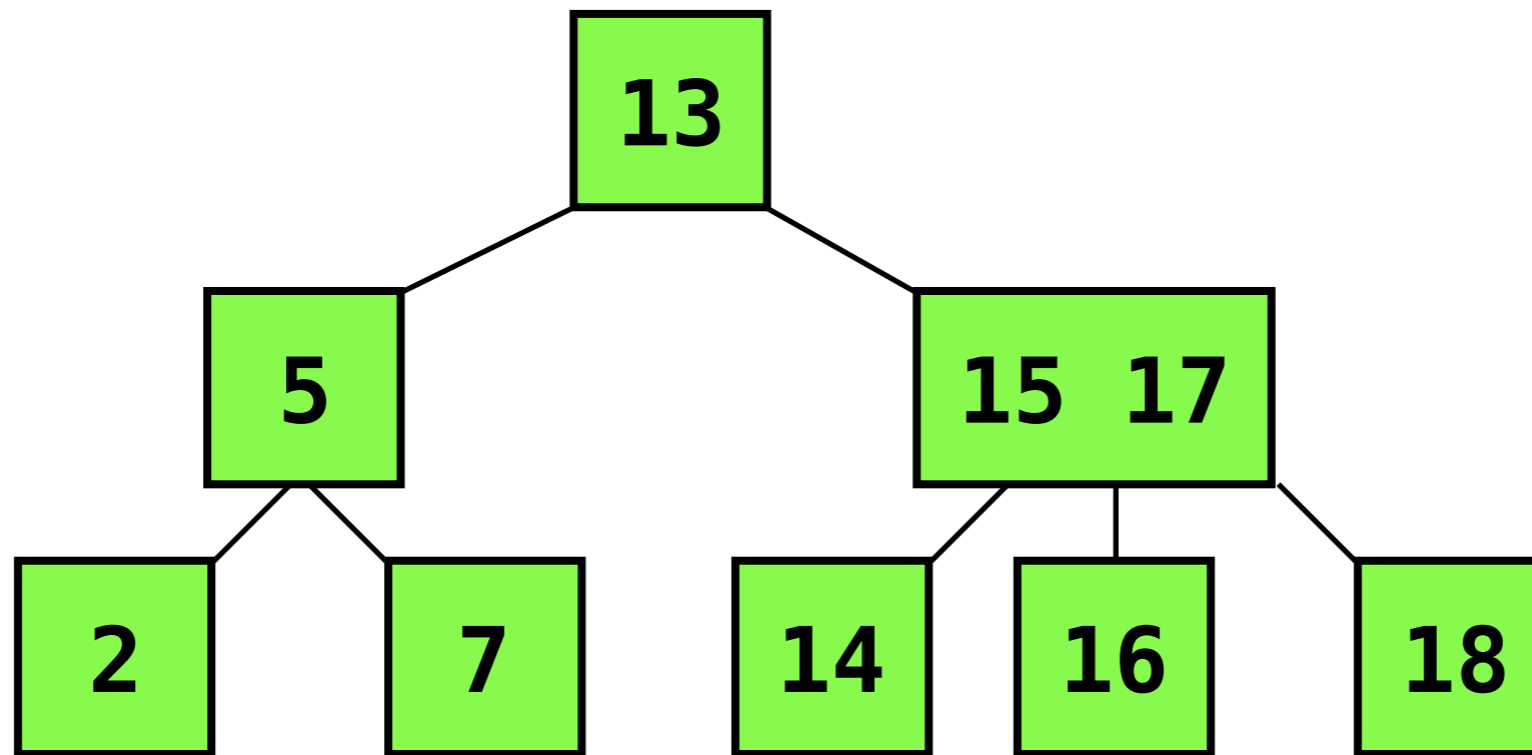
- Idea: add an item limit for each node.
- If the leaf we want to add to is at its item limit, give the middle item to its parent and split, then add.
- Example: Given limit 3, what happens when we insert 19?



# Adding 19 with Limit 3

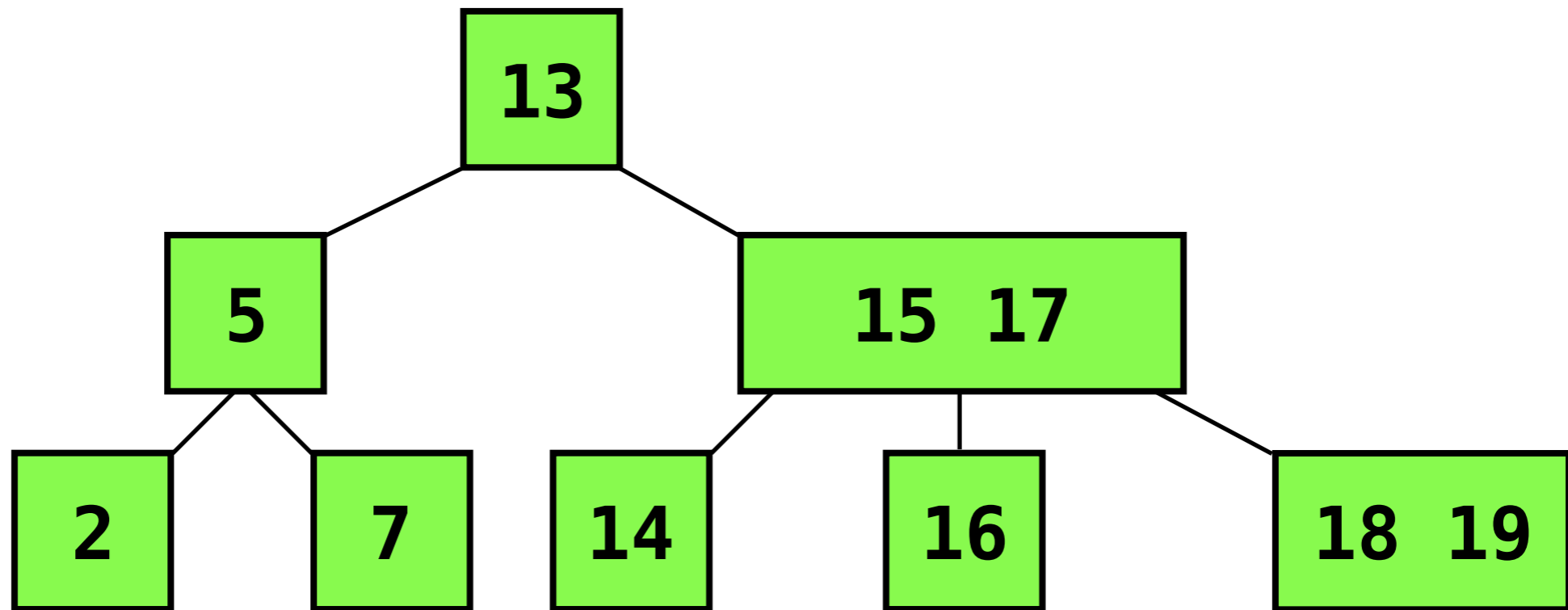


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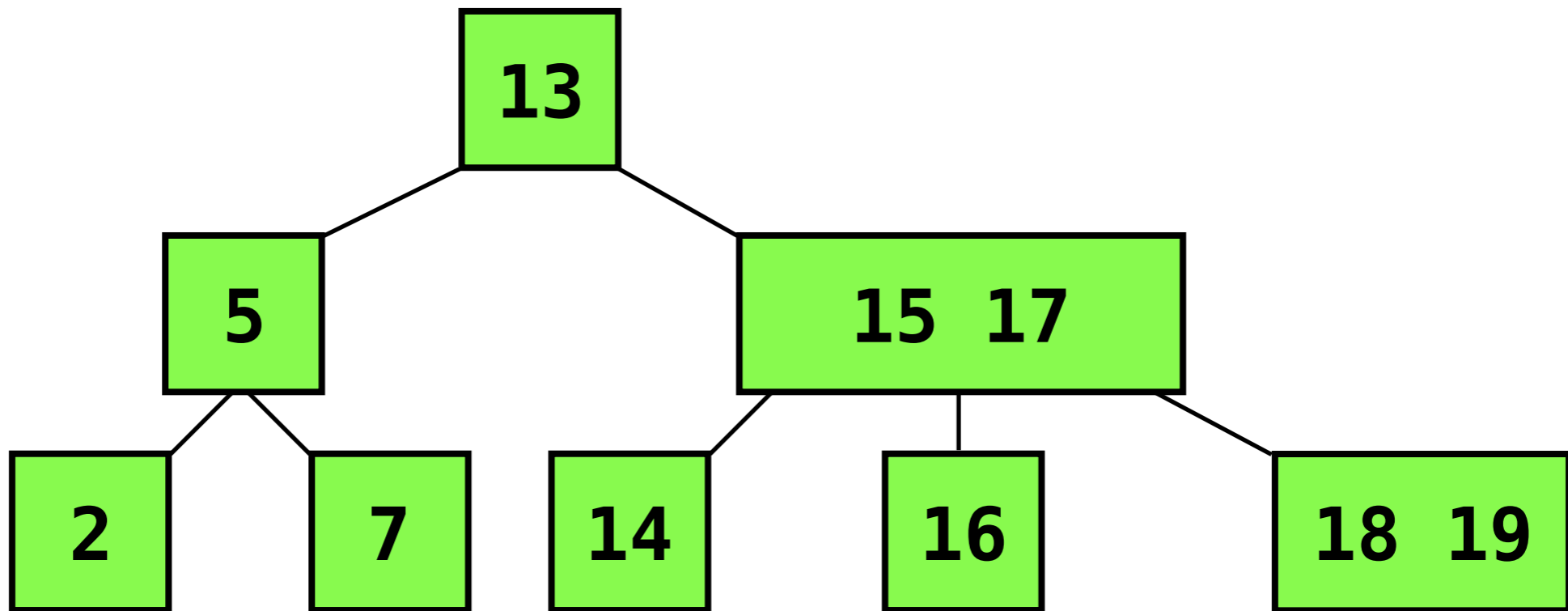


# Adding 19 with Limit 3

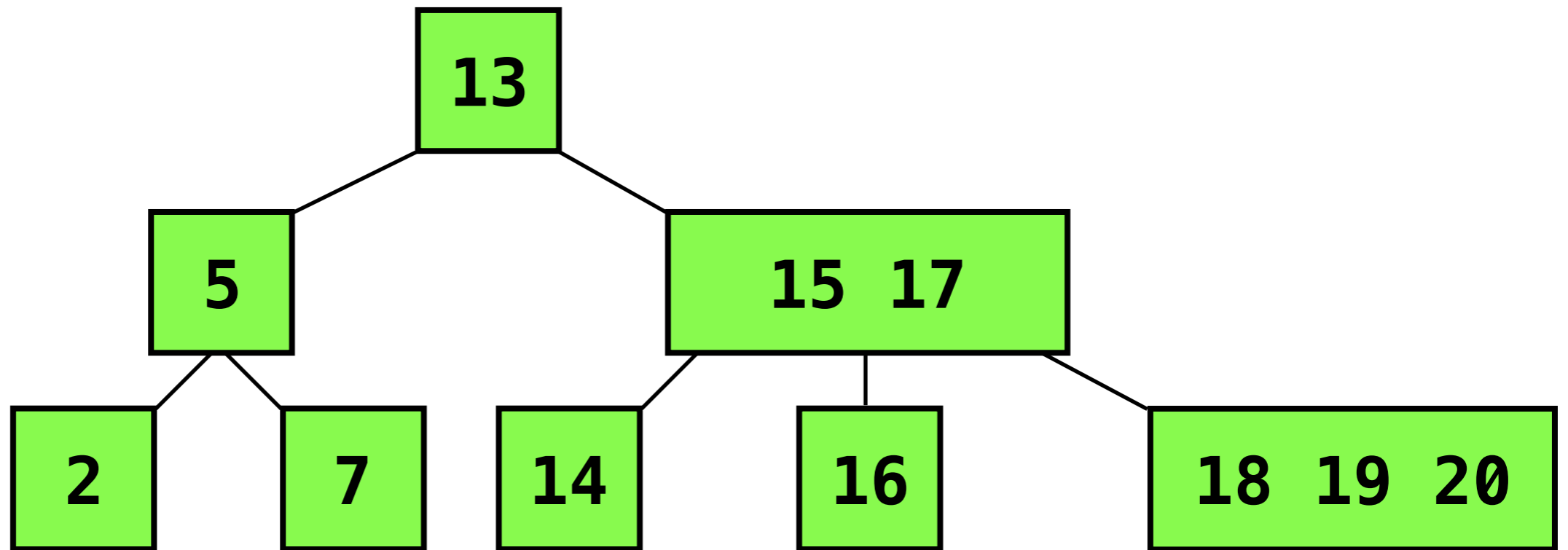


Now we can easily fit in 19.

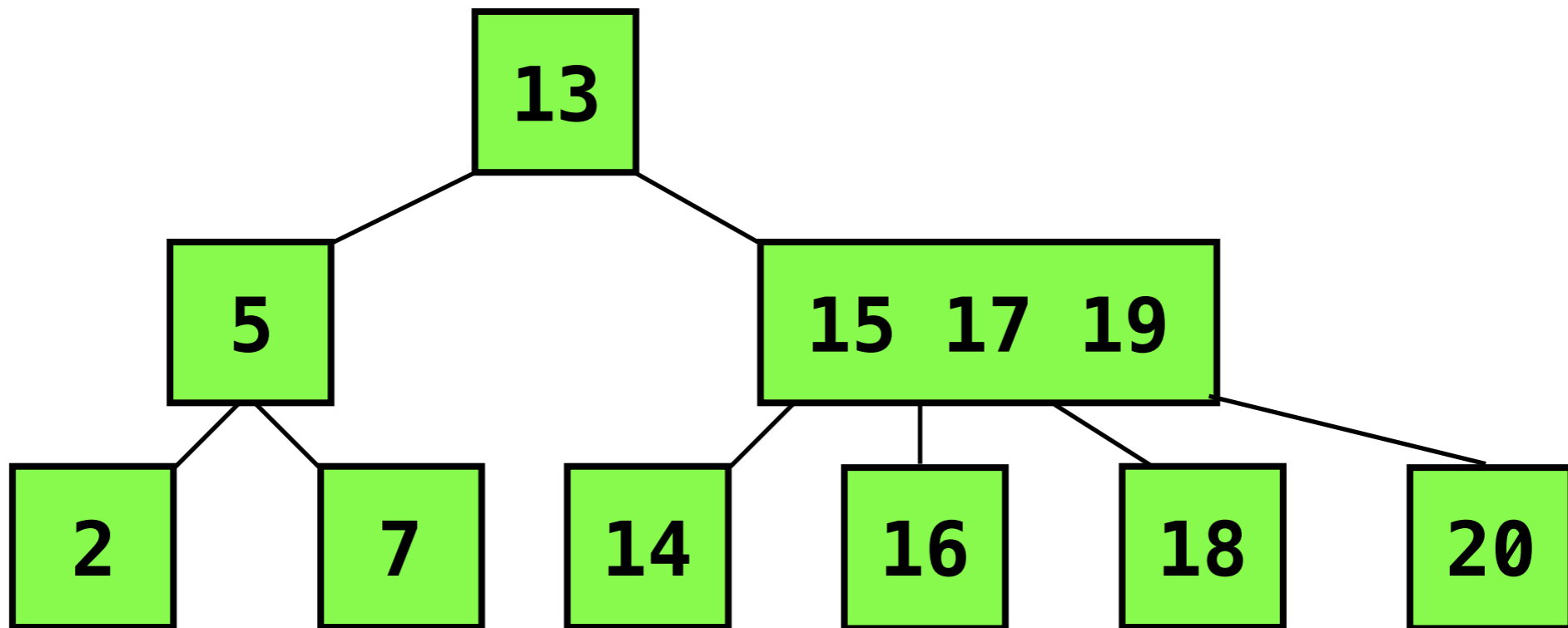
# Add 20



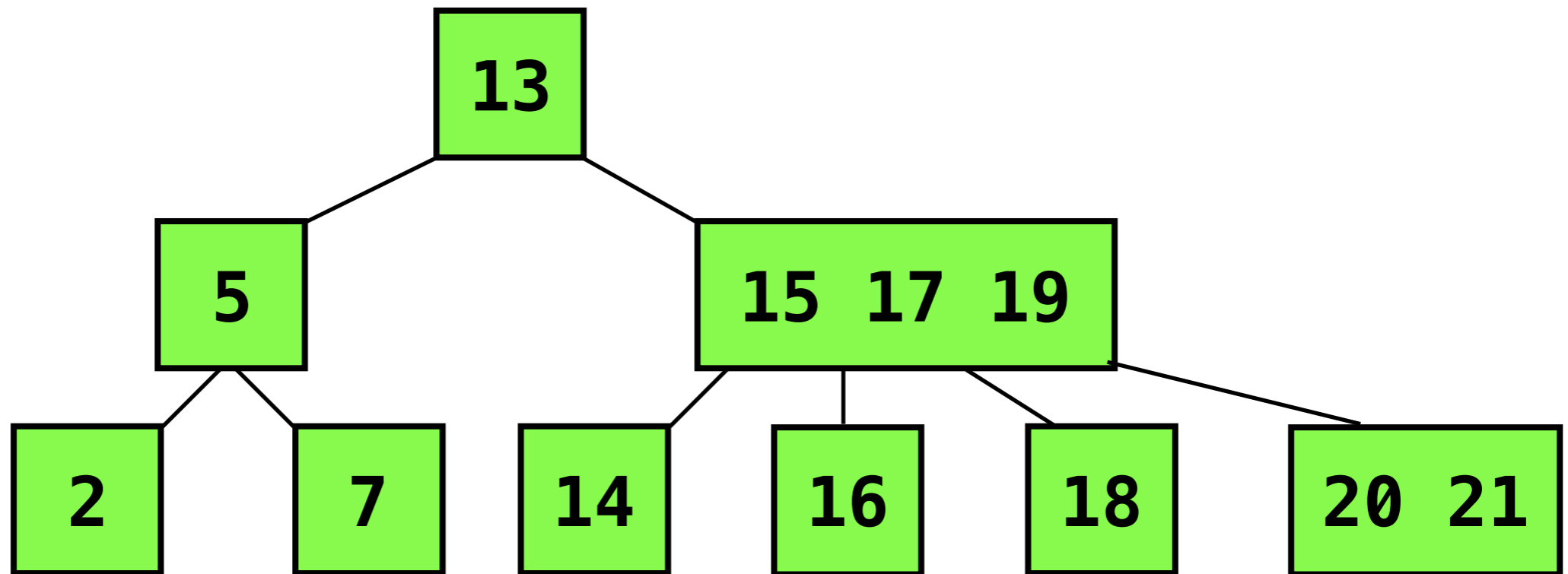
# Add 20



# Add 21

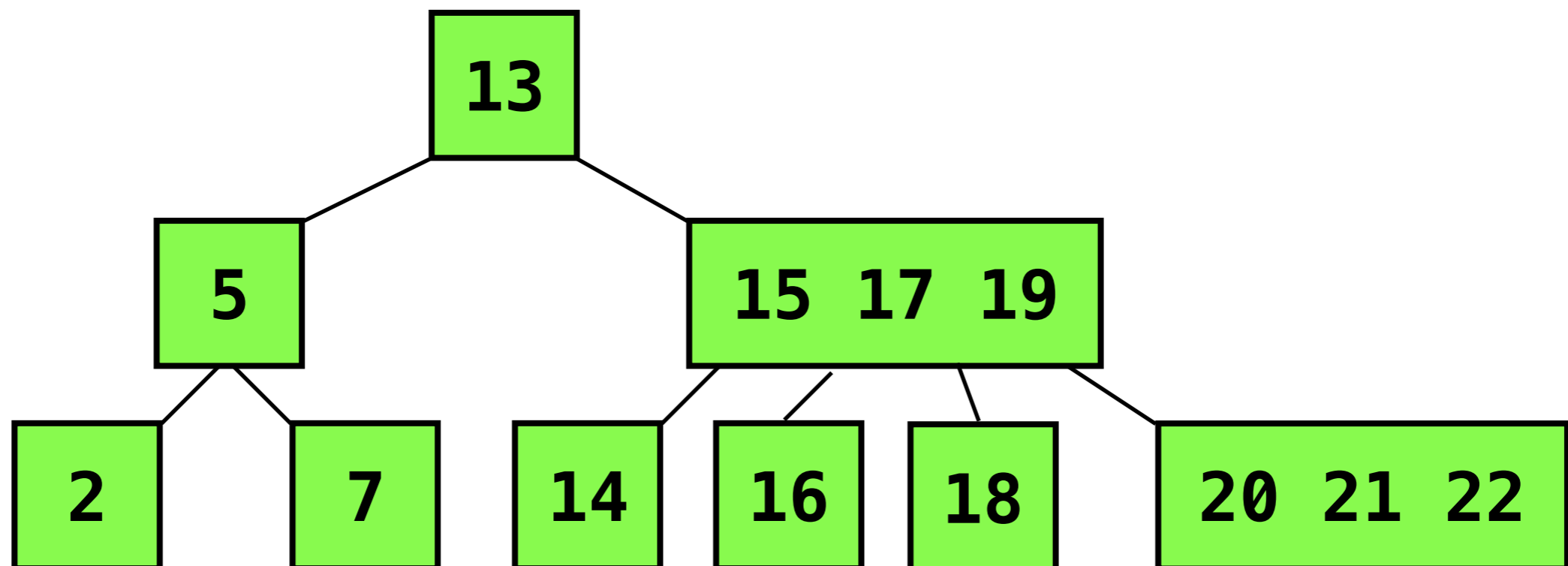


# Add 21



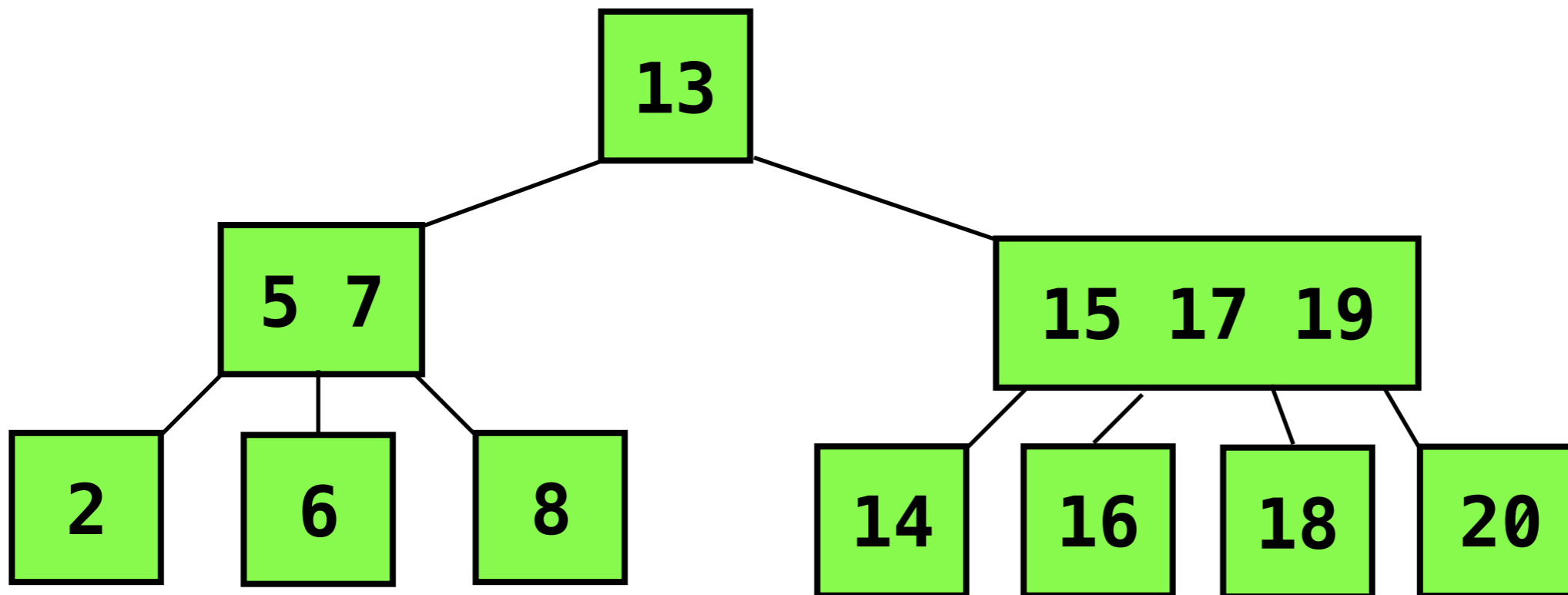
# This is called a B-Tree

- "What, if anything, the *B* stands for has never been established."  
--Wikipedia
- What we were working with before is a type of B-tree called a 2-3-4 tree, since each node can have two, three, or four children.



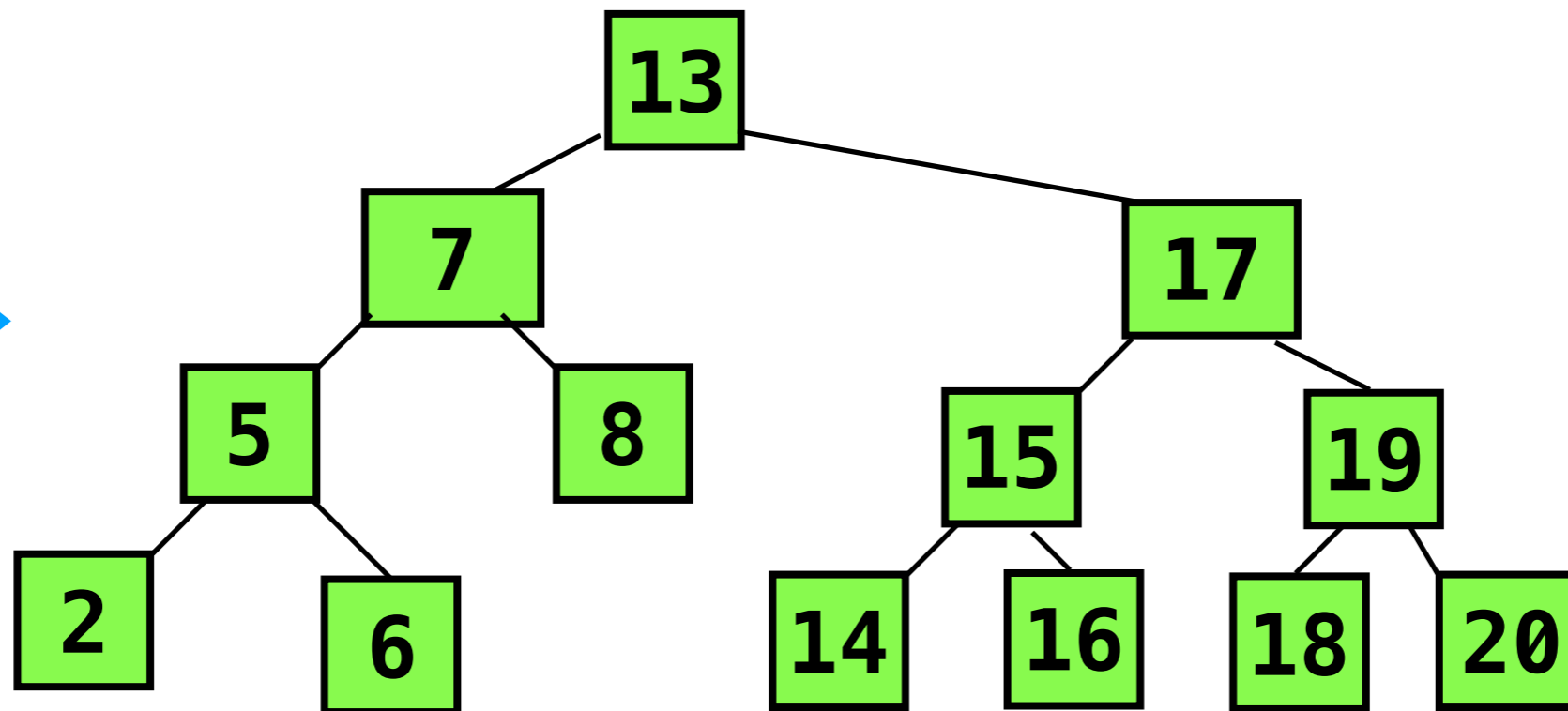
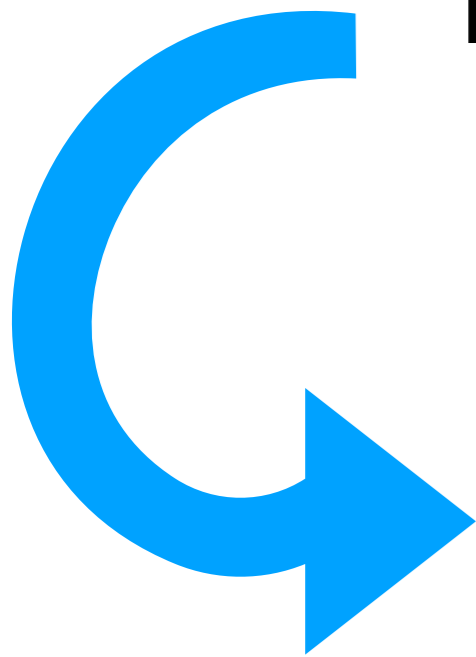
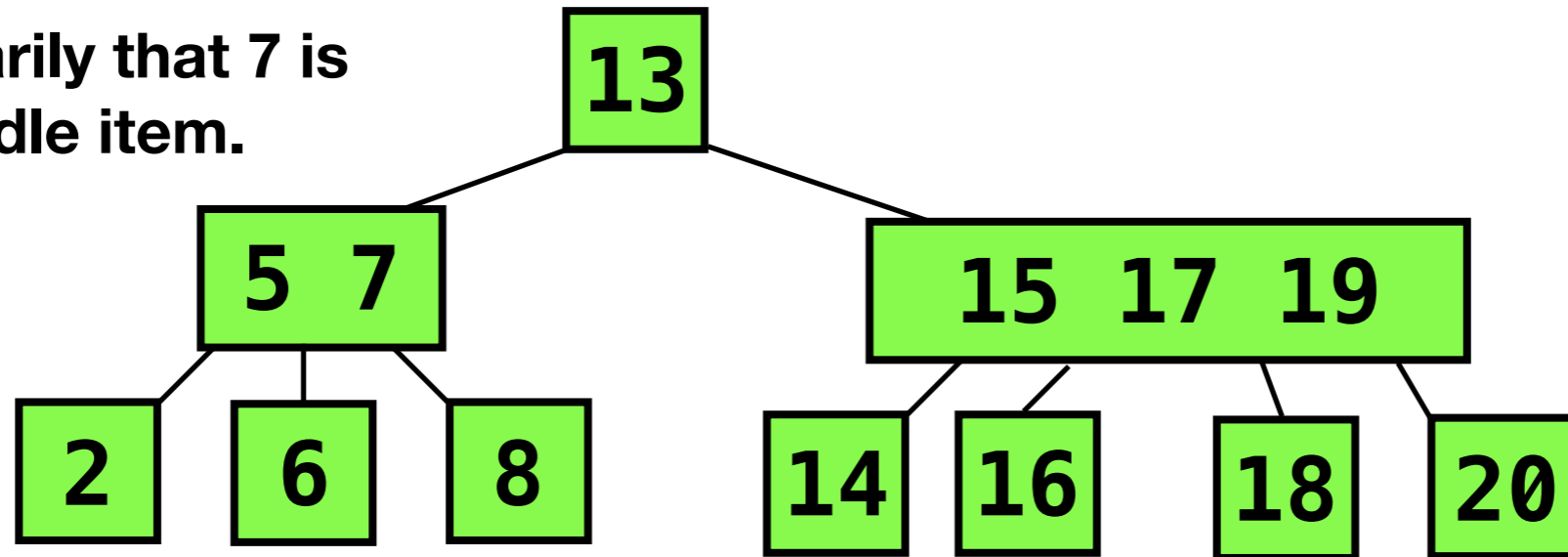
# **B**-Tree limitations

- B-trees are annoying to implement, but give us guaranteed  $\log N$  height.
- Can we convert the 2-3-4 trees into binary trees and simplify our logic?



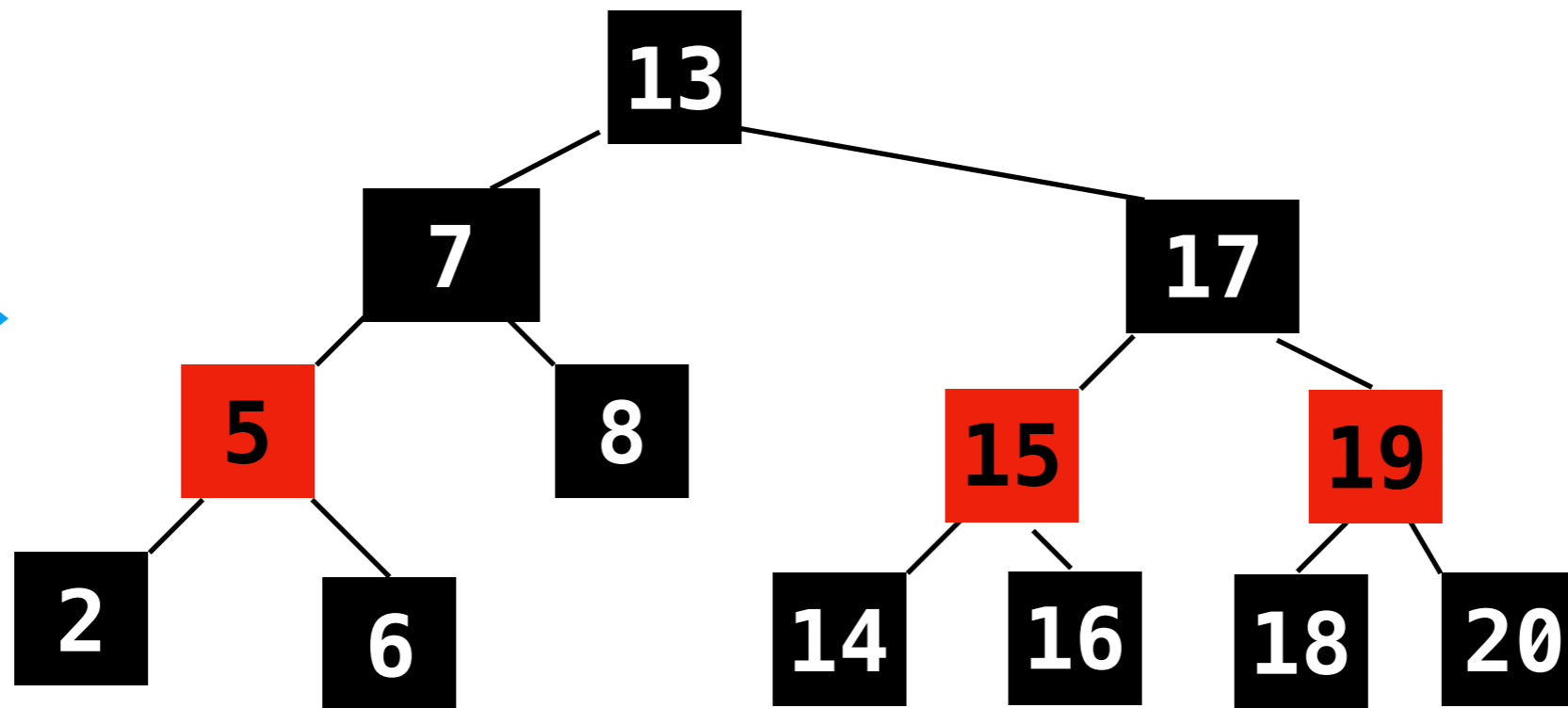
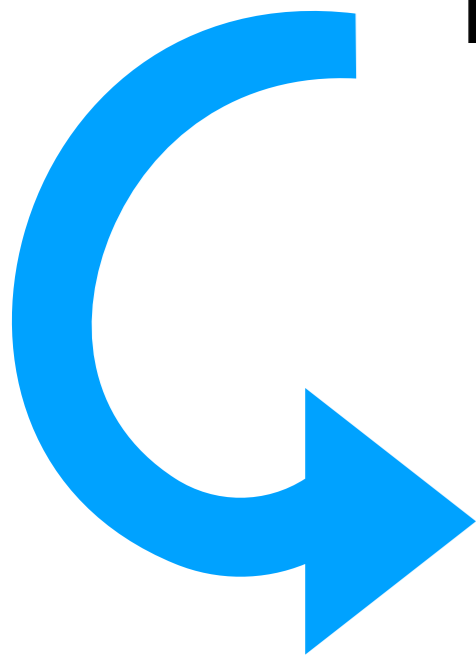
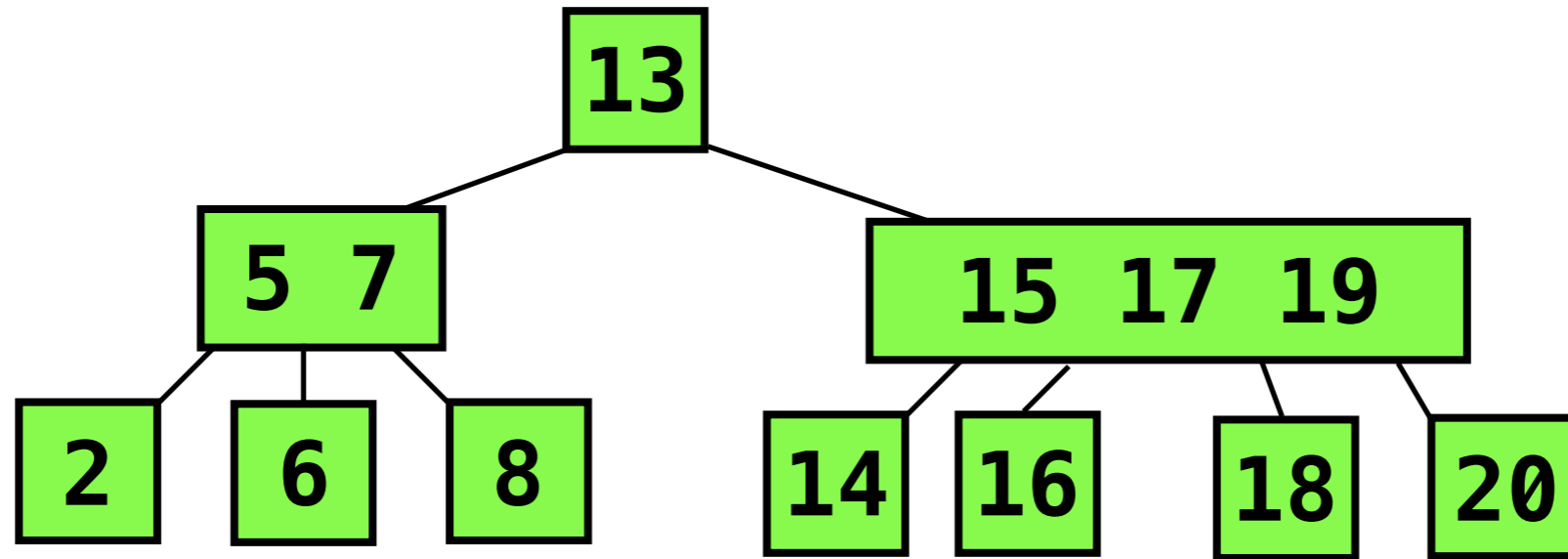
# Idea: "Drop down" the Left and Right Items

Pick arbitrarily that 7 is the middle item.

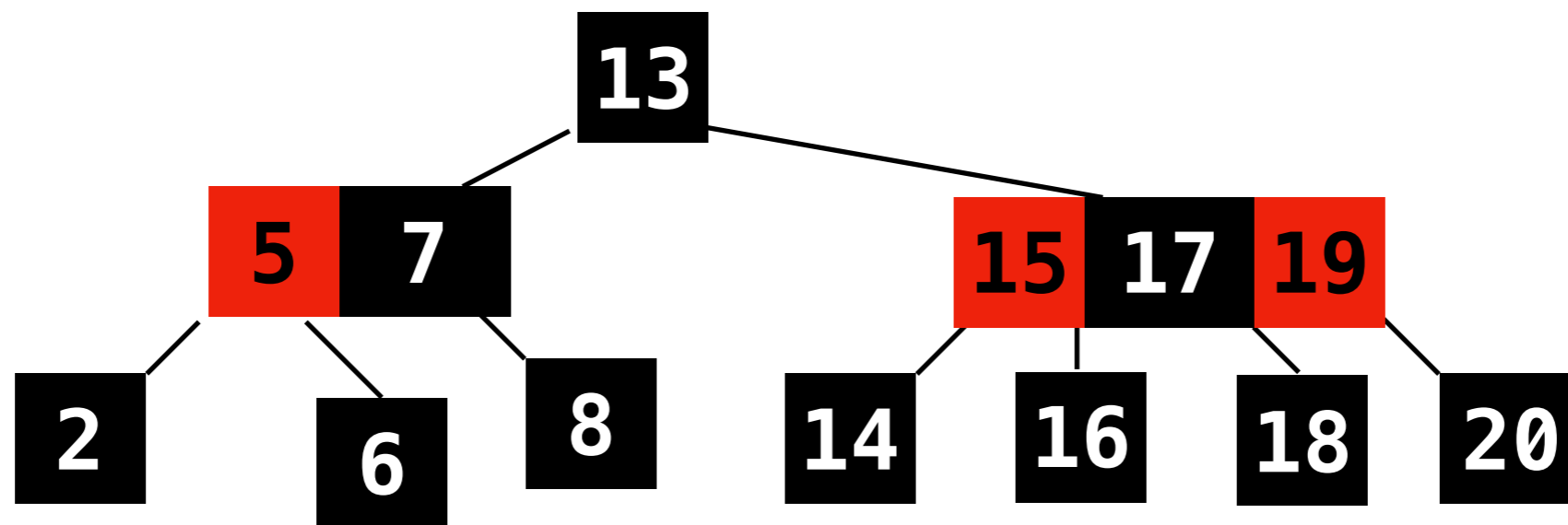
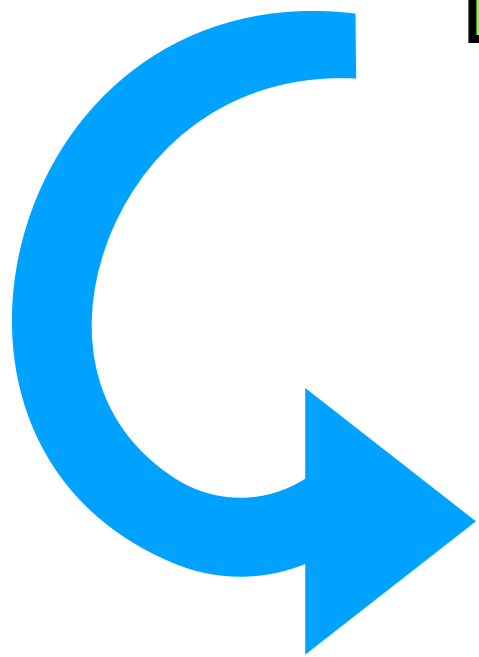
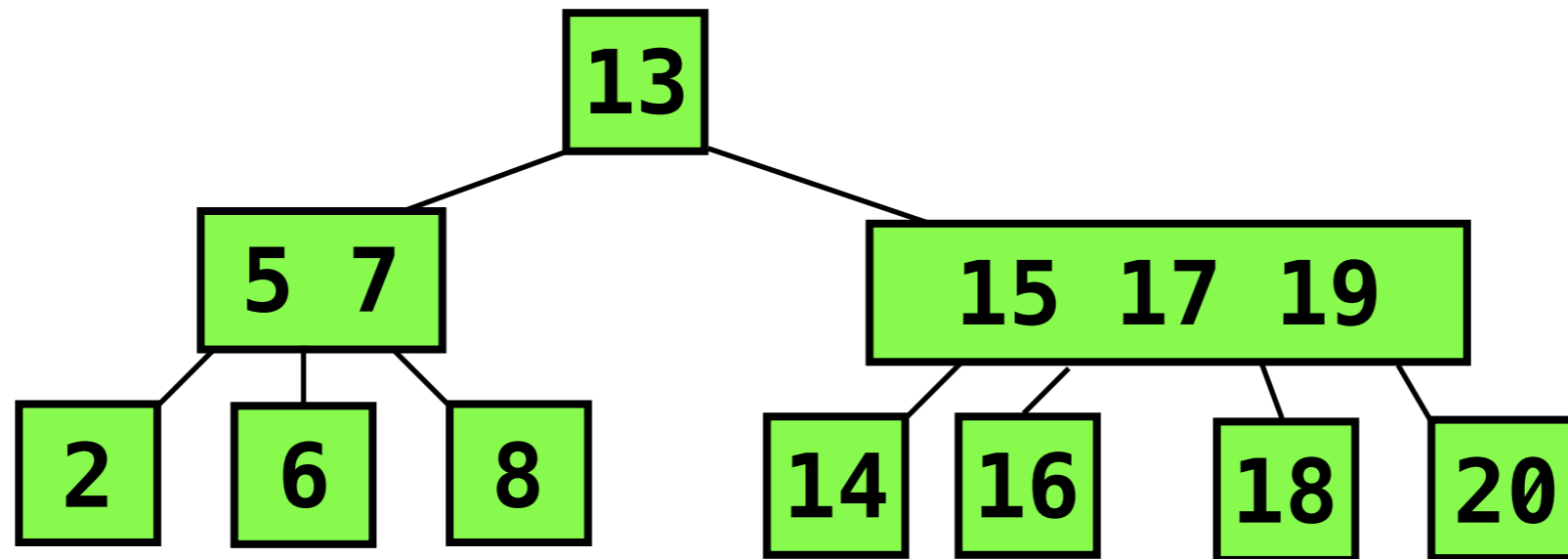




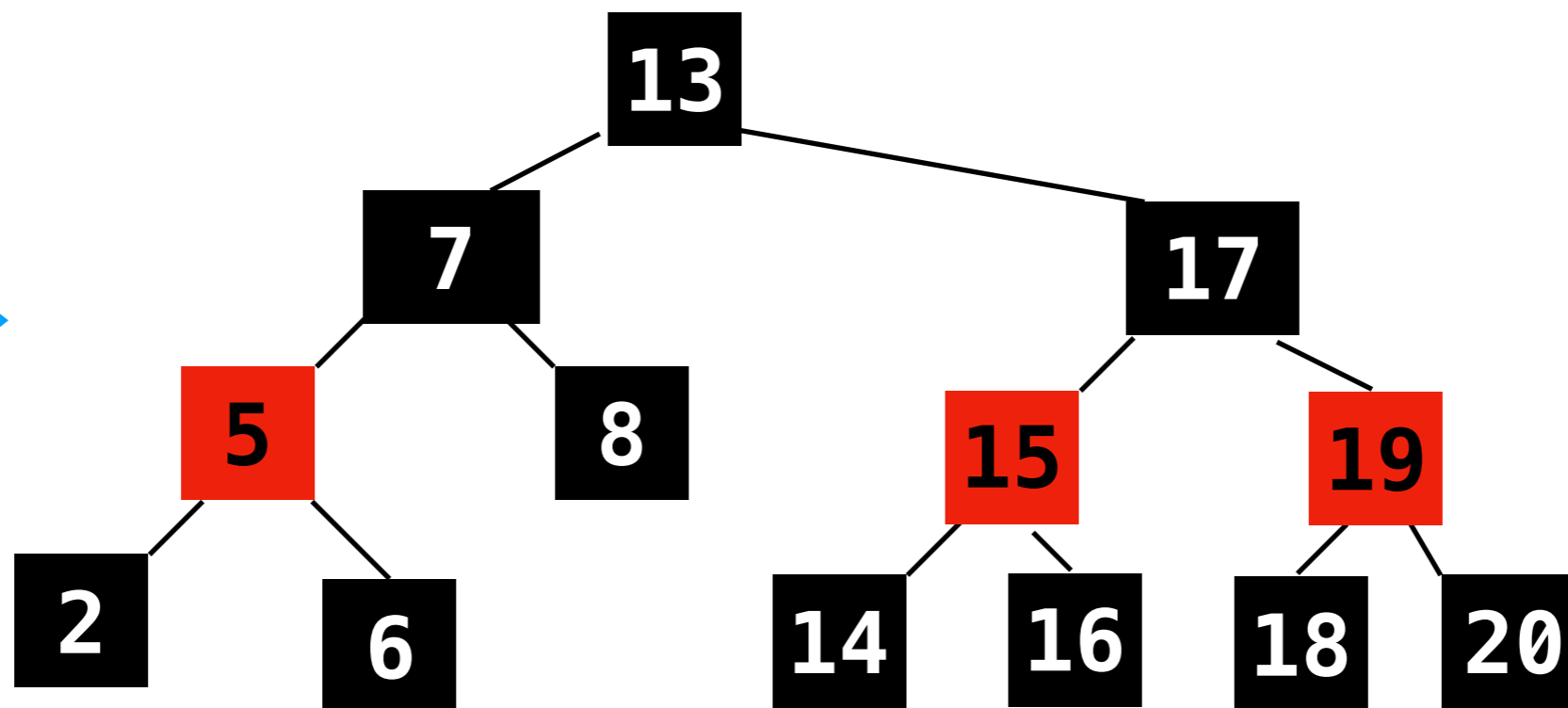
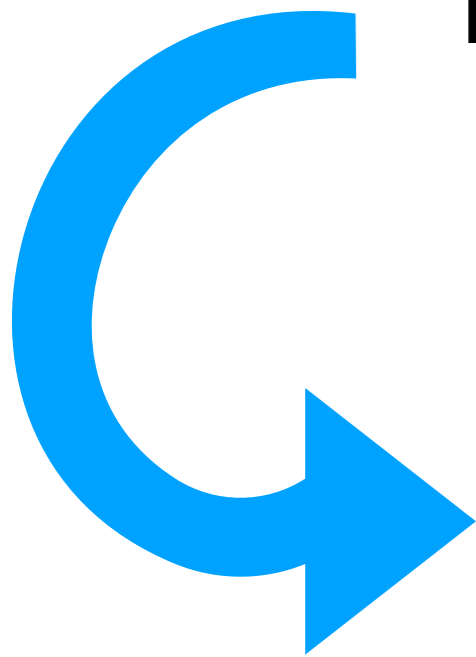
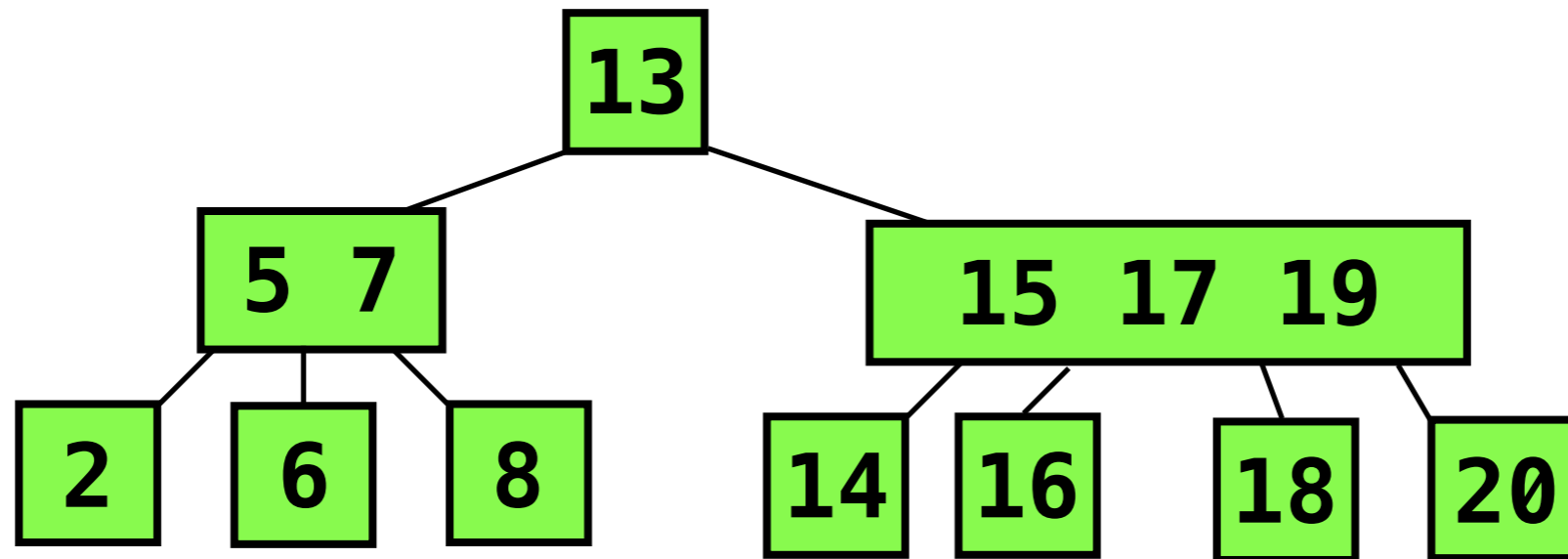
Color nodes that we dropped down red.  
Color everything else black.



To convert back, just move all the red elements up.

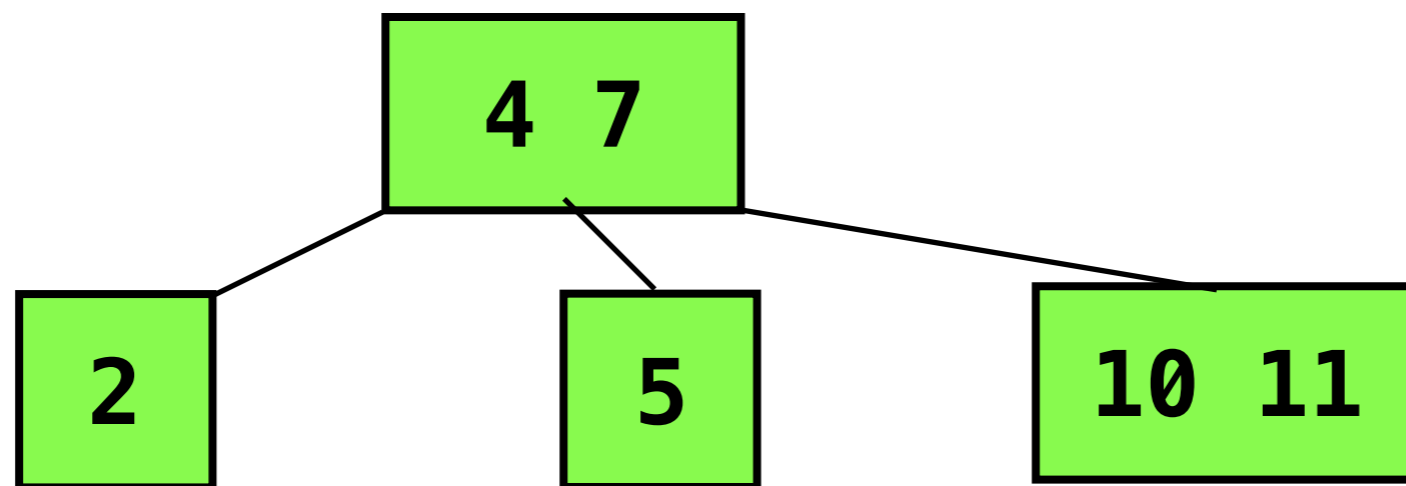


# This is called a "Red-Black Tree"

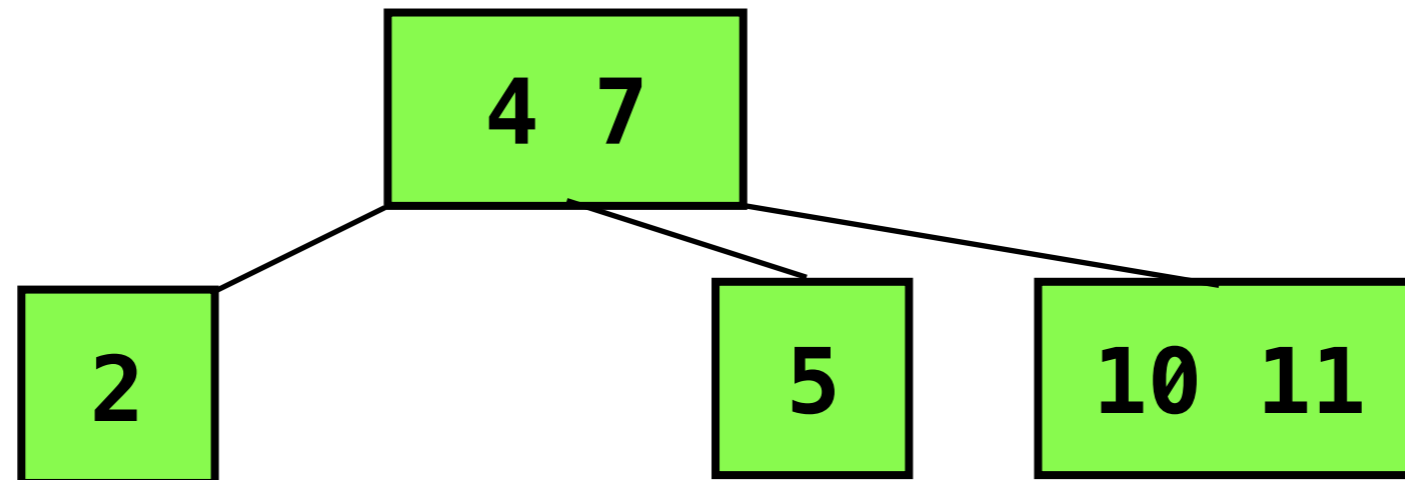


# Now we can have a balanced binary search tree!

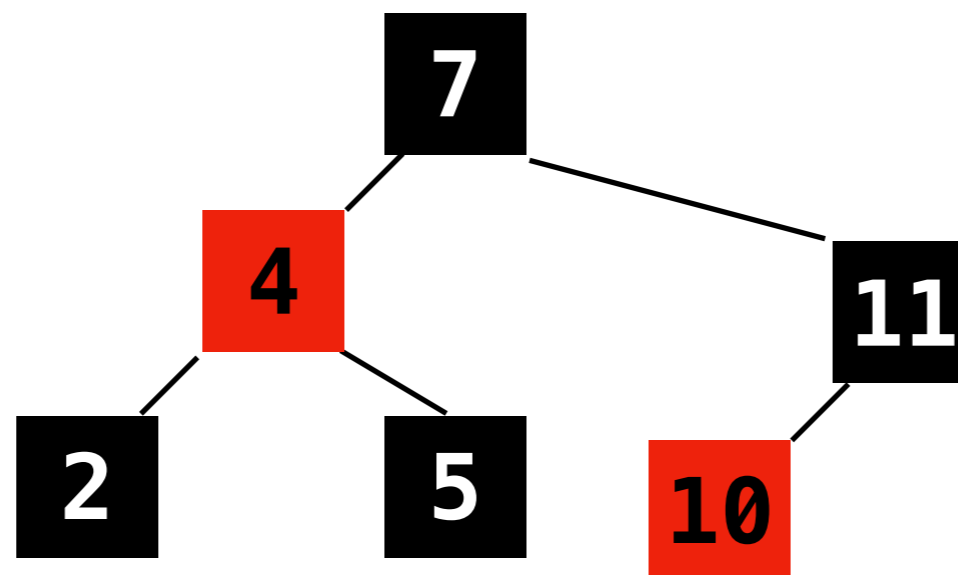
- Just use a red-black tree, and do the equivalent operation that a 2-3-4 tree would do.
- ...But this is hard to implement. Let's simplify further by using a 2-3 tree instead.
  - In a 2-3 tree, a node can have one or two elements
  - And each node has either 2 or 3 children.



Like before, we can make a red-black tree with it.



Always drop the left element down: this creates a left-leaning red-black tree.



# Facts for Left-Leaning Red-Black Trees

- By convention, the root is always black.
- Since we always drop the left element down, a red node *MUST* be the left-child of a black node
- A black node *CANNOT* have two red children since that would correspond to a node with 3 elements

